

# Introduction to Digital Logic

## Lecture 19: State Machines State Machine Analysis

# Characteristic Equations

- With latches and flip-flops we can come up with an equation for the **next value of Q** ( $Q^*$ ) in terms of the **current value of Q** and the **inputs**

S	R	$Q^*$
0	0	Q
1	0	1
0	1	0
1	1	illegal

**Function Table**  
( $Q^*$  listed in terms of Q)

S	R	Q	$Q^*$
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

**Truth Table**  
( $Q^*$  in terms of 0 or 1)

# Characteristic Equations

- With latches and flip-flops we can come up with an equation for the **next value of Q** ( $Q^*$ ) in terms of the **current value of Q** and the **inputs**

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0	0	Q
1	0	1
0	1	0
1	1	illegal

**Function Table**  
( $Q^*$  listed in terms of Q)

S	R	Q	$Q^*$
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	d
1	1	1	d

**Truth Table**  
( $Q^*$  in terms of 0 or 1)

# Characteristic Equations

- For an SR-Latch make a truth table with S,R, and Q and show the next value of Q\*
- The use a K-Map to find an equation for Q\*
- This equation indicates what the next value of Q will be

S	R	Q	Q*
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	d
1	1	1	d

SR		00	01	11	10
		0	0	d	1
Q	0	0	0	d	1
	1	1	0	d	1

$$Q^* = S + R'Q$$

# Characteristic Equations

- For a D-Latch make a truth table with D, and Q and show the next value of  $Q^*$
- You may use a K-Map but we can eyeball it
- This equation indicates what the next value of Q will be

D	Q	$Q^*$
0	0	0
0	1	0
1	0	1
1	1	1

$$Q^* = D$$

# Characteristic Equation

- For a JK-FF make a truth table with J, K, and Q and show the next value of Q\*
- The use a K-Map to find an equation for Q\*
- This equation indicates what the next value of Q will be

J	K	Q	Q*
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	0

		JK			
		00	01	11	10
Q	0	0 <sup>0</sup> 0	2 <sup>2</sup> 0	6 <sup>6</sup> 1	4 <sup>4</sup> 1
	1	1 <sup>1</sup> 1	3 <sup>3</sup> 0	7 <sup>7</sup> 0	5 <sup>5</sup> 1

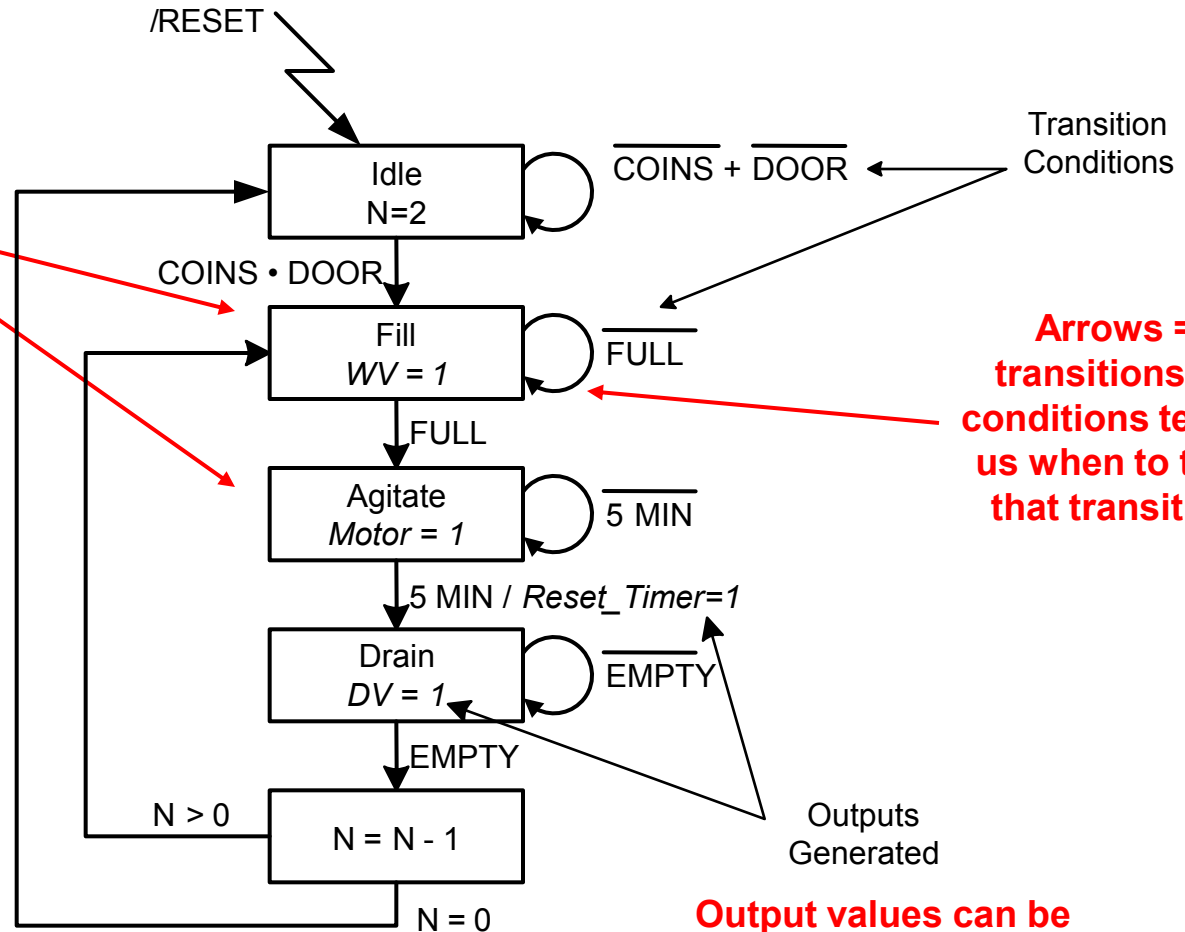
$$Q^* = JQ' + K'Q$$

# State Machines

- Provide the “brains” or control for electronic and electro-mechanical systems
- Implement a set of steps (or algorithm) to control or solve a problem
- **Goal is to generate output values at specific times**
- Combine Sequential and Combinational logic elements
  - Sequential Logic to remember what step (state) we’re in
    - Encodes everything that has happened in the past
  - Combinational Logic to produce outputs and find what state to go to next
    - Generates outputs based on what state we’re in and the input values
- Use state diagrams (a.k.a. flowcharts) to specify the operation of the corresponding state machine

# Washing Machine State Diagram

Boxes represent states or steps in the algorithm

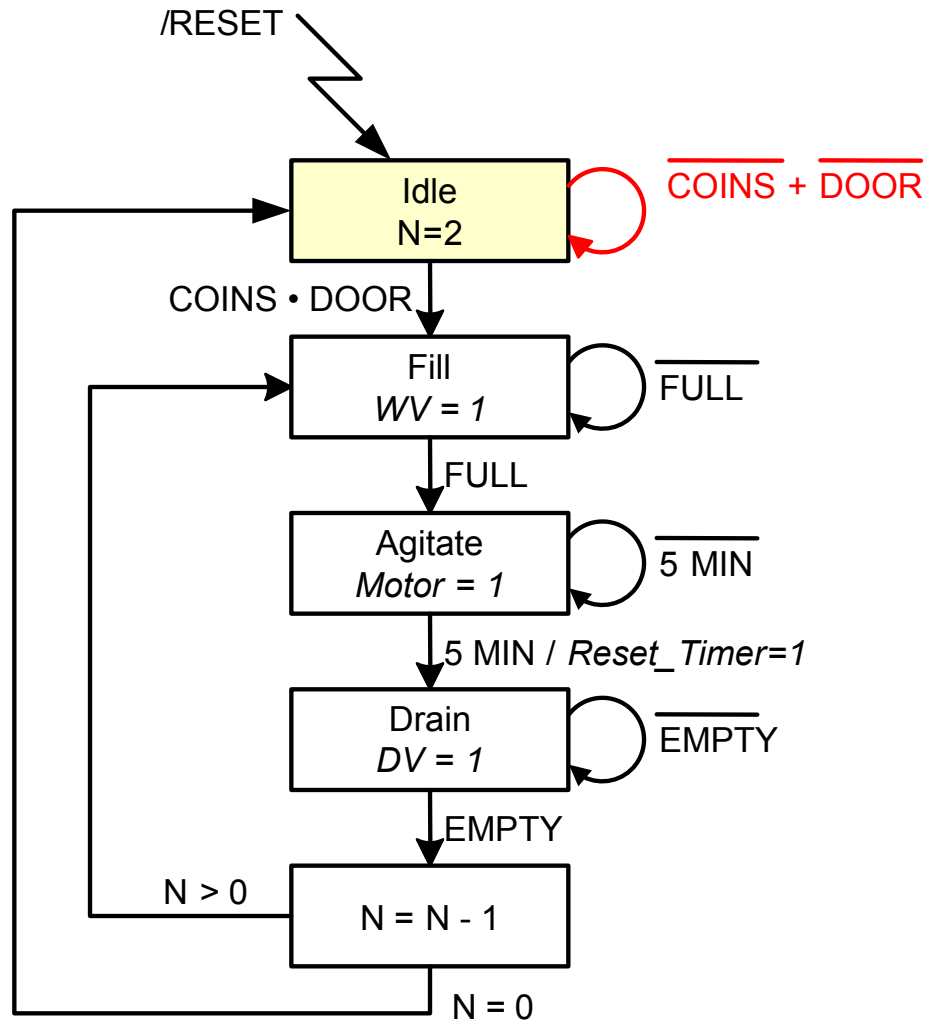


Arrows = transitions w/ conditions telling us when to take that transition

Output values can be associated with a state and transition or just with a state



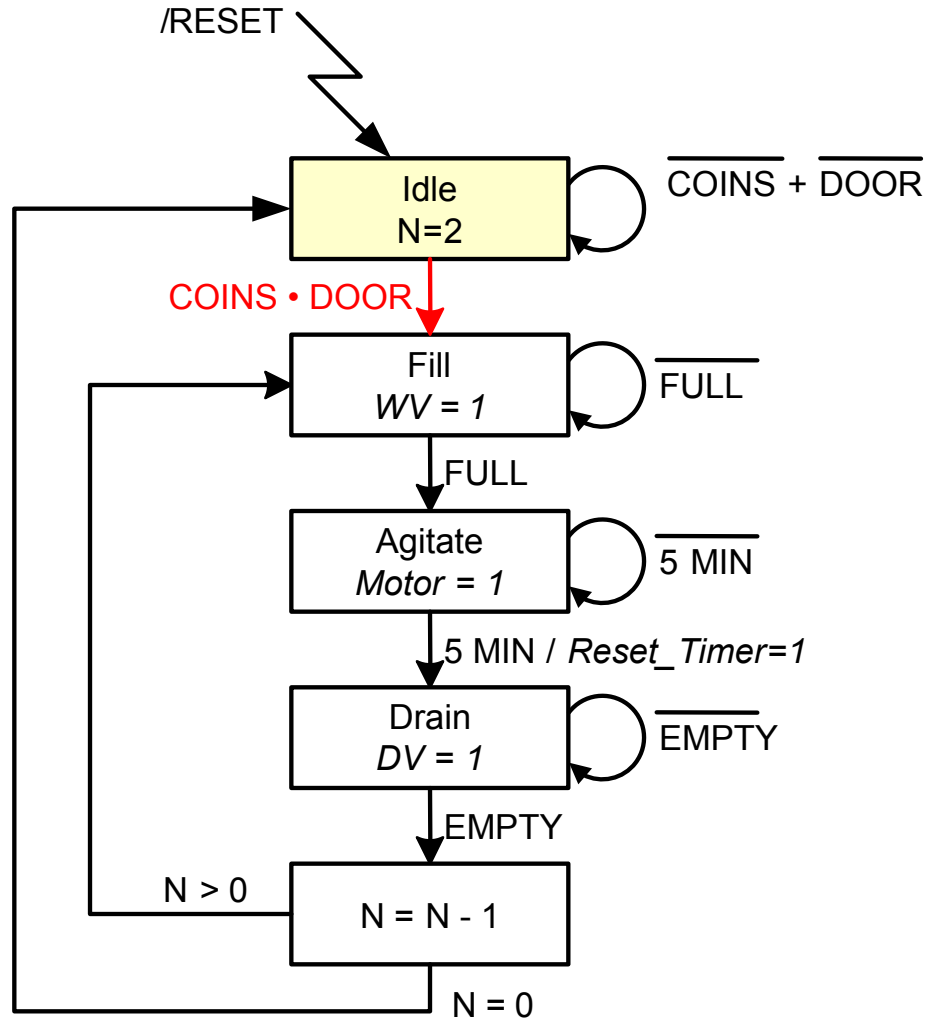
# Washing Machine State Diagram



Stay in the initial state until there is enough money (coins) and the door is closed

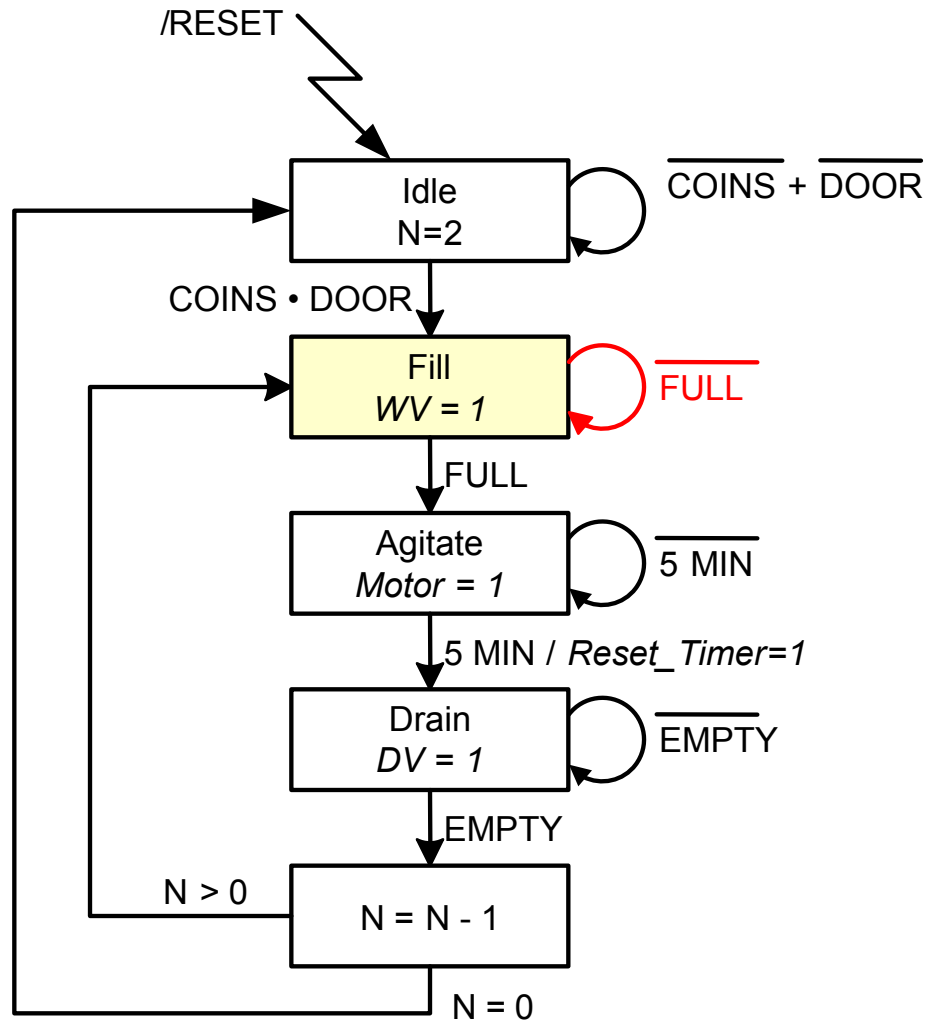
We move through the states based on the conditions. Outputs get asserted when the machine is in that state and the transition is true.

# Washing Machine State Diagram



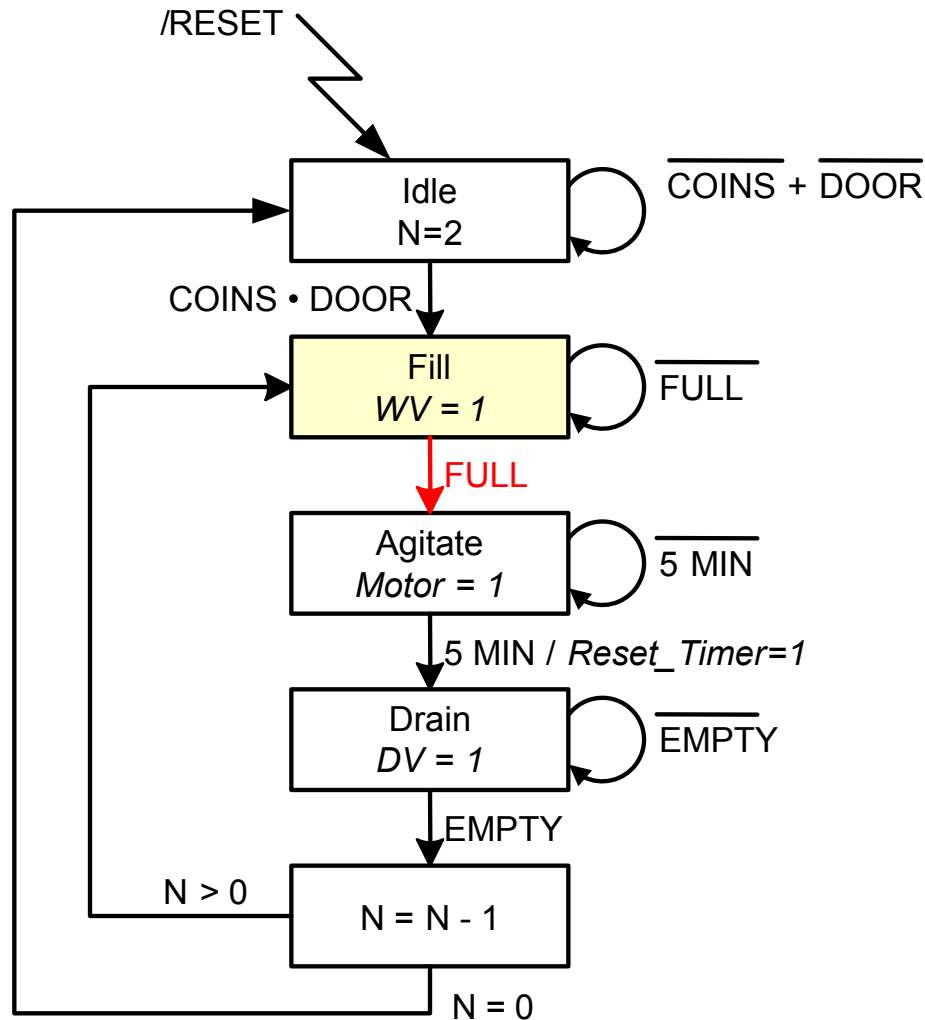
Move to the Fill state when there is enough money (coins) and the door is closed

# Washing Machine State Diagram



Stay in the Fill state until it is full...also set the Water Valve Open output to be true

# Washing Machine State Diagram



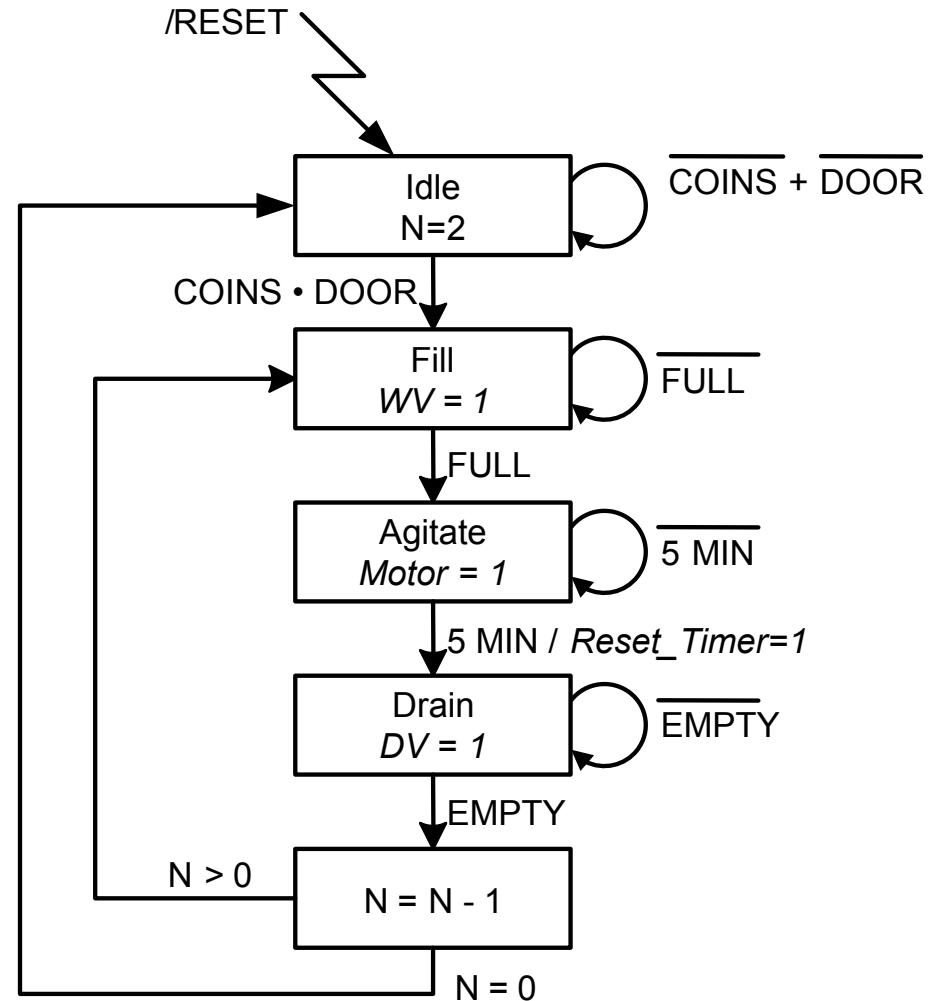
Move to the Agitate state after it is full

# State Machines

- Use sequential and combinational logic to implement a set of steps (i.e. an algorithm)
- Goal is to produce outputs at specific points of time
  - Combinational logic alone cannot do that because the outputs will change as soon as the inputs change (no notion of time)

# State Diagrams

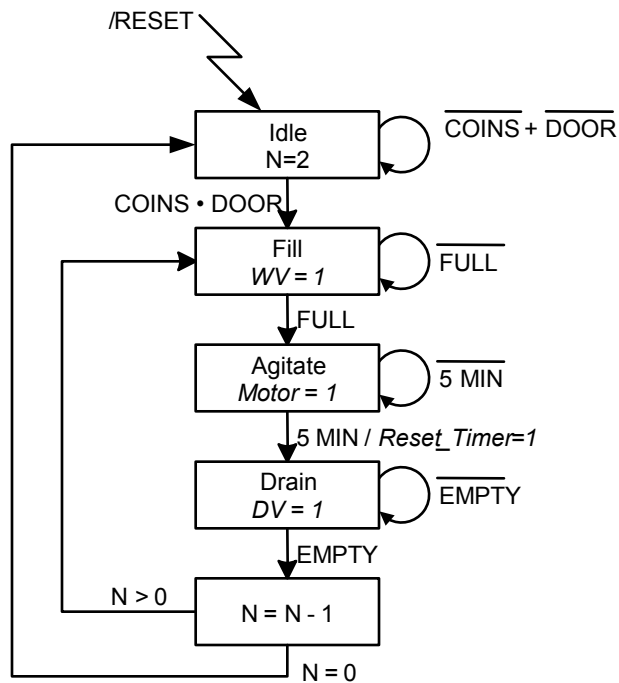
- Used to show operation or function of a state machine
- Like a flowchart but called a state diagram
- 3 parts
  - States
  - Transitions
  - Outputs



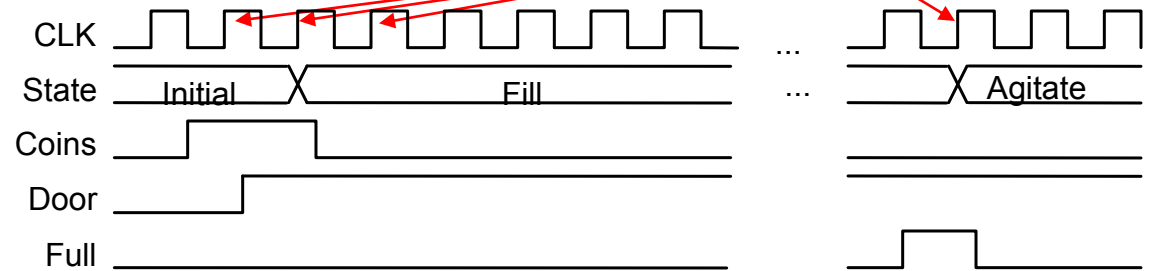
**State Diagram for a Washing Machine**

# State Diagrams

- One transition is made at each clock edge
  - Based on the current state and the conditions associated w/ the transitions



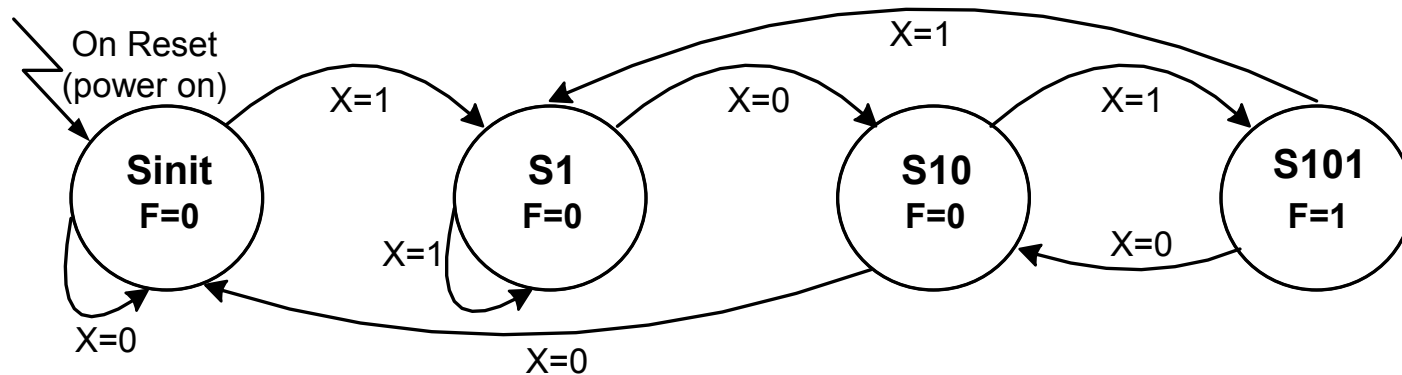
The transition conditions are checked at the edge of the clock



**State Diagram for a Washing Machine**

# Another State Diagram Example

- “101” Sequence Detector should output  $F=1$  when the sequence 101 is found in consecutive order

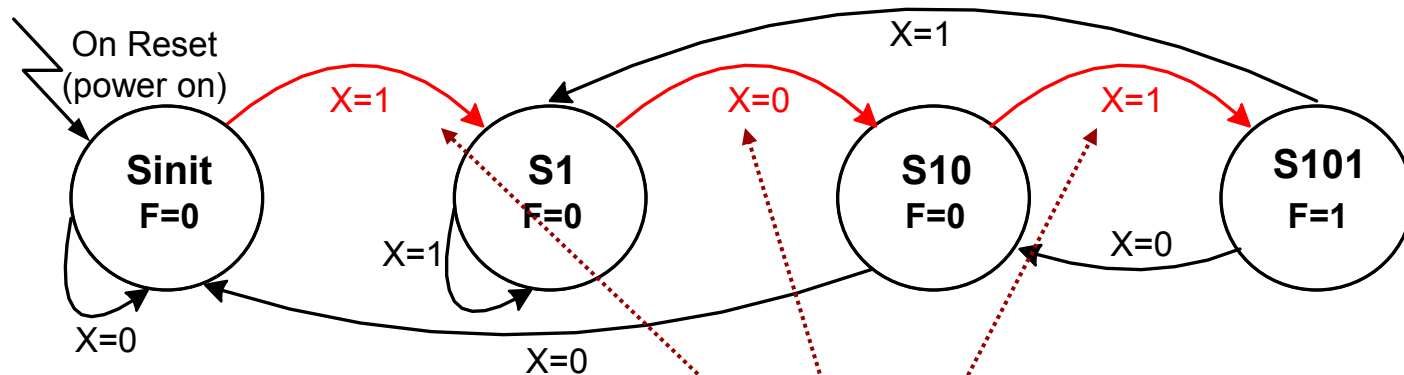


**State Diagram for “101”  
Sequence Detector**



# Another State Diagram Example

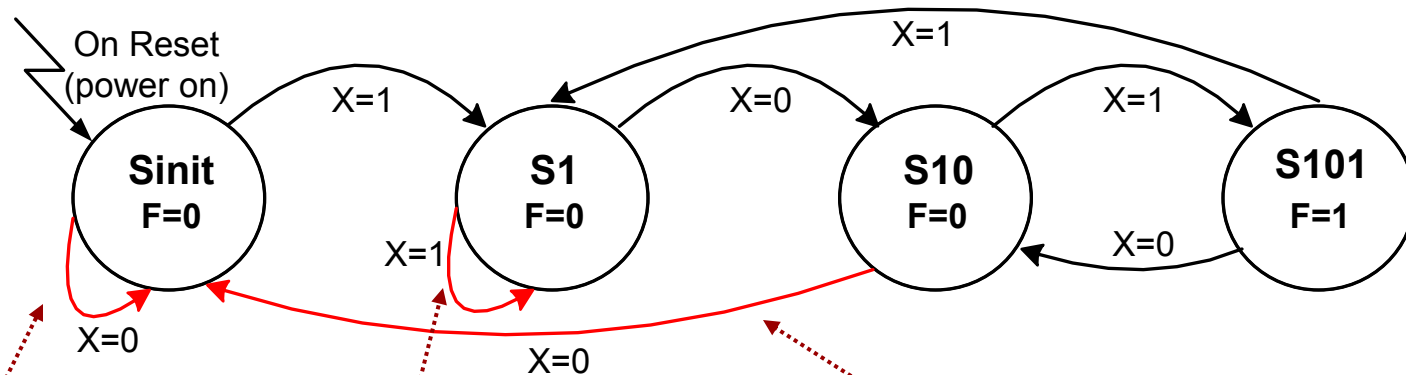
- “101” Sequence Detector should output  $F=1$  when the sequence 101 is found in consecutive order



**We have to remember the 1,0,1 along the way**

# Another State Diagram Example

- “101” Sequence Detector should output  $F=1$  when the sequence 101 is found in consecutive order



A '0' initially is not part of the sequence so stay in Sinit

Another '1' in S1 means you have 11, but that second '1' can be the start of the sequence

A '0' in S10 means you have 100 which can't be part of the sequence

# State Machines

- State Machines can be broken into 3 sections of logic
  - State Memory (SM)
    - Just FF's to remember the *current state*
  - Next State Logic (NSL)
    - Combo logic to determine the next state
    - Essentially implements the transition conditions
  - Output Function Logic (OFL)
    - Combo logic to produce the outputs

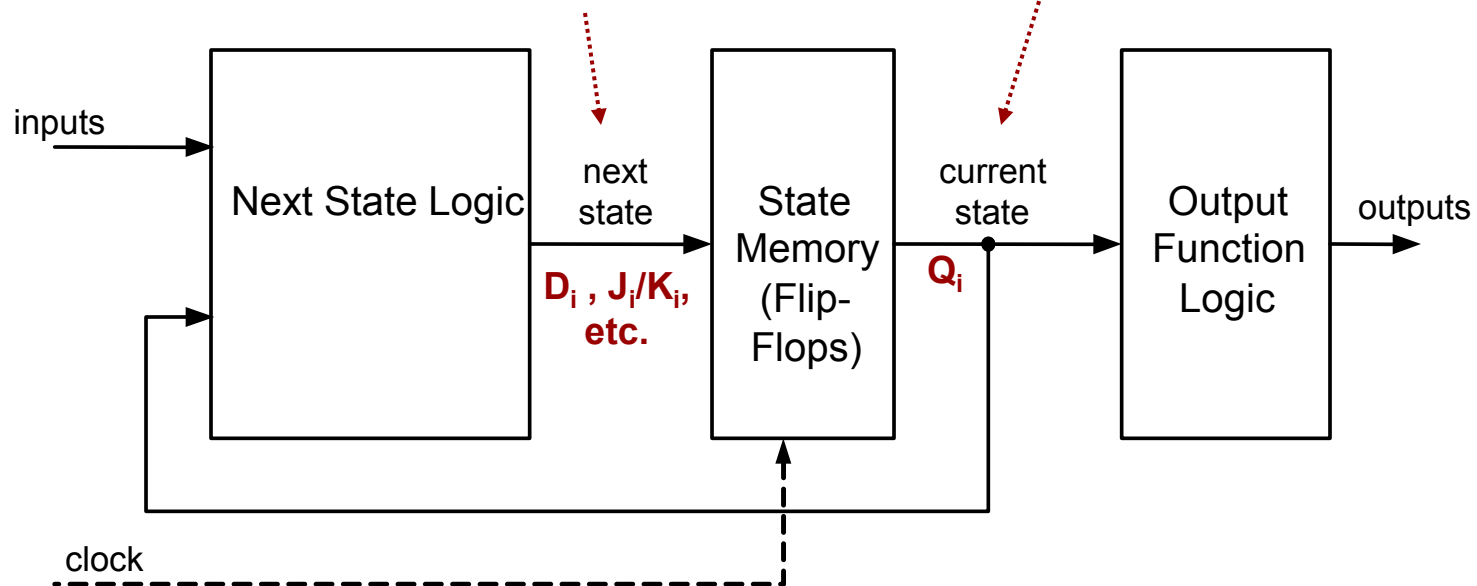
# State Machine

## NEXT STATE

The FF inputs will be the value of the next state (on the next clock edge the FF outputs will change based on the inputs)

## CURRENT STATE

The FF outputs represent the current state (the state we're in right now)

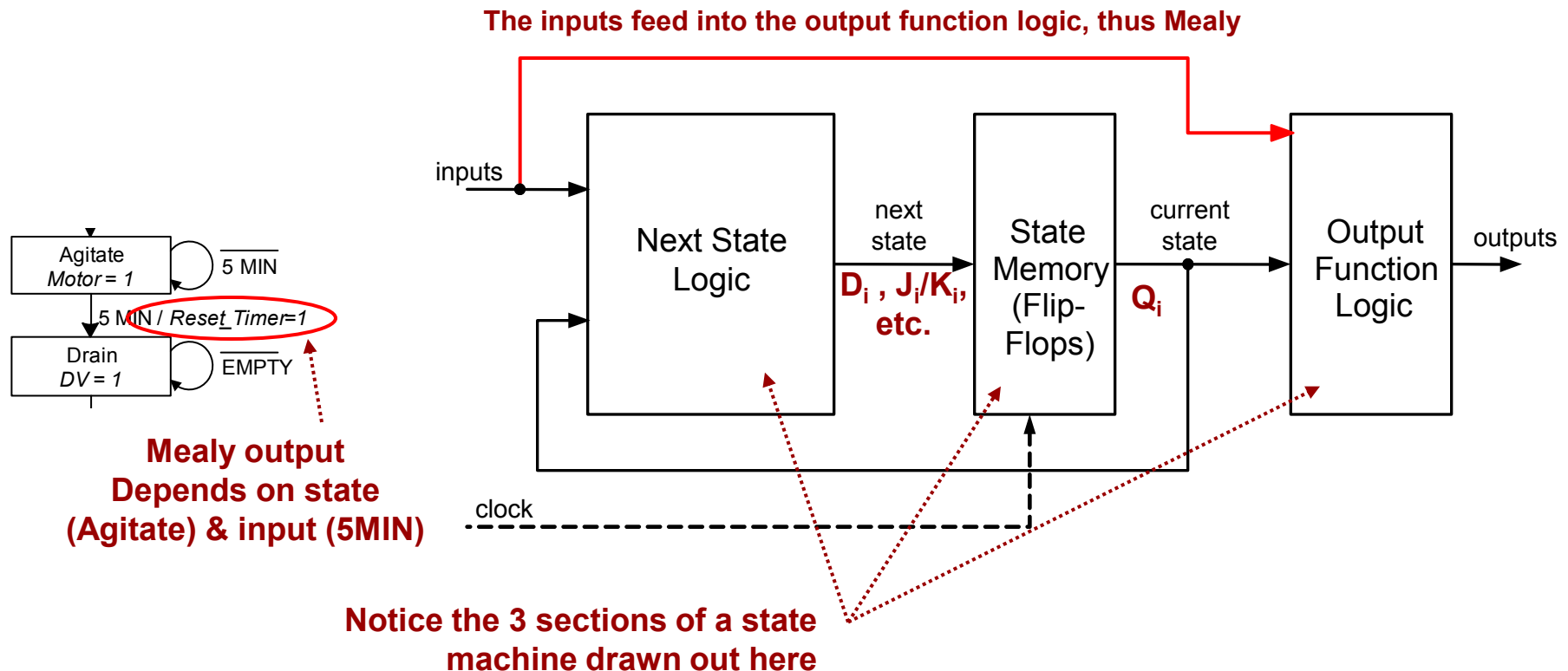


# State Machines

- State Machines can be classified according to how the outputs are produced
  - If  $Outputs = f(current\ state, external\ inputs)...$   
**MEALY Machine**
  - If  $Outputs = f(current\ state)...$   
**MOORE Machine**

# Mealy Machine

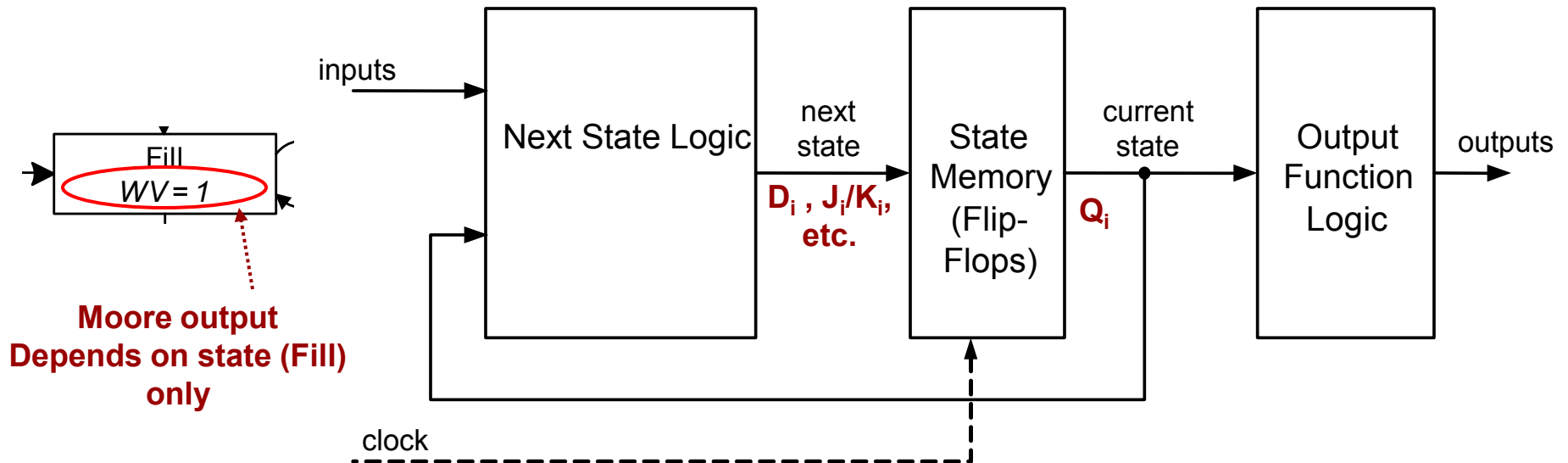
- In a Mealy Machine the outputs depend not only on the current state but the external inputs



# Moore Machine

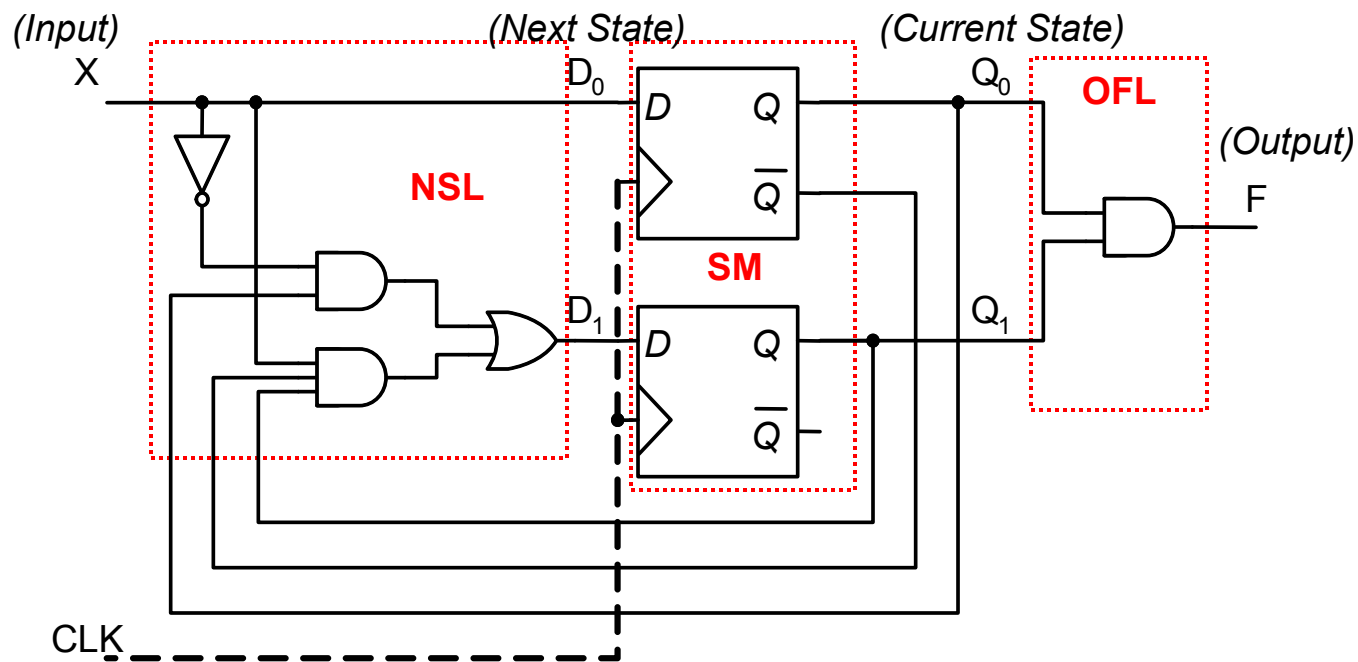
- In a Moore Machine the outputs only depend on the current state

The inputs do not feed into the OFL, thus Moore Machine



# State Machines

- Below is a circuit implementing a state machine, notice how it breaks into the 3 sections



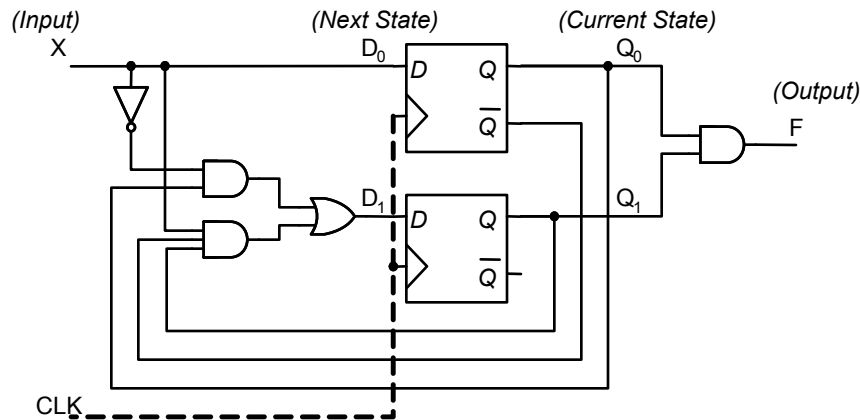


# State Machine Analysis

- In state machine analysis our goal is to take a circuit and figure out the state diagram that it implements

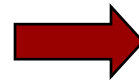
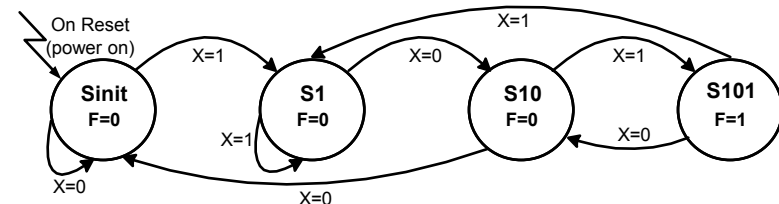
Convert...

Circuit



to

State Diagram



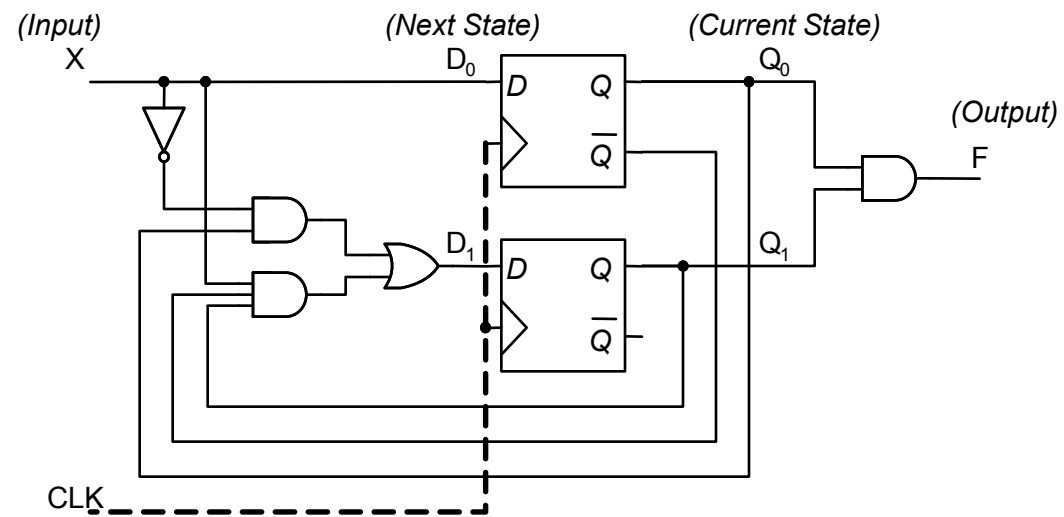
# State Machine Analysis

- 6 Steps to analyze
  - Excitation Equations
    - (eqn's for FF inputs)
  - Transition Equations
    - (use characteristic equation of FF and substitute excitation equations for the FF inputs)
  - Output Equations
  - Transition/Output Table
    - Make a table showing all combinations of current state and external inputs and then what each of the next state and output values will be for each of those combinations
  - State Name Assignment
    - (Symbolic names replace binary codes)
  - Draw the State Diagram

# Excitation Equations

Write equations for  
all FF inputs

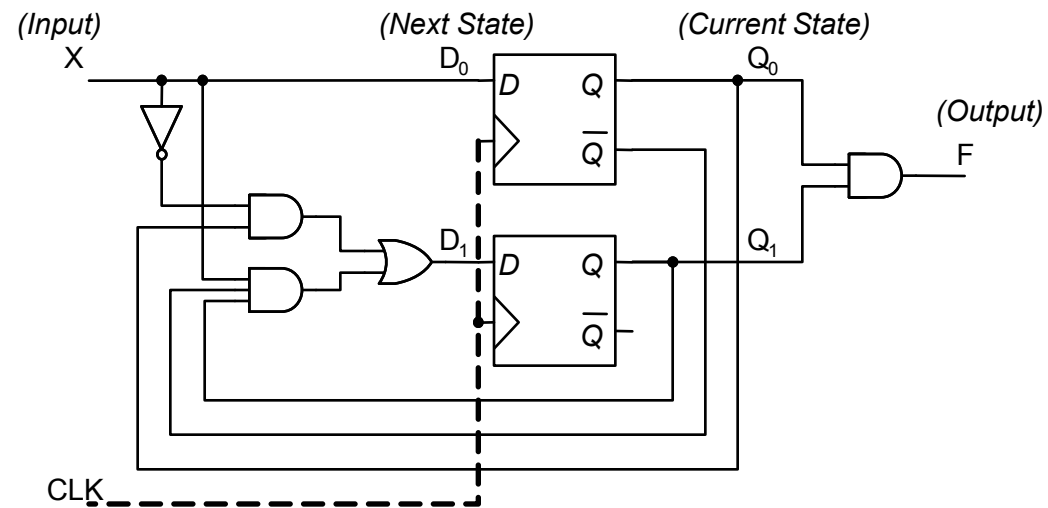
- $D_0 = X$
- $D_1 = X'Q_0 + XQ_1Q_0'$



# Transition Equations

Come up with equations for the next value of the Q's (use characteristic equation  $Q^* = D$ )

- $Q_0^* = D_0 = X$
- $Q_1^* = D_1 = X'Q_0 + XQ_1Q_0'$

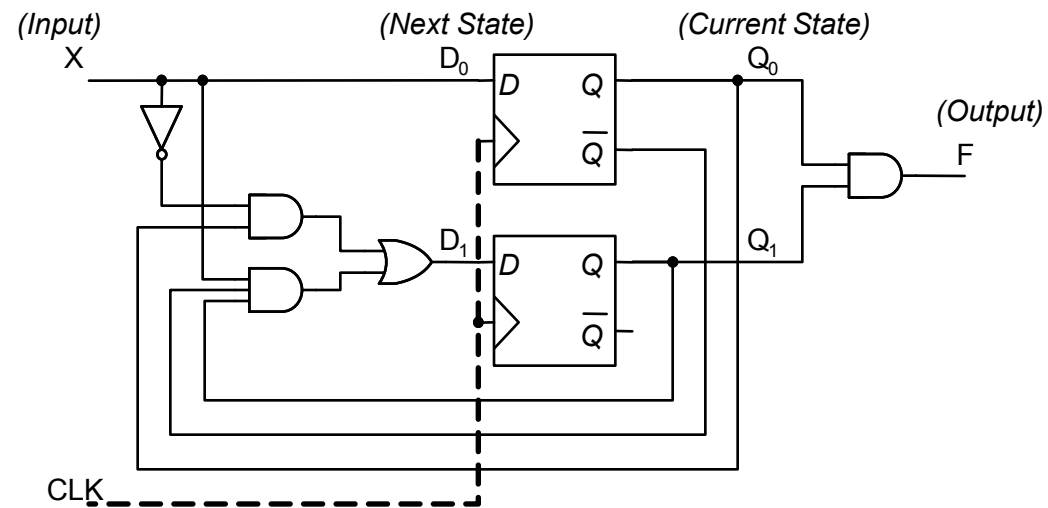


# Output Equations

Equations for all outputs

- $F = Q_1Q_0$

Notice that this is a Moore Machine since the outputs depend on only the current state



# Transition/Output Table

- Make a table of the current state, next state, and outputs
- Use the previous equations to fill out the transition output table

Current State		Next State				Outputs
		X = 0		X = 1		
Q1	Q0	Q1*	Q0*	Q1*	Q0*	F
0	0					
0	1					
1	0					
1	1					

$$Q_0^* = X$$

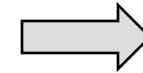
$$Q_1^* = X'Q_0 + XQ_1Q_0'$$

$$F = Q_1Q_0$$

Because this is a Moore Machine, we can make a separate column for F apart from the values of X

# Table Format

Current State		Next State				Outputs
		X = 0		X = 1		
Q1	Q0	Q1*	Q0*	Q1*	Q0*	F
0	0					
0	1					
1	0					
1	1					



X	Q1	Q0	Q1*	Q0*
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

# Transition/Output Table

- Make a table of the current state, next state, and outputs
- Use the previous equations to fill out the transition output table

Current State		Next State				Outputs
		X = 0		X = 1		
Q1	Q0	Q1*	Q0*	Q1*	Q0*	F
0	0		0		1	
0	1		0		1	
1	0		0		1	
1	1		0		1	

$$Q_0^* = X$$

$$Q_1^* = X'Q_0 + XQ_1Q_0'$$

$$F = Q_1Q_0$$



# Transition/Output Table

- Make a table of the current state, next state, and outputs
- Use the previous equations to fill out the transition output table

Current State		Next State				Outputs
		X = 0		X = 1		
Q1	Q0	Q1*	Q0*	Q1*	Q0*	F
0	0	0	0	0	1	
0	1	1	0	0	1	
1	0	0	0	1	1	
1	1	1	0	0	1	

$$Q_0^* = X$$

$$Q_1^* = X'Q_0 + XQ_1Q_0'$$

$$F = Q_1Q_0$$

# Transition/Output Table

- Make a table of the current state, next state, and outputs
- Use the previous equations to fill out the transition output table

Current State		Next State				Outputs
		X = 0		X = 1		
Q1	Q0	Q1*	Q0*	Q1*	Q0*	F
0	0	0	0	0	1	0
0	1	1	0	0	1	0
1	0	0	0	1	1	0
1	1	1	0	0	1	1

$$Q_0^* = X$$

$$Q_1^* = X'Q_0 + XQ_1Q_0'$$

$$F = Q_1Q_0$$

# Transition/Output Table

- Make a table of the current state, next state, and outputs
- Use the previous equations to fill out the transition output table

Current State		Next State				Outputs
		X = 0		X = 1		
Q1	Q0	Q1*	Q0*	Q1*	Q0*	F
0	0	0	0	0	1	0
0	1	1	0	0	1	0
1	0	0	0	1	1	0
1	1	1	0	0	1	1

If current state  $(Q_1, Q_0) = 00$  and  $X = 0$  then we'll stay in state 00

If current state  $(Q_1, Q_0) = 00$  and  $X = 1$  then we'll go to state 01

# State Name Assignment

- Just give a symbolic name to each state (i.e. 00 = SA, 01 = SB, etc.)
- In this case use the following names...

Current State			Next State						Output
			X = 0			X = 1			
Q <sub>1</sub>	Q <sub>0</sub>	State	Q <sub>1</sub> *	Q <sub>0</sub> *	State	Q <sub>1</sub> *	Q <sub>0</sub> *	State	F
0	0	Sinit	0	0	Sinit	0	1	S1	0
0	1	S1	1	0	S10	0	1	S1	0
1	0	S10	0	0	Sinit	1	1	S101	0
1	1	S101	1	0	S10	0	1	S1	1

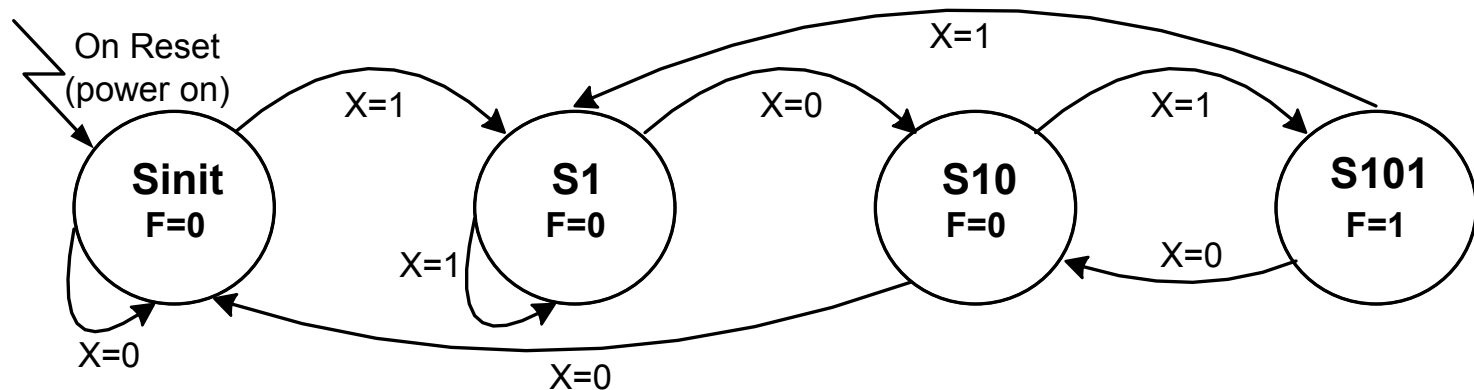
**We call state 01 = S1**

**Replace 01 with S1**

# State Diagram

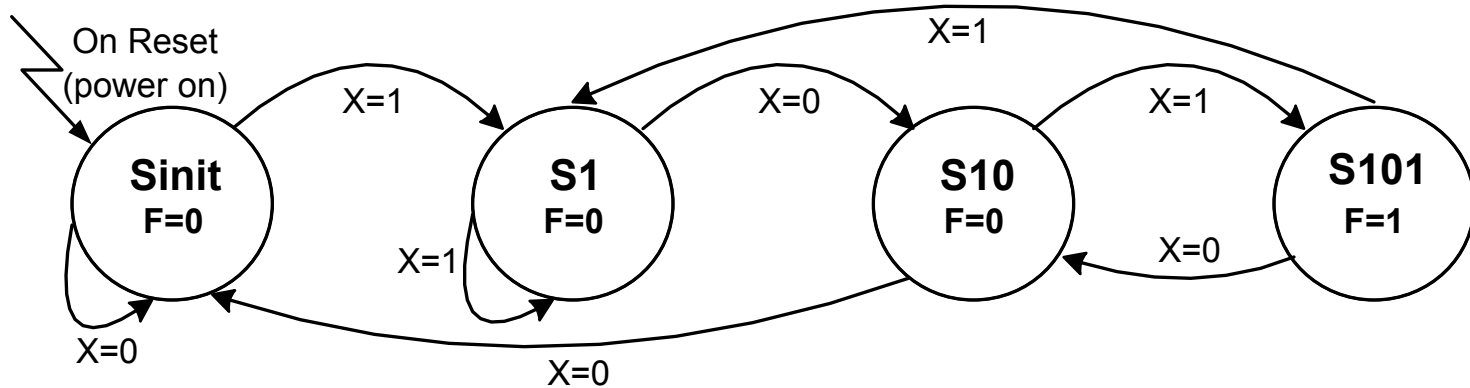
- Translate table to State Diagram

Current State			Next State						Output
			X = 0			X = 1			
$Q_1$	$Q_0$	State	$Q_1^*$	$Q_0^*$	State	$Q_1^*$	$Q_0^*$	State	F
0	0	Sinit	0	0	Sinit	0	1	S1	0
0	1	S1	1	0	S10	0	1	S1	0
1	0	S10	0	0	Sinit	1	1	S101	0
1	1	S101	1	0	S10	0	1	S1	1



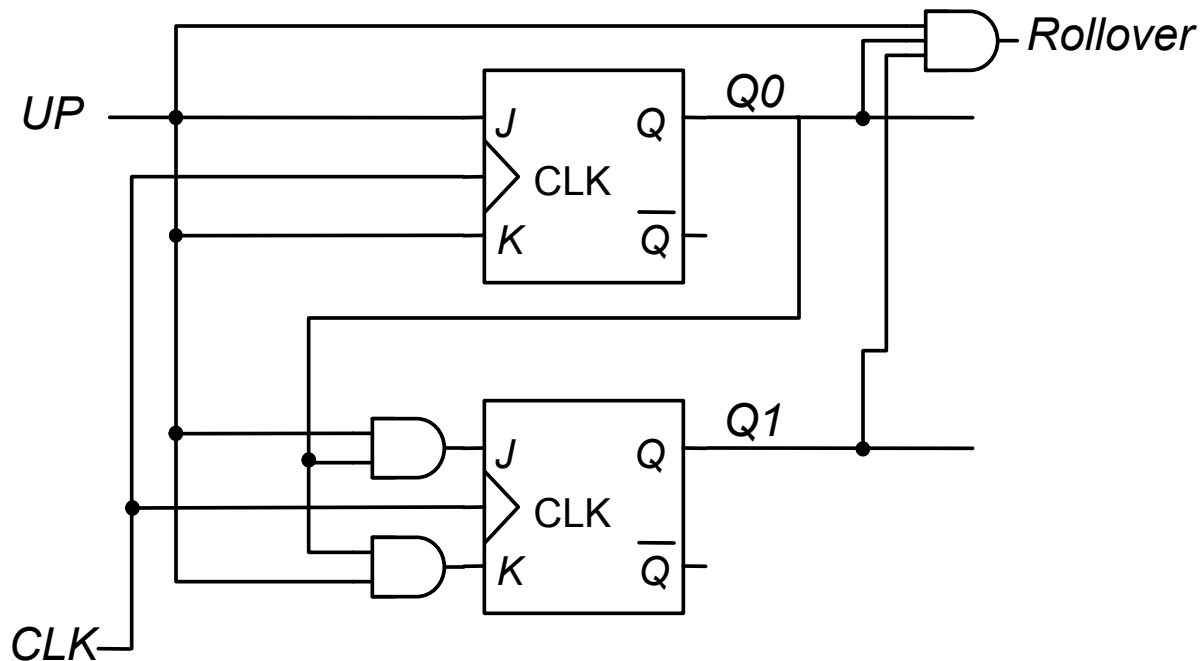
# State Diagram

- This state diagram implements the “101” sequence detector



# State Machine Analysis

- Consider the following circuit
- We now use JK FF's and a Mealy output



# State Machine Analysis

## Excitation Equations

- $J_0 = K_0 = Up$
- $J_1 = K_1 = Up \cdot Q_0$

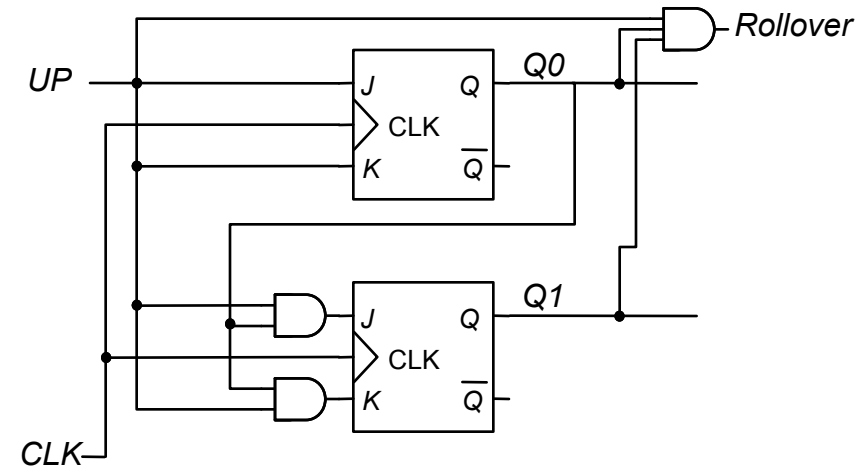
## Transition Equations

(char. eqn for JK FF:  $Q^* = JQ' + K'Q$ )

- $Q_0^* = Up \cdot Q_0' + Up' \cdot Q_0$
- $Q_1^* = Up \cdot Q_0 \cdot Q_1' + \overline{Up \cdot Q_0} \cdot Q_1$   
 $= Up \cdot Q_1' \cdot Q_0 + Up' \cdot Q_1 + Up \cdot Q_0 \cdot Q_1'$

## Output Equations

- $R = Up \cdot Q_1 \cdot Q_0$  (*Mealy output*)





# Transition Output Table

Current State			Next State							
			Up = 0				Up = 1			
Q1	Q0	State	Q1*	Q0*	State	R	Q1*	Q0*	State	R
0	0	S0	0	0	S0	0	0	1	S1	0
0	1	S1	0	1	S1	0	1	0	S2	0
1	0	S2	1	0	S2	0	1	1	S3	0
1	1	S3	1	1	S3	0	0	0	S0	1

Mealy outputs must be shown for each sub column of the inputs

$$Q0^* = Up \cdot Q0' + Up' \cdot Q0$$

$$Q1^* = Up \cdot Q1' \cdot Q0 + Up' \cdot Q1 + \cdot Q1Q0'$$

$$R = Up \cdot Q1 \cdot Q0$$

# Transition Output Table

Current State			Next State							
			Up = 0				Up = 1			
Q1	Q0	State	Q1*	Q0*	State	R	Q1*	Q0*	State	R
0	0	S0	0	0	S0	0	0	1	S1	0
0	1	S1	0	1	S1	0	1	0	S2	0
1	0	S2	1	0	S2	0	1	1	S3	0
1	1	S3	1	1	S3	0	0	0	S0	1

Notice that...

- When Up=1, we count up
- When Up=0 we stay in the current state
- Rollover = 1 only when we're about to rollover from S3 to S0