

CS103 Unit 8

Recursion

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Recursion

- Defining an object, mathematical function, or computer function in terms of *itself*



GNU

- Makers of gedit, g++ compiler, etc.
- GNU = GNU is Not Unix

GNU is Not Unix

GNU is Not Unix

... is Not Unix is not Unix is Not Unix

Recursion

- Problem in which the solution can be expressed in terms of itself (usually a smaller instance/input of the same problem) **and a base/terminating case**
- Usually takes the place of a loop
- Input to the problem must be categorized as a:
 - Base case: Solution known beforehand or easily computable (no recursion needed)
 - Recursive case: Solution can be described using solutions to smaller problems of the same type
 - Keeping putting in terms of something smaller until we reach the base case
- Factorial: $n! = n * (n-1) * (n-2) * \dots * 2 * 1$
 - $n! = n * (n-1)!$
 - Base case: $n = 1$
 - Recursive case: $n > 1 \Rightarrow n * (n-1)!$

Recursive Functions

- Recall the system stack essentially provides separate areas of memory for each 'instance' of a function
- Thus each **local variable** and **actual parameter** of a function has its own value within that particular function instance's memory space

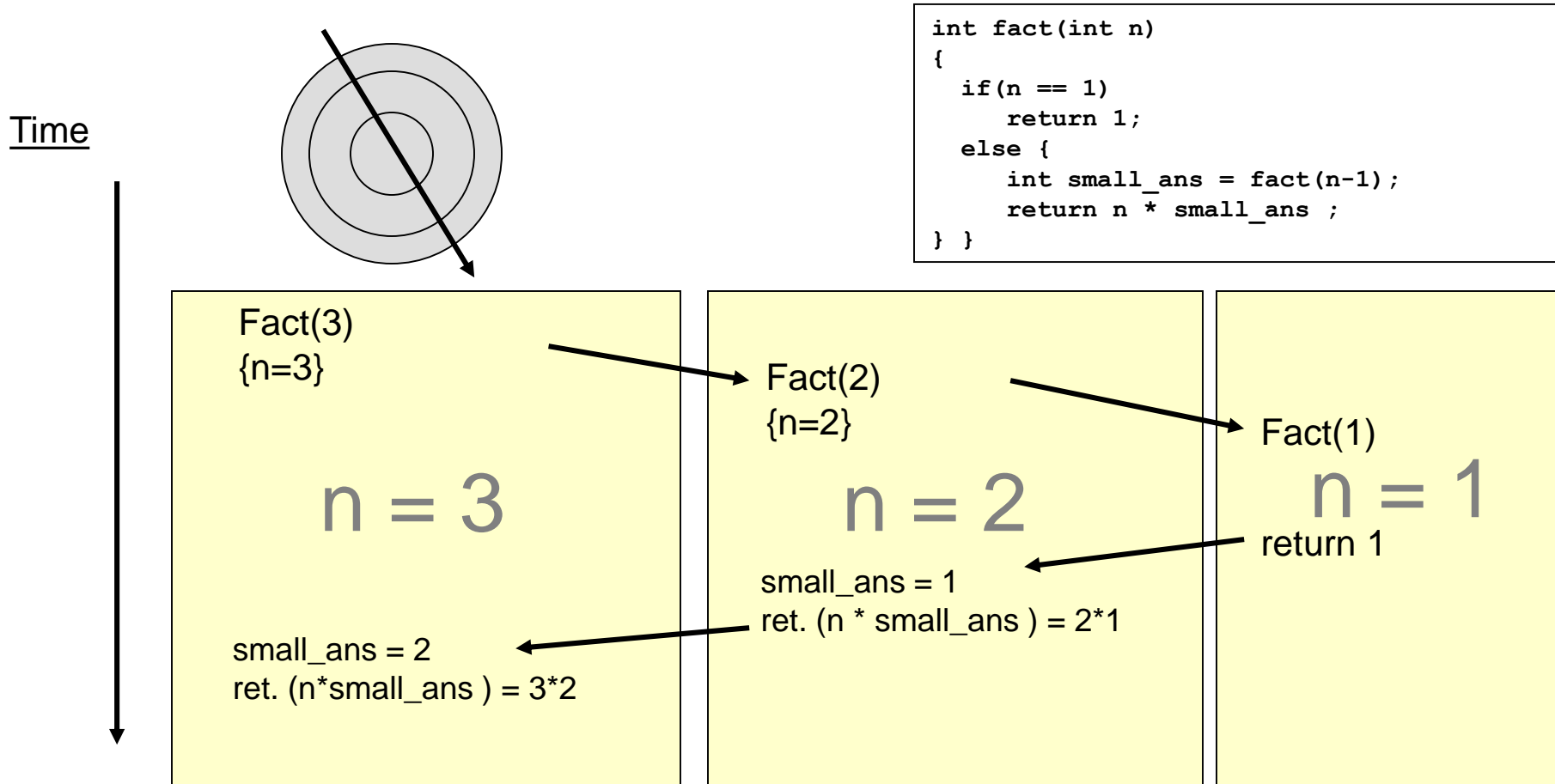
C Code:

```
int fact(int n)
{
    // base case
    if(n == 1)
        return 1;

    // recursive case
    else {
        // calculate (n-1)!
        int small_ans = fact(n-1);

        // now ans = (n-1)!
        // so calculate n!
        return n * small_ans;
    }
}
```

Recursive Call Timeline



- Value/version of `n` is implicitly “saved” and “restored” as we move from one instance of the ‘fact’ function to the next

Head vs. Tail Recursion

- Head Recursion: Recursive call is made before the real work is performed in the function body
- Tail Recursion: Some work is performed and then the recursive call is made

Tail Recursion

```
void doit(int n)
{
    if(n == 1) cout << "Stop";
    else {
        cout << "Go" << endl;
        doit(n-1);
    }
}
```

Head Recursion

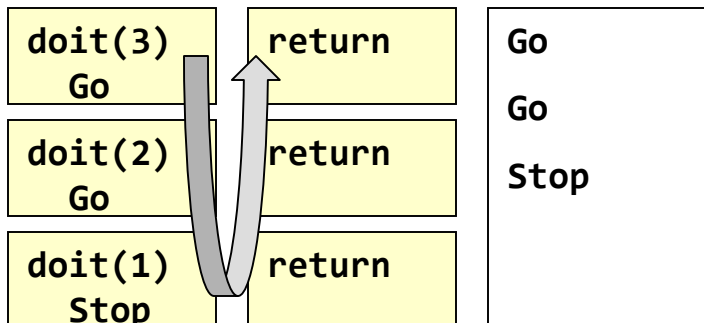
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        cout << "Go" << endl;
    }
}
```

Head vs. Tail Recursion

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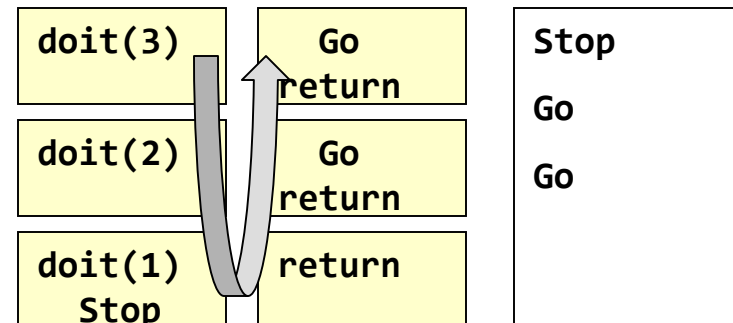
Tail Recursion

```
Void doit(int n)
{
    if(n == 1) cout << "Stop";
    else {
        cout << "Go" << endl;
        doit(n-1);
    }
}
```



Head Recursion

```
Void doit(int n)
{
    if(n == 1) cout << "Stop";
    else {
        doit(n-1);
        cout << "Go" << endl;
    }
}
```



Recursive Functions

- Recall the system stack essentially provides separate areas of memory for each 'instance' of a function
- Thus each **local variable** and **actual parameter** of a function has its own value within that particular function instance's memory space

C Code:

```
int main()
{
    int data[4] = {8, 6, 7, 9};
    int sum1 = isum_it(data, 4);
    int sum2 = rsum_it(data, 4);
}

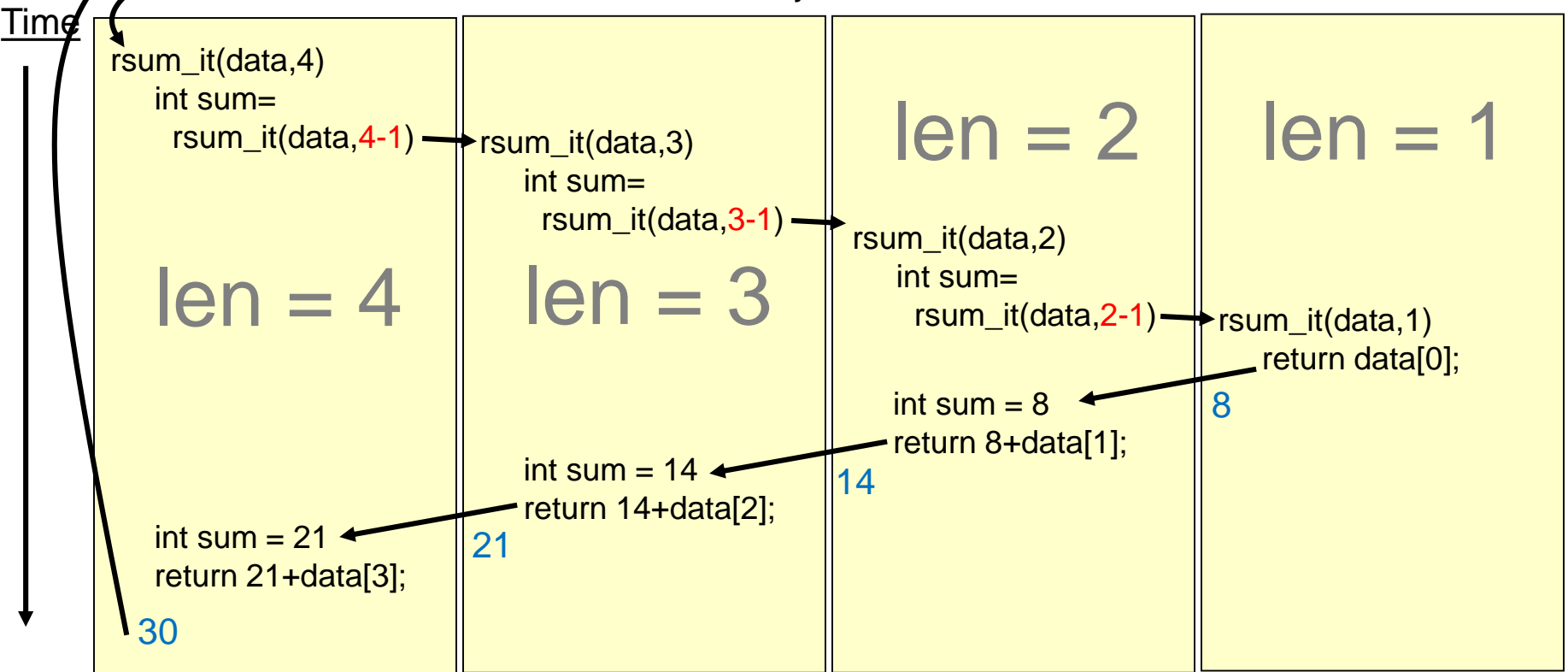
int isum_it(int data[], int len)
{
    sum = data[0];
    for(int i=1; i < len; i++){
        sum += data[i];
    }
}

int rsum_it(int data[], int len)
{
    if(len == 1)
        return data[0];
    else
        int sum = rsum_it(data, len-1);
        return sum + data[len-1];
}
```


Recursive Call Timeline

```
int main(){
    int data[4] = {8, 6, 7, 9};
    int sum2 = rsum_it(data, 4);
    ...
}
```

```
int rsum_it(int data[], int len)
{
    if(len == 1)
        return data[0];
    else
        int sum = rsum_it(data, len-1);
        return sum + data[len-1];
}
```



Each instance of `rsum_it` has its own `len` argument and `sum` variable

Every instance of a function has its own copy of local variables

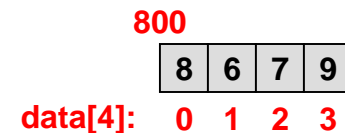
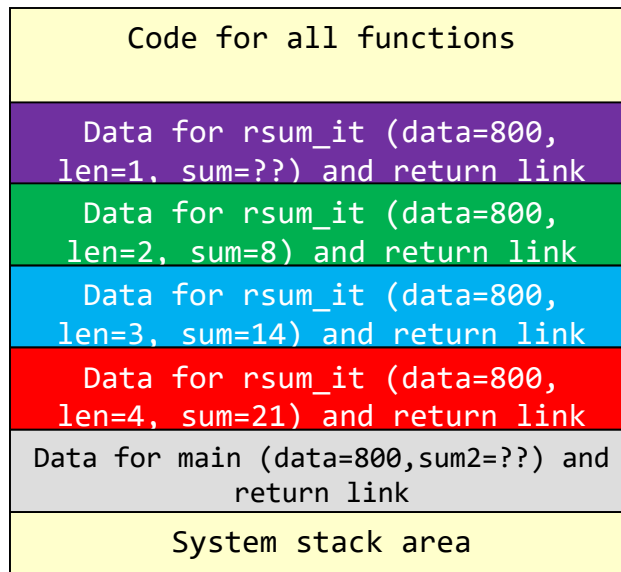
System Stack & Recursion

- The system stack makes recursion possible by providing separate memory storage for the local variables of each running instance of the function

```
int main()
{
    int data[4] = {8, 6, 7, 9};
    int sum2 = rsum_it(data, 4);
}

int rsum_it(int data[], int len)
{
    if(len == 1)
        return data[0];
    else
        int sum =
            rsum_it(data, len-1);
        return sum + data[len-1];
}
```

**System
Memory
(RAM)**



Exercise

- Exercises
 - Count-down
 - Count-up

Recursion Double Check

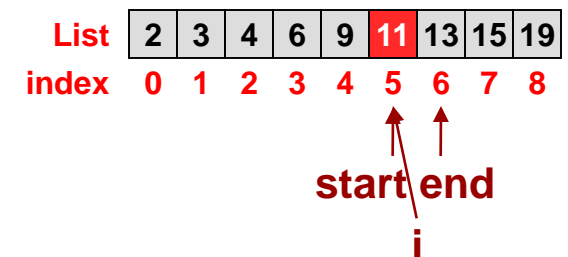
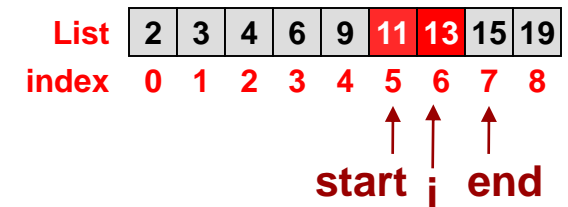
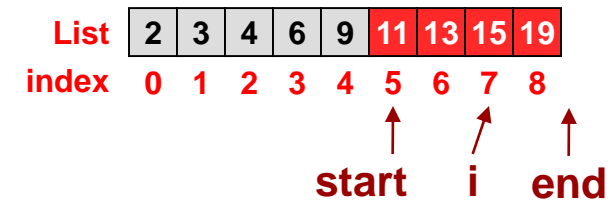
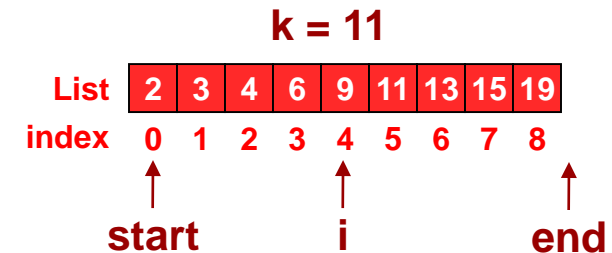
- When you write a recursive routine:
 - Check that you have appropriate base cases
 - Need to check for these first before recursive cases
 - Check that each recursive call makes progress toward the base case
 - Otherwise you'll get an infinite loop and stack overflow
 - Check that you use a 'return' statement at each level to return appropriate values back to each recursive call
 - You have to return back up through every level of recursion, so make sure you are returning something (the appropriate thing)

Loops & Recursion

- Is it better to use recursion or iteration?
 - ANY problem that can be solved using recursion can also be solved with iteration and other appropriate data structures
- Why use recursion?
 - Usually clean & elegant. Easier to read.
 - Sometimes generates much simpler code than iteration would
 - Sometimes iteration will be almost impossible
 - The power of recursion often comes when each function instance makes **multiple** recursive calls
- How do you choose?
 - Iteration is usually faster and uses less memory
 - However, if iteration produces a very complex solution, consider recursion

Recursive Binary Search

- Assume remaining items = [start, end)
 - start is inclusive index of start item in remaining list
 - End is exclusive index of start item in remaining list
- binSearch(target, List[], start, end)
 - Perform base check (empty list)
 - Return NOT FOUND (-1)
 - Pick mid item
 - Based on comparison of k with List[mid]
 - EQ => Found => return mid
 - LT => return answer to BinSearch[start,mid)
 - GT => return answer to BinSearch[mid+1,end)



Sorting

- If we have an unordered list, sequential search becomes our only choice
- If we will perform a lot of searches it may be beneficial to sort the list, then use binary search
- Many sorting algorithms of differing complexity (i.e. faster or slower)
- Bubble Sort (simple though not terribly efficient)
 - On each pass through thru the list, pick up the maximum element and place it at the end of the list. Then repeat using a list of size $n-1$ (i.e. w/o the newly placed maximum value)

List

7	3	8	6	5	1
---	---	---	---	---	---

index 0 1 2 3 4 5
Original

List

3	7	6	5	1	8
---	---	---	---	---	---

index 0 1 2 3 4 5
After Pass 1

List

3	6	5	1	7	8
---	---	---	---	---	---

index 0 1 2 3 4 5
After Pass 2

List

3	5	1	6	7	8
---	---	---	---	---	---

index 0 1 2 3 4 5
After Pass 3

List

3	1	5	6	7	8
---	---	---	---	---	---

index 0 1 2 3 4 5
After Pass 4

List

1	3	5	6	7	8
---	---	---	---	---	---

index 0 1 2 3 4 5
After Pass 5

Exercise

- Exercises
 - Text-based fractal

Flood Fill

- Imagine you are given an image with outlines of shapes (boxes and circles) and you had to write a program to shade (make black) the inside of one of the shapes. How would you do it?
- Flood fill is a recursive approach
- Given a pixel
 - Base case: If it is black already, stop!
 - Recursive case: Call floodfill on each neighbor pixel
 - Hidden base case: If pixel out of bounds, stop!

