CSCI 104
Runtime Complexity

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REVIEW FROM CS 170
Steps for Deriving $T(n)$

- Considering an input of size $n$ that requires the maximum runtime, go through each line of the algorithm or code.
- Assume elementary operations such as incrementing a variable occur in constant time.
- If sequential blocks of code have runtime $T_1(n)$ and $T_2(n)$ respectively, then their total runtime will be their sum $T_1(n) + T_2(n)$.
- When we encounter loops, sum the runtime for each iteration of the loop, $T_i(n)$, to get the total runtime for the loop.
  - Nested loops often lead to summations of summations, etc.
Asymptotic Notation

• $T(n)$ is said to be $O(f(n))$ if...
  - $T(n) < a \cdot f(n)$ for $n > n_0$ (where $a$ and $n_0$ are constants)
  - Essentially an upper-bound
  - We'll focus on big-O for the worst case

• $T(n)$ is said to be $\Omega(f(n))$ if...
  - $T(n) > a \cdot f(n)$ for $n > n_0$ (where $a$ and $n_0$ are constants)
  - Essentially a lower-bound

• $T(n)$ is said to be $\Theta(f(n))$ if...
  - $T(n)$ is both $O(f(n))$ AND $\Omega(f(n))$
Helpful Common Summations

• \[ \sum_{i=1}^{n} i = \frac{n(n+1)}{2} = \theta(n^2) \]
  – This is called the arithmetic series

• \[ \sum_{i=1}^{n} \theta(i^p) = \theta(n^{p+1}) \]
  – This is a general form of the arithmetic series

• \[ \sum_{i=0}^{n} c^i = \frac{c^{n+1}-1}{c-1} = \theta(c^n) \]
  – This is called the geometric series

• \[ \sum_{i=1}^{n} \frac{1}{i} = \theta(\log n) \]
  – This is called the harmonic series
Runtime Practice #4

- It may seem like you can just look for nested loops and then raise n to that power
  - 2 nested for loops => $O(n^2)$
- But be careful!!
- Find $T(n)$ for this example

\[
\sum_{i=0}^{n-1} \sum_{j=0}^{\log_2(n)} \theta(1)
\]

= 

- Use the geometric sum eqn.

\[
\sum_{i=0}^{n-1} a^i = \frac{1-a^n}{1-a}
\]
- So our answer is...

```cpp
for (int i = 0; i <= log2(n); i ++)
    for (int j=0; j < (int) pow(2,i); j++)
        cout << j << endl;
```

Hint: Geometric series
Runtime Practice #3

- \( T(n) = \sum_{i=1}^{n} \left( \theta(1) + O\left(\sum_{j=0}^{n-1} \theta(1)\right) \right) \)
  - big-O indicates we have not considered the 'if' statement but are setting an upper bound

- \( T(n) = \sum_{i=1}^{n} \theta(1) + \sum_{i} \sum_{j=0}^{n-1} \theta(1) \) but we need to user our own analysis skills to find the actual values of \( i \) that will cause the 'if' to be true?
  - Use some actual values of \( n \) (e.g. \( n=9 \) or 16). Write out a table to find the pattern.
  - If \( n=9 \), the 'if' will trigger ___ times for \( i = \) ______________
  - If \( n=16 \), the 'if' will trigger ___ times for \( i = \) ______________
  - The dummy variable of a summation must increment ____ at a time
  - Thus, make a table with some dummy variable (k) that increments 1 at a time and find a relationship to the actual variable, \( i \), for when the if statement will trigger.
  - Solve for upper bound of \( k \)
    - Stop when \( i = \), but \( i = \) ____ so we stop when __________thus solve for \( k \) to find that the upper-bound for \( k = \) ________

- \( T(n) = 

```c++
for (int i = 1; i <= n; i++)
{
    int m = sqrt(n);
    if( i % m == 0){
        for (int j=0; j < n; j++)
            cout << j << " ";
    }
    cout << endl;
}
```
Key Skill

• The dummy variable (say k) of a summation runs from 1 to an **UPPER_BOUND** incrementing 1 at a time

• Often our code does work at some other interval such as \( i = \{1\sqrt{n}, 2\sqrt{n}, 3\sqrt{n} \ldots\} \) (or actual values that are not incrementing by 1 at a time)

• You must use your own analytical abilities to find a relationship that converts the dummy variable (\( k=1,2,3,\ldots \)) to the actual values [eg. \( i = f(k) = k\sqrt{n} \)], usually by making a table of the dummy variable (\( k \)) and the actual code values/variables (\( i \))

<table>
<thead>
<tr>
<th>k</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>...</th>
<th>Arbitrary k</th>
<th>Stop when k =??</th>
</tr>
</thead>
<tbody>
<tr>
<td>i</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>i = _______</td>
<td>Stop when i =______</td>
</tr>
</tbody>
</table>

• Then use that relationship to find the **UPPER_BOUND** of the dummy variable
  – In the previous example, we stopped when \( i = n \), thus we would stop when our dummy variable is \( \sqrt{n} \). This then is the upper bound.

• The key skill is to relate the dummy variable to the actual variable values and then find the **UPPER BOUND** of the dummy variable
Motivation

• You are given a large data set with $n = 500,000$ genetic markers for 5000 patients and you want to examine that data for genetic markers that maybe correlated to a disease that the patients have.
• You are given two algorithms, Algorithm A and Algorithm B, to solve this problem. You are given the implementation, code, and description of each algorithm.
• You need a solution as soon as possible to give medical professionals more data to advise patients and apply for grants for more funding.
• How would you determine which algorithm runs faster?
Runtime

• It is hard to compare the run time of an algorithm on actual hardware
  – Time may vary based on speed of the HW, etc.
    • The same program may take 1 sec. on your laptop but 0.5 second on a high performance server
• If we want to compare 2 algorithms that perform the same task we could try to count operations (regardless of how fast the operation can execute on given hardware)...
  – But what is an operation?
  – How many operations is: i++?
  – i++ actually requires grabbing the value of i from memory and bringing it to the processor, then adding 1, then putting it back in memory. Should that be 3 operations or 1?
    – Its painful to count 'exact' numbers operations
• Big-O, Big-Ω, and Θ notation allows us to be more general (or "sloppy" as you may prefer)
Complexity Analysis

• To find upper or lower bounds on the complexity, we must consider the set of all possible inputs, I, of size, n
• Derive an expression, \( T(n) \), in terms of the input size, \( n \), for the number of operations/steps that are required to solve the problem of a given input, \( i \)
  – Some algorithms depend on \( i \) and \( n \)
    • Find(3) in the list shown vs. Find(2)
  – Others just depend on \( n \)
    • Push_back / Append

• Which inputs though?
  – Best, worst, or "typical/average" case?
• We will always apply it to the "worst case"
  – That's usually what people care about

Note: Running time of an algorithm is not just based on input size (n), BUT input size (n) and its value (i)
Time Complexity Analysis

- Case Analysis is when you determine which input must be used to define the runtime function, $T(n)$, for inputs of size $n$.

- **Best-case analysis**: Find the input of size $n$ that takes the minimum amount of time.

- **Average-case analysis**: Find the runtime for all inputs of size $n$ and take the average of all of the runtimes. (This assumes a distribution over the inputs, but uniform is a reasonable choice.)

- **Worst-case analysis**: Find the input, $i$, of size $n$ that takes the maximum amount of time.

- Our focus will be on worst-case analysis, but for many examples, the runtime is the same on any input of size $n$. Please consider this as we study them.
Steps for Performing Runtime Analysis of Algorithms

• We perform **worst-case analysis** in determining the runtime function on inputs of size $n$, $T(n)$.

• To do so, we need to find at least one input of size $n$ that will require the **maximum** runtime of the algorithm.
  – In many of the examples we will examine, the algorithm will take the same amount of running time on any input (i.e. only depend on $n$)

• Using that input, **express the runtime of the algorithm** (on that input case) as a function of $n$, $T(n)$.
  – This is done by **stepping through the code and counting the steps** that will be done.

• Once we have a function for the runtime, $T(n)$, we **apply asymptotic notation to that function** in order to find the order of growth of the runtime function, $T(n)$.
Asymptotic Notation

- **T(n) is said to be $O(f(n))$ if...**
  - $T(n) < a \cdot f(n)$ for $n > n_0$ (where $a$ and $n_0$ are constants)
  - Essentially an upper-bound
  - We'll focus on big-O for the worst case

- **T(n) is said to be $\Omega(f(n))$ if...**
  - $T(n) > a \cdot f(n)$ for $n > n_0$ (where $a$ and $n_0$ are constants)
  - Essentially a lower-bound

- **T(n) is said to be $\Theta(f(n))$ if...**
  - $T(n)$ is both $O(f(n))$ AND $\Omega(f(n))$
Worst Case and Big-Ω

• What's the lower bound on List::find(val)
  – Is it \( \Omega(1) \) since we might find the given value on the first element?
  – Well it could be if we are finding a lower bound on the 'best case'

• Big-Ω does **NOT** have to be **synonymous** with 'best case'
  – Though many times it mistakenly is

• You can have:
  – Big-O for the best, average, worst cases
  – Big-Ω for the best, average, worst cases
  – Big-Θ for the best, average, worst cases

• Note:
  – Big-O and Big-Ω analysis are **ONLY** necessary when the runtime of the algorithm is **data-dependent** (i.e. function of inputs / \( T(n,i) \)).
  – If the code is **NOT data-dependent** then your analysis is valid for any input and thus is already a tight bound (big- Θ)
Worst Case and Big-$$\Omega$$

- The key idea is an algorithm may perform differently for different input cases
  - Imagine an algorithm that processes an array of size n but depends on what data is in the array
- Big-$$O$$ for the worst-case says for REGARDLESS of possible inputs the runtime is bound (at-most) by $$O(f(n))$$
- Big-$$\Omega$$ for the worst-case is attempting to establish a lower bound (at-least) for the worst case (the worst case is just one of the possible input scenarios)
  - If we look at the first data combination in the array and it takes $$n$$ steps then we can say the algorithm is $$\Omega(n)$$.
  - Now we look at the next data combination in the array and the algorithm takes $$n^{1.5}$$. We can now say worst case is $$\Omega(n^{1.5})$$.
- To arrive at $$\Omega(f(n))$$ for the worst-case requires you simply to find AN input case (i.e. the worst case) that requires at least $$f(n)$$ steps
- Cost analogy...

```c
int i; j;
for(i=0; i < n; i++){
    if(a[i][0] == 0){
        for(j=0; j<n; j++)
            a[i][j] = i*j;
    }
}
```

Consider the effect of the 'if' statement. Can it be true for each value of $$i$$? If we don't want to (or can't) determine this we can assume it will be true and say that the upper bound for the runtime is $$O(n^2)$$. To prove it is $$\Theta(n^2)$$ we'd need to prove there is a set of inputs for the a matrix that makes the 'if' true on each iteration (i.e. $$\Omega(n^2))$$. 
Steps for Deriving $T(n)$

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- If sequential blocks of code have runtime $T_1(n)$ and $T_2(n)$ respectively, then their total runtime will be their sum $T_1(n)+T_2(n)$.
- When we encounter loops, sum the runtime for each iteration of the loop, $T_i(n)$, to get the total runtime for the loop.
  - Nested loops often lead to summations of summations, etc.
Helpful Common Summations

- $\sum_{i=1}^{n} i = \frac{n(n+1)}{2} = \theta(n^2)$
  - This is called the arithmetic series

- $\sum_{i=1}^{n} \theta(i^p) = \theta(n^{p+1})$
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- $\sum_{i=0}^{n} c^i = \frac{c^{n+1}-1}{c-1} = \theta(c^n)$
  - This is called the geometric series

- $\sum_{i=1}^{n} \frac{1}{i} = \theta(\log n)$
  - This is called the harmonic series
Deriving T(n)

- Derive an expression, T(n), in terms of the input size for the number of operations/steps that are required to solve a problem
- If is true => 4 "steps"
- Else if is true => 5 "steps"
- Worst case => T(n) = $\theta(1)$

```cpp
#include <iostream>

using namespace std;

int main(int argc, char* argv[])
{
    int i = argc;  // 1
    int x = 5;    // 1

    if(i < x){    // 1
        x--;    // 1
    }
    else if(i > x){  // 1
        x += 2;    // 1
    }
    return 0;
}
```
Deriving $T(n)$

- Since loops repeat you have to take the sum of the steps that get executed over all iterations

- $T(n) = \sum_{i=0}^{n-1} 4 = 4 + 4 + \cdots + 4 = 4 \cdot n$

- $= \theta(n)$

```cpp
#include <iostream>
using namespace std;

int main()
{
    int x;
    for(int i=0; i < N; i++){
        cin >> x;
        if(i < x){
            x--;
        }
        else if(i > x){
            x += 2;
        }
    }
    return 0;
}
```

This code does nothing useful and is just illustrative.
Skills To Gain

- To solve these runtime problems try to break the problem into 3 parts:
  - FIRST, setup the expression (or recurrence relationship) for the number of operations, $T(n)$
  - SECOND, solve to get a closed form for $T(n)$
    - Unwind the recurrence relationship
    - Develop a series summation
    - Solve the series summation
  - THIRD, determine the asymptotic bound for $T(n)$
Loops 1

- Derive an expression, \( T(n) \), in terms of the input size for the number of operations/steps that are required to solve a problem

\[
T(n) = \sum_{i=0}^{n-1} \sum_{j=0}^{n-1} \theta(1) = \sum_{i=0}^{n-1} \theta(n) = \Theta(n^2)
\]
Matrix Multiply

- Derive an expression, $T(n)$, in terms of the input size for the number of operations/steps that are required to solve a problem
- $T(n) = \sum_{i=0}^{n-1} \sum_{j=0}^{n-1} \sum_{k=0}^{n-1} \theta(1) = \theta(n^3)$

```cpp
#include <iostream>
using namespace std;

const int n = 256;
int a[n][n], b[n][n], c[n][n];
int main()
{
    for(int i=0; i < n; i++)
        for(int j=0; j < n; j++)
            c[i][j] = 0;
    for(int k=0; k < n; k++)
        c[i][j] += a[i][k]*b[k][j];
    return 0;
}
```
Sequential Loops

- Is this also $n^3$?
- _____________
  - 3 for loops, ______________

```cpp
#include <iostream>
using namespace std;

const int n = /* large constant */;

unsigned char image[n][n]

int main()
{
    for(int i=0; i < n; i++){
        image[0][i] = 5;
    }
    for(int j=0; j < n; j++){
        image[1][j] = 5;
    }
    for(int k=0; k < n; k++){
        image[2][k] = 5;
    }
    return 0;
}
```
Runtime Practice #1

• Count steps here...
  – Think about how many times if statement will evaluate true

  ```java
  for(int i=0; i < n; i++){
      if (a[i][0] == 0){
          for (int j = 0; j < i; j++){
              a[i][j] = i*j;
          }
      }
  }
  ```

  **Hint: Arithmetic series**

• \( T(n) = \) ____________________________ May start with big-O and not worry about input values affecting how many times if statement executes

• \( T(n) = \sum_{i=0}^{n-1} (\theta(1)) + \sum_{i} (\theta(i)) \) Distribute to deal with 'if' separately. Not sure which values of i will trigger the for loop that incurs i steps
  – In the worst case, how many times can the 'if' statement be true? _________________

• \( T(n) = \)
Runtime Practice #2

- $T(n) =$

- $T(n) = \sum_{i=0}^{n-1} \left( \theta(1) + O(\sum_{j=1}^{n} \theta(1)) \right)$ Use big-O since unsure of how many times if statement executes
  - Important: How many times will the 'if' statement be true?

- $T(n) = \sum_{i=0}^{n-1} \theta(1) + \sum_{i} \sum_{j=1}^{n} \theta(1)$
  - The 'if' statement only triggers once! So the inner loop executes only once

- $T(n) =$

```java
for(int i=0; i < n; i++){
    if (i == 0){
        for (int j = 0; j < n; j++){
            a[i][j] = i*j;
        }
    }
}
```
Runtime Practice #3

- $T(n) = \sum_{i=1}^{n} \left( \theta(1) + O(\sum_{j=0}^{n-1} \theta(1)) \right)$
  - big-O indicates we have not considered the 'if' statement but are setting an upper bound

- $T(n) = \sum_{i=1}^{n} \theta(1) + \sum_i \sum_{j=0}^{n-1} \theta(1)$ but we need to use our own analysis skills to find the actual values of $i$ that will cause the 'if' to be true?
  - Use some actual values of $n$ (e.g. $n=9$ or $16$). Write out a table to find the pattern.
  - If $n=9$, the 'if' will trigger ___ times for $i = \______________$
  - If $n=16$, the 'if' will trigger ___ times for $i = \______________$
  - The dummy variable of a summation must increment ____ at a time
  - Thus, make a table with some dummy variable ($k$) that increments 1 at a time and find a relationship to the actual variable, $i$, for when the if statement will trigger.
  - Solve for upper bound of $k$
    - Stop when $i = __$, but $i = ___$ so we stop when __________thus solve for $k$ to find that the upper-bound for $k = ______$

- $T(n) =$

```cpp
for (int i = 1; i <= n; i++)
{
    int m = sqrt(n);
    if( i % m == 0){
        for (int j=0; j < n; j++)
            cout << j << " ";
    }
    cout << endl;
}
```

| $k$ | 1 | 2 | 3 | ... | Arbitrary k | Stop when $k =??$
|-----|---|---|---|-----|------------|------------------|
| $i$ |   |   |   |     | ... $i = ______$ | Stop when $i = ______$
Key Skill

• The dummy variable (say k) of a summation runs from 1 to an \textbf{UPPER\_BOUND} incrementing 1 at a time

• Often our code does work at some other interval such as \(i = \{1\sqrt{n}, 2\sqrt{n}, 3\sqrt{n} \ldots\} \) (or actual values that are not incrementing by 1 at a time)

• You must use your own analytical abilities to find a relationship that converts the dummy variable (k=1,2,3,...) to the actual values [eg. \(i = f(k) = k\sqrt{n}\)], usually by making a table of the dummy variable (k) and the actual code values/variables (i)

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<tbody>
<tr>
<td>i</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>i = _______</td>
<td>Stop when i = _______</td>
</tr>
</tbody>
</table>

• Then use that relationship to find the \textbf{UPPER\_BOUND} of the dummy variable
  
  – In the previous example, we stopped when \(i = n\), thus we would stop when our dummy variable is \(\sqrt{n}\). This then is the upper bound.

• The key skill is to relate the dummy variable to the actual variable values and then find the \textbf{UPPER BOUND} of the dummy variable
It may seem like you can just look for nested loops and then raise n to that power

- 2 nested for loops => $O(n^2)$

But be careful!!

Find $T(n)$ for this example

$\sigma_i = \sum_{j=0}^{n-1} \theta(1)$

$= \sum_{i=0}^{n-1} a^i = \frac{1-a^n}{1-a}$

So our answer is...

```cpp
for (int i = 0; i <= log2(n); i++)
    for (int j=0; j < (int) pow(2,i); j++)
        cout << j << endl;
```

Hint: Geometric series
Runtime Practice #5

• \( T(n) = \)

\[
T(n) = \sum_{i=1}^{n} (\theta(1) + \sum_{j} \theta(1)) = \theta(n) + \sum_{i=1}^{n} \sum_{j} \theta(1)
\]

• Manually, determine how many times the j-loop iterates:
  – When \( i=1 \), j takes on values: __________________________ [Total = _____ iters]
  – When \( i=2 \), j takes on values: __________________________ [Total = _____ iters]
  – When \( i=3 \), j takes on values: __________________________ [Total = _____ iters]

• \( T(n) = \theta(n) + \)

```java
for(int i=1; i <= n; i++){
    for (int j = 0; j < n; j += i){
        a[i][j] = i*j;
    }
}
```

Hint: Harmonic series
Runtime Practice #6

• You have to count steps
  – Look at the update statement
  – Outer loop increments by 1 each time so it will iterate N times
  – Inner loop updates by dividing x in half each iteration?
  – After 1\textsuperscript{st} iteration => x=____
  – After 2\textsuperscript{nd} iteration => x=____
  – After 3\textsuperscript{rd} iteration => x=____
  – Say k\textsuperscript{th} iteration is last => x = ______ = 1
  – Solve for k
  – k = __________ iterations
  – \( \theta(__________) \)

#include <iostream>
using namespace std;
const int n = /* Some constant */;

int main()
{
    for(int i=0; i < n; i++){
        int y=0;
        for(int x=n; x != 1; x=x/2){
            y++;
        }
        cout << y << endl;
    }
    return 0;
}
Importance of Complexity

<table>
<thead>
<tr>
<th>N</th>
<th>O(1)</th>
<th>O(log₂n)</th>
<th>O(n)</th>
<th>O(n*log₂n)</th>
<th>O(n²)</th>
<th>O(2^n)</th>
</tr>
</thead>
<tbody>
<tr>
<td>2</td>
<td>1</td>
<td>1</td>
<td>2</td>
<td>2</td>
<td>4</td>
<td>4</td>
</tr>
<tr>
<td>20</td>
<td>1</td>
<td>4.3</td>
<td>20</td>
<td>86.4</td>
<td>400</td>
<td>1,048,576</td>
</tr>
<tr>
<td>200</td>
<td>1</td>
<td>7.6</td>
<td>200</td>
<td>1,528.8</td>
<td>40,000</td>
<td>1.60694E+60</td>
</tr>
<tr>
<td>2000</td>
<td>1</td>
<td>11.0</td>
<td>2000</td>
<td>21,931.6</td>
<td>4,000,000</td>
<td>#NUM!</td>
</tr>
</tbody>
</table>
EXTRAS
Runtime Practice #7

- \( T(n) = \sum_{i=1}^{n} \left( \theta(1) + O\left(\sum_{j=1}^{i} \theta(1)\right) \right) \)

- Important: How many times will the 'if' statement be true?
- \( T(n) = \sum_{i=1}^{n} \left( \theta(1) \right) + \sum_{i} \sum_{j=1}^{n} \theta(1) \)
  - Find a relationship between a dummy variable, \( k \), that increments by 1 and the values of \( i \) that cause the if statement to trigger

\[
\begin{array}{cccccc}
\begin{array}{c|c|c|c|c|c|c|}
 k & 1 & 2 & 3 & \ldots & \text{Arbitrary } k & \text{Stop when } k = (n/2)+1 \\
 i & 0 & 2 & 4 & \ldots & i = \_\_\_\_\_ & \text{Stop when } i = \_\_\_\_\_ \\
\end{array}
\end{array}
\]

- \( T(n) = \)

Recall: \( \sum_{i=1}^{n} i = \frac{n(n+1)}{2} = \theta(n^2) \)
Runtime Practice #8

• \( T(n) = \)

• \( T(n) = \sum_i \left( \sum_{j=0}^{i-1} \theta(1) \right) = \sum_i \left( \theta(i) \right) \)

• The number of iterations of the outer loop requires derivation:

\[
\begin{array}{|c|c|c|c|c|c|c|}
\hline
\text{Iter}, k & 1 & 2 & 3 & 4 & \ldots & k & \text{Stop at: (log_2 n)} \\
\hline
\text{i after iteration} & 2 & 4 & 8 & 16 & \ldots & 2^k & \text{Stop at: n} \\
\hline
\end{array}
\]

• \( T(n) = \sum_{k=1}^{\log_2(n)} \theta(2^k) \)

• \( T(n) = \theta \left( \frac{2^{\log_2(n)+1}-1}{2-1} \right) = \theta \left( \frac{2^{\log_2(n)2^1}-1}{1} \right) = \theta(2n - 1) = \theta(n) \)

for(int i=1; i <= n; i*=2){
    for (int j = 0; j < i; j++){
        a[i][j] = i*j;
    }
}

Hint: Geometric series
Iterative Binary Search

- Assume n is total array size and let $L = (\text{end-start})$
  - $L = \#$ of items to be searched
- $T(n) = \sum_k \theta(1)$
  - $k$ is the $\#$ of iterations required
- After 1$^{\text{st}}$ iteration $L = n/2$
- After 2$^{\text{nd}}$ iteration $L = n/4$
- After 3$^{\text{rd}}$ iteration $L = n/8$
- ...
- After $k$th iteration $L = n/2^k$
- We stop when we reach size 0 or 1...when $k = \log_2(n)$
- $T(n) = \sum_{k=1}^{\log_2(n)} \theta(1) = \theta(\log_2(n))$

```c
int main()
{
    int data[4] = {1, 6, 7, 9};
    it_bsearch(3, data, 4);
}

int it_bsearch(int target, int data[], int len)
{
    int start = 0, end = len, mid;

    while (start < end) {
        mid = (start+end)/2;
        if (data[mid] == target){
            return mid;
        } else if (target < data[mid]){
            end = mid-1;
        } else {
            start = mid+1;
        }
    }
    return -1;
}
```
SOLUTIONS
Sequential Loops

• Is this also $n^3$?
• No!
  – 3 for loops, but not nested
  – $O(n) + O(n) + O(n) = 3*O(n) = O(n)$

```cpp
#include <iostream>
using namespace std;

const int n = /* large constant */;

unsigned char image[n][n]

int main()
{
    for(int i=0; i < n; i++){
        image[0][i] = 5;
    }
    for(int j=0; j < n; j++){
        image[1][j] = 5;
    }
    for(int k=0; k < n; k++){
        image[2][k] = 5;
    }
    return 0;
}
```
Runtime Practice #1

• Count steps here...
  – Think about how many times if statement will evaluate true

\[ T(n) = \sum_{i=0}^{n-1} \left( \theta(1) + O(\sum_{j=1}^{i} \theta(1)) \right) \]

May start with big-O and not worry about input values affecting how many times if statement executes

• \[ T(n) = \sum_{i=0}^{n-1} (\theta(1)) + \sum_{i=0}^{n-1} (\theta(i)) \]
  Distribute to deal with 'if' separately. Not sure which values of i will trigger the for loop that incurs i steps
  – In the worst case, how many times can the 'if' statement be true? Each iteration (i.e. all n values of i)

• \[ T(n) = \sum_{i=0}^{n-1} (\theta(1)) + \sum_{i=0}^{n-1} (\theta(i)) \]

• \[ T(n) = \theta(n) + \sum_{i=0}^{n-1} (\theta(i)) = \theta(n) + \theta \left( \frac{n(n-1)}{2} \right) = \theta(n^2) \]

for(int i=0; i < n; i++){
    if (a[i][0] == 0){
        for (int j = 0; j < i; j++){
            a[i][j] = i*j;
        }
    }
}

Hint: Arithmetic series
Runtime Practice #2

- $T(n) = \sum_{i=0}^{n-1} \left( \theta(1) + O\left(\sum_{j=0}^{n-1} \theta(1)\right) \right)$ Use big-O since unsure of how many times if statement executes
  - Important: How many times will the 'if' statement be true?
- $T(n) = \sum_{i=0}^{n-1} \theta(1) + \sum_i \sum_{j=0}^{n-1} \theta(1)$
  - The 'if' statement only triggers once! So the inner loop executes only once
- $T(n) = \theta(n) + 1 \cdot \sum_{j=0}^{n-1} \theta(1) = \theta(n) + \theta(n) = \theta(n)$
Runtime Practice #3

- \( T(n) = \sum_{i=1}^{n} \left( \theta(1) + O \left( \sum_{j=0}^{n-1} \theta(1) \right) \right) \)
  - big-O indicates we have not considered the 'if' statement but are setting an upper bound

- \( T(n) = \sum_{i=1}^{n} \theta(1) + \sum_{i} \sum_{j=0}^{n-1} \theta(1) \) but we need to use our own analysis skills to find the actual values of \( i \) that will cause the 'if' to be true?
  - Use some actual values of \( n \) (e.g. \( n=9 \) or \( 16 \)). Write out a table to find the pattern.
  - If \( n=9 \), the 'if' will trigger 3 times for \( i = 3, 6, 9 \)
  - If \( n=16 \), the 'if' will trigger 4 times for \( i = 4, 8, 12, 16 \)
  - The dummy variable of a summation must increment 1 at a time
  - Thus, make a table with some dummy variable (\( k \)) that increments 1 at a time and find a relationship to the actual variable, \( i \), for when the if statement will trigger.
  - Solve for upper bound of \( k \)
    - Stop when \( i = n \), but \( i = k\sqrt{n} \) so we stop when \( k\sqrt{n} = n \) thus solve for \( k \) to find that the upper-bound for \( k = \sqrt{n} \)

- \( T(n) = \theta(n) + \sum_{k=1}^{\sqrt{n}} \sum_{j=0}^{n-1} \theta(1) = \theta(n) + \sum_{k=1}^{\sqrt{n}} \theta(n) = \theta(n) + \theta(n \cdot \sqrt{n}) = \theta(n^{3/2}) \)

```cpp
for (int i = 1; i <= n; i++)
{
  int m = sqrt(n);
  if (i % m == 0){
    for (int j=0; j < n; j++)
      cout << j << " ";
  }
  cout << endl;
}
```
Runtime Practice #4

• It may seem like you can just look for nested loops and then raise n to that power
  — 2 nested for loops => $O(n^2)$
• But be careful!!
• Find $T(n)$ for this example

$$\sum_{i=0}^{\log(n)} \sum_{j=0}^{2^i-1} \theta(1)$$

$$= \sum_{i=0}^{\log(n)} \theta(2^i)$$

• Use the geometric sum eqn.
  $$= \sum_{i=0}^{n-1} a^i = \frac{1-a^n}{1-a}$$
• So our answer is...
  $$= \frac{1-2^{\log(n)+1}}{1-2} = \frac{1-2*2^n}{-1} = \theta(n)$$

```
for (int i = 0; i <= log2(n); i++)
  for (int j=0; j < (int) pow(2,i); j++)
    cout << j << endl;
```

*Hint: Geometric series*
Runtime Practice #5

- \( T(n) = \)

- \( T(n) = \sum_{i=1}^{n}(\theta(1) + \sum_{j} \theta(1)) = \theta(n) + \sum_{i=1}^{n} \sum_{j} \theta(1) \)

- Manually, determine how many times the j-loop iterates:
  - When \( i=1 \), j takes on values: 0, 1, 2, 3, … , n-1 [Total = n iters]
  - When \( i=2 \), j takes on values: 0, 2, 4, 6, … , n-2 or n-1 [Total = n/2 iters]
  - When \( i=3 \), j takes on values: 0, 3, 6, 9, … [Total = n/3 iters]

- \( T(n) = \theta(n) + \left[ \frac{n}{1} + \frac{n}{2} + \frac{n}{3} + \cdots + \frac{n}{n} \right] \theta(1) \)
  \[ = \theta(n) + \left[ \frac{1}{1} + \frac{1}{2} + \frac{1}{3} + \cdots + \frac{1}{n} \right] \theta(n) \]
  \[ = \theta(n) + \left( \sum_{i=1}^{n} \frac{1}{i} \right) \cdot \theta(n) = \theta(n) + \log n \cdot \theta(n) = \theta(n \cdot \log n) \]

for(int i=1; i <= n; i++){
    for (int j = 0; j < n; j += i){
        a[i][j] = i*j;
    }
}  
**Hint: Harmonic series**
Runtime Practice #6

- You have to count steps
  - Look at the update statement
  - Outer loop increments by 1 each time so it will iterate N times
  - Inner loop updates by dividing x in half each iteration?
  - After 1\textsuperscript{st} iteration => x=n/2
  - After 2\textsuperscript{nd} iteration => x=n/4
  - After 3\textsuperscript{rd} iteration => x=n/8
  - Say k\textsuperscript{th} iteration is last => x = n/2^k = 1
  - Solve for k
  - k = \log_2(n) iterations
  - \( \theta(n \log(n)) \)

```cpp
#include <iostream>
using namespace std;
const int n = /* Some constant */;
int main()
{
    for(int i=0; i < n; i++){
        int y=0;
        for(int x=n; x != 1; x=x/2){
            y++;
        }
        cout << y << endl;
    }
    return 0;
}
```
Runtime Practice #7

- \( T(n) = \sum_{i=1}^{n} \left( \theta(1) + O\left( \sum_{j=1}^{i} \theta(1) \right) \right) \)

- Important: How many times will the 'if' statement be true?

- \( T(n) = \sum_{i=1}^{n} \left( \theta(1) \right) + \sum_{i} \sum_{j=1}^{n} \theta(1) \)

  - Find a relationship between a dummy variable, \( k \), that increments by 1 and the values of \( i \) that cause the if statement to trigger

<table>
<thead>
<tr>
<th>( k )</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>...</th>
<th>Arbitrary ( k )</th>
<th>Stop when ( k = (n/2)+1 )</th>
</tr>
</thead>
<tbody>
<tr>
<td>( i )</td>
<td>0</td>
<td>2</td>
<td>4</td>
<td>...</td>
<td>( i = 2(k-1) )</td>
<td>Stop when ( i = n )</td>
</tr>
</tbody>
</table>

- \( T(n) = \theta(n) + \sum_{k=1}^{n+1} \sum_{j=1}^{2(k-1)} \theta(1) = \theta(n) + \sum_{k=1}^{n+1} \theta(2k - 1) = \theta(n)+2 \cdot \sum_{k=1}^{n+1} \theta(k) = \theta(n) + 2 \cdot \theta \left( \left[ \frac{n}{2} + 1 \right]^2 \right) = \theta(n^2) \)

Recall: \( \sum_{i=1}^{n} i = \frac{n(n+1)}{2} = \theta(n^2) \)
Runtime Practice #8

- \( T(n) = \)

- \( T(n) = \sum_i (\sum_{j=0}^{i-1} \theta(1)) = \sum_i (\theta(i)) \)

- The number of iterations of the outer loop requires derivation:

<table>
<thead>
<tr>
<th>Iter, k</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>...</th>
<th>k</th>
<th>(log₂n)</th>
</tr>
</thead>
<tbody>
<tr>
<td>i after iteration</td>
<td>2</td>
<td>4</td>
<td>8</td>
<td>16</td>
<td>...</td>
<td>( 2^k )</td>
<td>n</td>
</tr>
</tbody>
</table>

- \( T(n) = \sum_{k=1}^{\log_2(n)} \theta(2^k) \)

- \( T(n) = \theta \left( \frac{2^{\log_2(n)+1} - 1}{2-1} \right) = \theta \left( \frac{2^{\log_2(n)}2^1 - 1}{1} \right) = \theta(2n - 1) = \theta(n) \)

for(int i=1; i <= n; i*=2){
    for (int j = 0; j < i; j++){
        a[i][j] = i*j;
    }
}