EE 457 Unit 0

Class Introduction

Basic Hardware Organization

EE 457 – Computer Systems Organization

- Computer Architecture class
- Focus on CPU Design
  - Microarchitecture
  - General Digital System Design
- Focus on Memory Hierarchy
  - Cache
  - Virtual Memory
- Focus on Computer Arithmetic
  - Fast Adders
  - Fast Multipliers

Course Info

- Lecture:
  - Prof. Redekopp (redekopp@usc.edu)
- Discussion:
  - TA: See website
- Website:

  http://bytes.usc.edu/ee457
  https://courses.uscden.net/d2l/home

- Midterm (30%):
- Final (31%):
- Homework Assignments (14%): Individual
- Lab Assignments (25%): Individual and Teams of 2
  - Contact TA

Prerequisites

- EE 354L "Introduction to Digital Circuits"
  - Logic design
  - State machine implementation
  - Datapath/control unit implementation
  - Verilog HDL
- EE 109/354 "Basic Computer Organization"
  - Assembly language programming
  - Basic hardware organization and structures
- C or similar high-level programming knowledge
- Familiarity with Verilog HDL

EE 109/354 Required Knowledge

- You must know and understand the following terms and concepts; please review them as necessary
  - Bit, Nibble (four bit word), Byte, Word (16- or 32-bit value)
  - CPU, ALU, CU (Control Unit), ROM, RAM (RWM), Word length of a computer, System Bus (Address, Data, Control)
  - General Purpose Registers, Instruction Register (IR), Program Counter (PC), Stack, Stack Pointer (SP) Subroutine calls, Flag register (or Condition Code Register or Processor Status Word), Microprogramming
  - Instruction Set, Addressing Modes, Machine Language, Assembly Language, Assembler, High Level Language, Compiler, Linker, Object code, Loader
  - Interrupts, Exceptions, Interrupt Vector, Vectored Interrupts, Traps

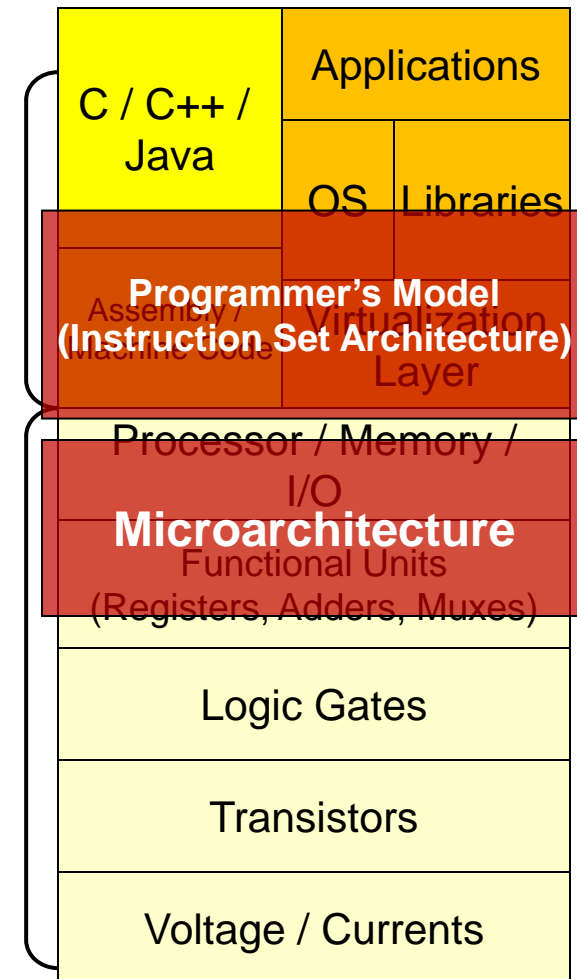# EE 354L Requisite Knowledge

- You must know and understand the following terms and concepts; please review them as necessary
  - Combinational design of functions specified by truth tables and function tables
  - Design of adders, comparators, multiplexers, decoders, demultiplexers
  - Tri-state outputs and buses
  - Sequential Logic components: D-Latches and D-Flip-Flops, counters, registers
  - State Machine Design: State diagrams, Mealy vs. Moore-style outputs, Input Function Logic, Next State Logic, State Memory, Output Function Logic, power-on reset state
  - State Machine Design using encoded state assignments vs. one-hot state assignment
  - Drawing, interpretation, and analysis of waveform diagrams

Computer Arithmetic Requisite Knowledge

- You must know and understand the following terms and concepts; please review them as necessary
  - Unsigned and Signed (2's complement representation) Numbers
  - Unsigned and signed addition and subtraction
  - Overflow in addition and subtraction
  - Multiplication
  - Booth's algorithm for multiplications of signed numbers
  - Restoring or Non-Restoring Division for unsigned numbers
  - Hardware implementations for adders and multipliers

Levels of Architecture

- System architecture
  - High-level HW org.
- Instruction Set Architecture
  - A contract or agreement about what the HW will support and how the programmer can write SW for the HW
  - Vocabulary that the HW understands and SW is composed of
- Microarchitecture
  - HW implementation for executing instructions
  - Usually transparent to SW programs but not program performance
  - Example: Intel and AMD have different microarchitectures but support essentially the same instruction set

SW

HW

C / C++ / Java

Applications

OS | Libraries

Assembly / Machine Code

Virtualization Layer

Programmer's Model (Instruction Set Architecture)

Processor / Memory / I/O

Microarchitecture

Functional Units (Registers, Adders, Muxes)
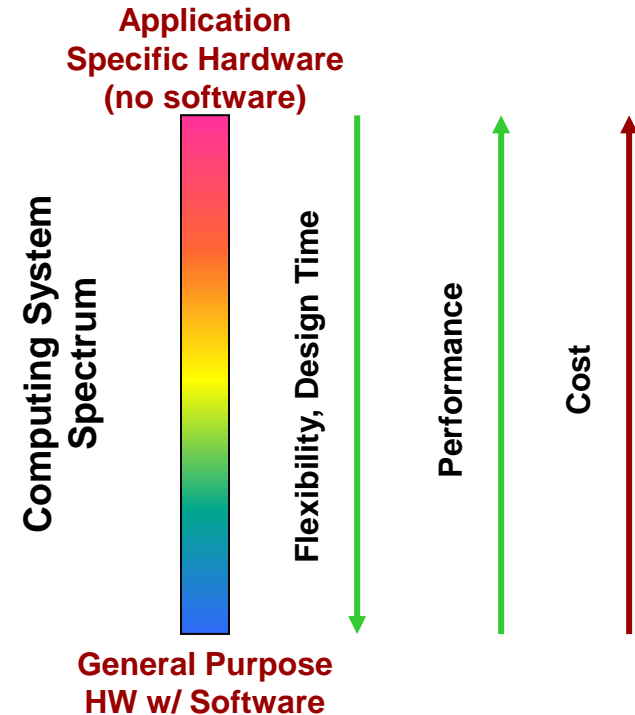
Logic Gates

Transistors

Voltage / Currents

Why is Architecture Important

- Enabling ever more capable computers
- Different systems require different architectures
  - PC's
  - Servers
  - Embedded Systems
    - Simple control devices like ATM's, toys, appliances
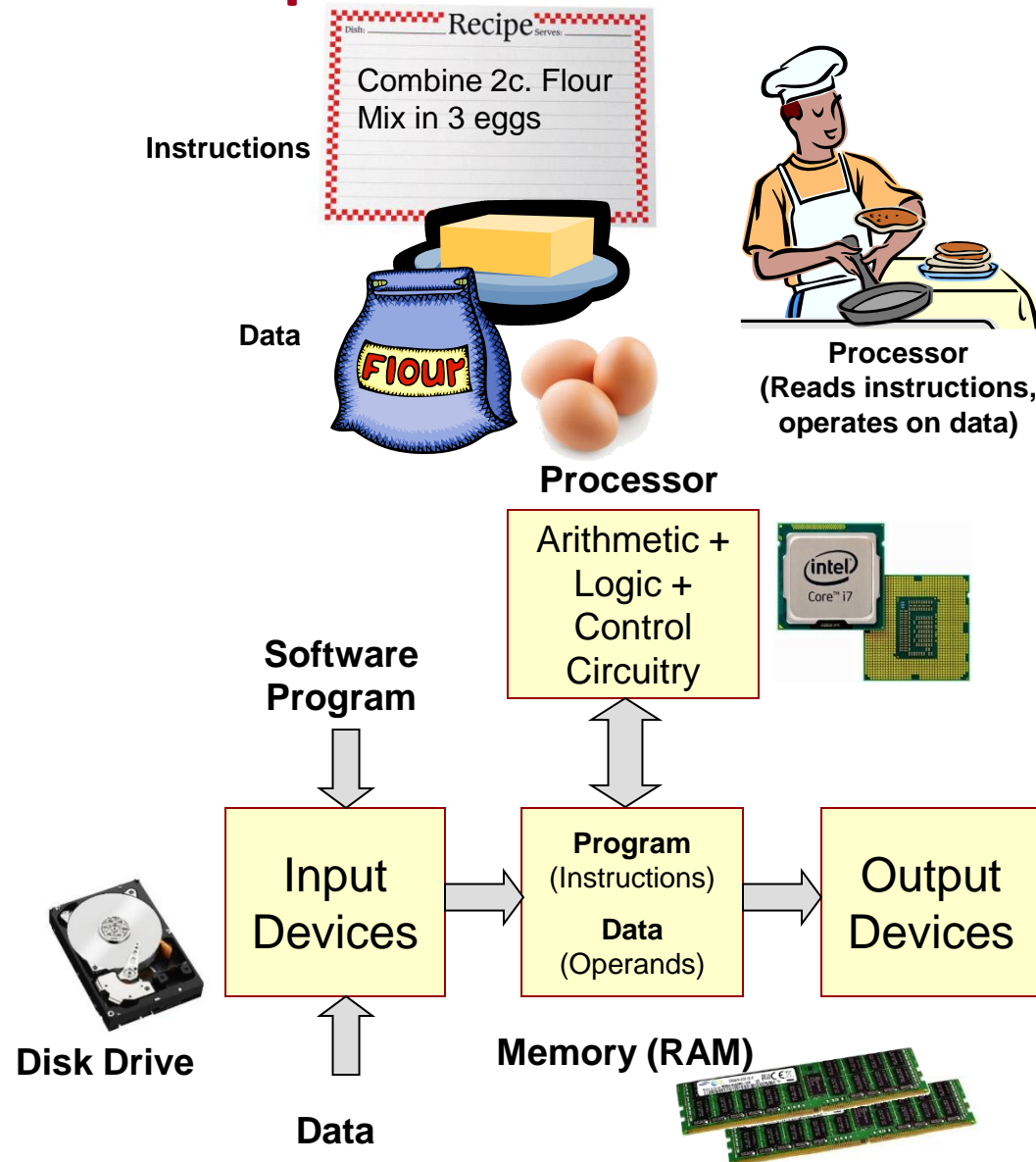   - Media systems like game consoles and MP3 players
    - Robotics

Digital System Spectrum

- Key idea: Any "algorithm" can be implemented in HW or SW or some mixture of both
- A digital systems can be located anywhere in a spectrum of:
  - ALL HW: (a.k.a. Application-Specific IC's)
  - ALL SW: An embedded computer system
- Advantages of application specific HW
  - Faster, less power
- Advantages of an embedded computer system (i.e. general purpose HW for executing SW)
  - Reprogrammable (i.e. make a mistake, fix it)
  - Less expensive than a dedicated hardware system (single computer system can be used for multiple designs)
- Phone: System-on-Chip (SoC) approach
  - Some dedicated HW for sound, graphics, and other operations
  - Programmable processor for UI & other simple tasks

Application Specific Hardware (no software)

Computing System Spectrum

Flexibility, Design Time

Performance

Cost

General Purpose HW w/ Software

Computer Components

- Processor
  - Executes the program and performs all the operations
- Main Memory
  - Stores *data* and *program* (*instructions)*
  - Different forms:
    - RAM = read and write but volatile (lose values when power off)
    - ROM = read-only but non-volatile (maintains values when power off)
  - Significantly slower than the processor speeds
- Input / Output Devices
  - Generate and consume data from the system
  - MUCH, MUCH slower than the processor

Instructions

Combine 2c. Flour
Mix in 3 eggs

Data

**Processor
(Reads instructions, operates on data)**

Processor

Arithmetic + Logic + Control Circuitry

Software Program

Program
(Instructions)
Data
(Operands)

Input Devices

Output Devices

Disk Drive

Memory (RAM)

Data

Drivers and Trends

ARCHITECTURE OVERVIEW

Architecture Issues

- Fundamentally, architecture is all about the different ways of answering the question:

  "What do we do with the ever-increasing number of transistors available to us"

- Goal of a computer architect is to take increasing transistor budgets of a chip (i.e. Moore's Law) and produce an equivalent increase in computational ability

Moore's Law, Computer Architecture & Real-Estate Planning

- Moore's Law = Number of transistors able to be fabricated on a chip grows exponentially with time

- Computer architects decide, "What should we do with all of this capability?"

- Similarly real-estate developers ask, "How do we make best use of the land area given to us?"
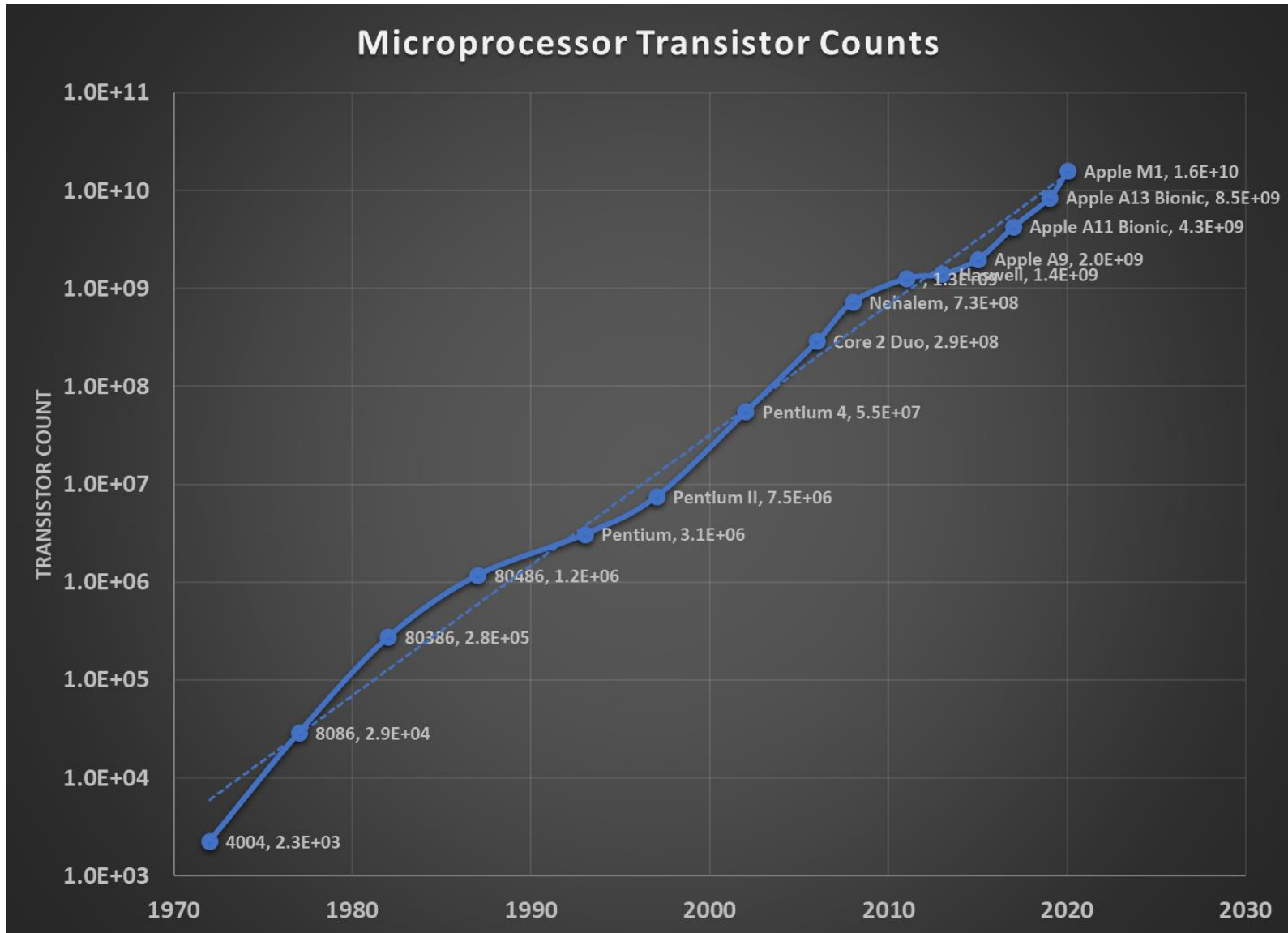

USC University Park Development Master Plan
http://re.usc.edu/docs/University%20Park%20Development%20Project.pdf

Transistor Physics

- Cross-section of transistors on an IC

- Moore's Law is founded on our ability to keep shrinking transistor sizes
  - Gate/channel width shrinks
  - Gate oxide shrinks

- Transistor feature size is referred to as the implementation "technology node"

Technology Nodes

Process Technology Progression
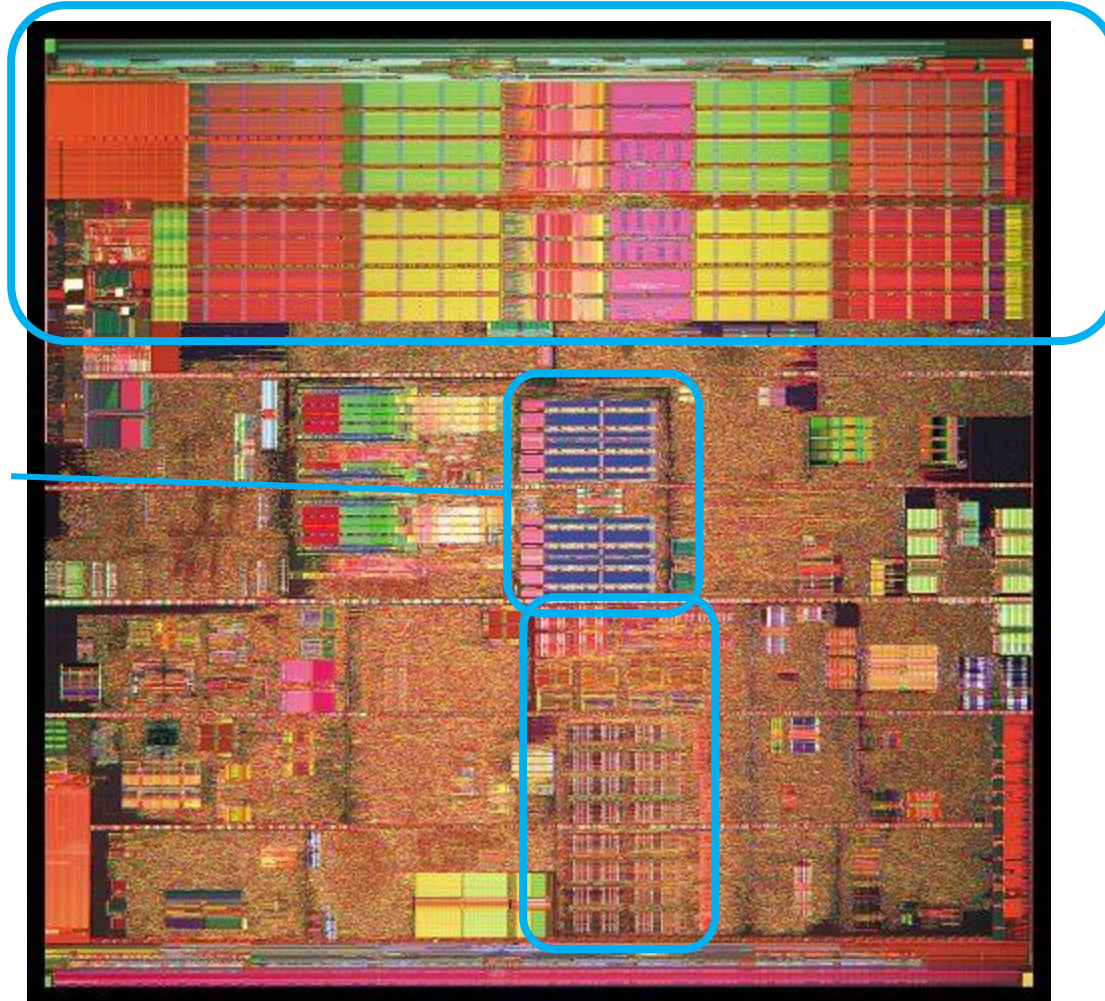
Growth of Transistors on Chip

Microprocessor Transistor Counts

Implications of Moore's Law

- What should we do with all these transistors
  - Put additional simple cores on a chip
  - Use transistors to make cores execute instructions faster
  - Use transistors for more on-chip cache memory
    - Cache is an on-chip memory used to store data the processor is likely to need
    - Faster than main-memory which is on a separate chip and much larger (thus slower)
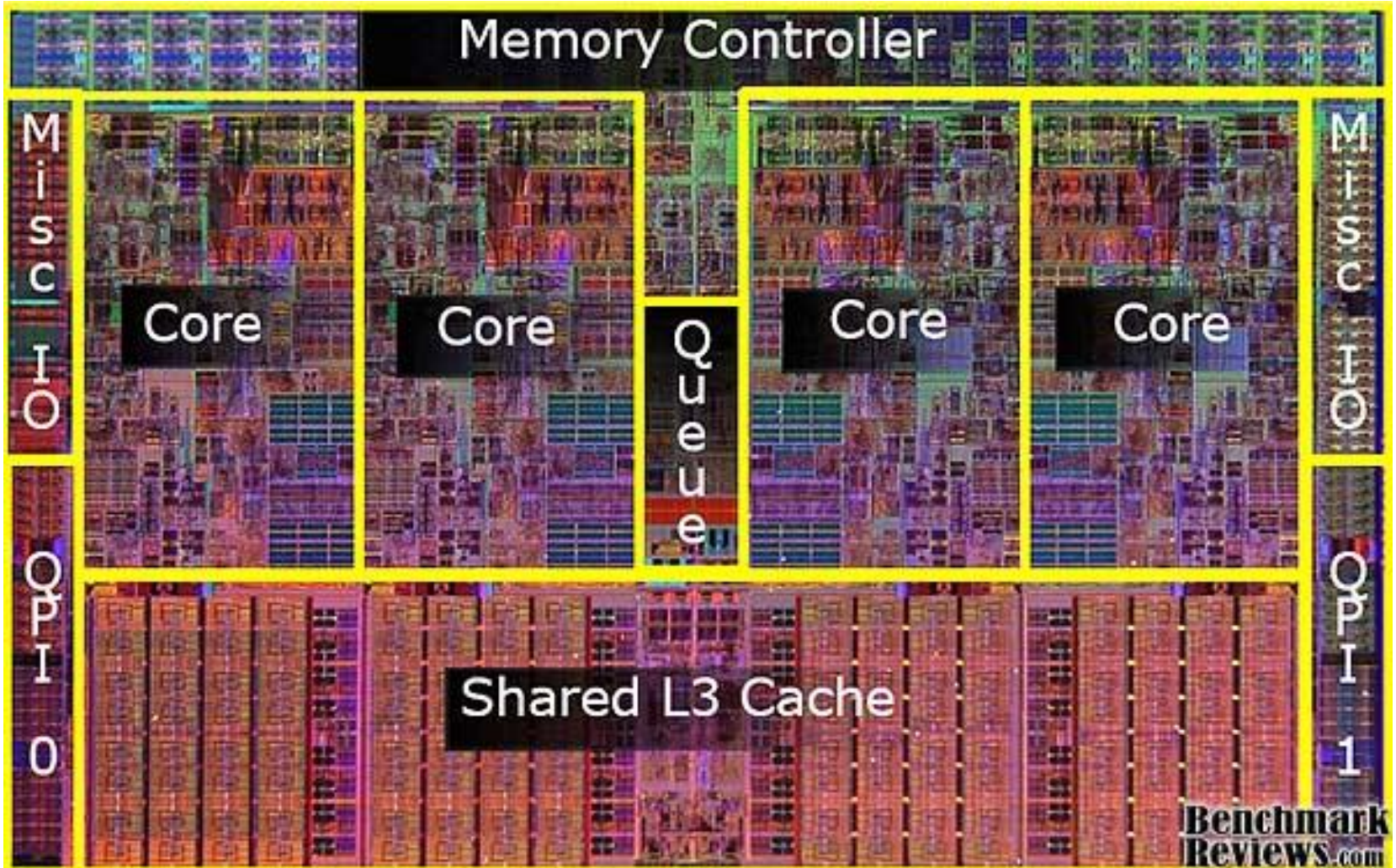
Pentium 4

L2 Cache

L1 Data

L1 Instruc.

Increase in Clock Frequency

Microprocessor Clock Frequencies

What happened here?

Intel Nehalem Quad Core

Progression to Parallel Systems

- If power begins to limit clock frequency, how can we continue to achieve more and more operations per second?
  - By running several processor cores in parallel at lower frequencies
  - Two cores @ 2 GHz vs. 1 core @ 4 GHz yield the same theoretical maximum ops./sec.
- We'll end our semester by examining (briefly) a few parallel architectures
  - Chip multiprocessors (multicore)
  - Graphics Processor Units (SIMT)

Flynn's Taxonomy

- Categorize architectures based on relationship between program (instructions) and data

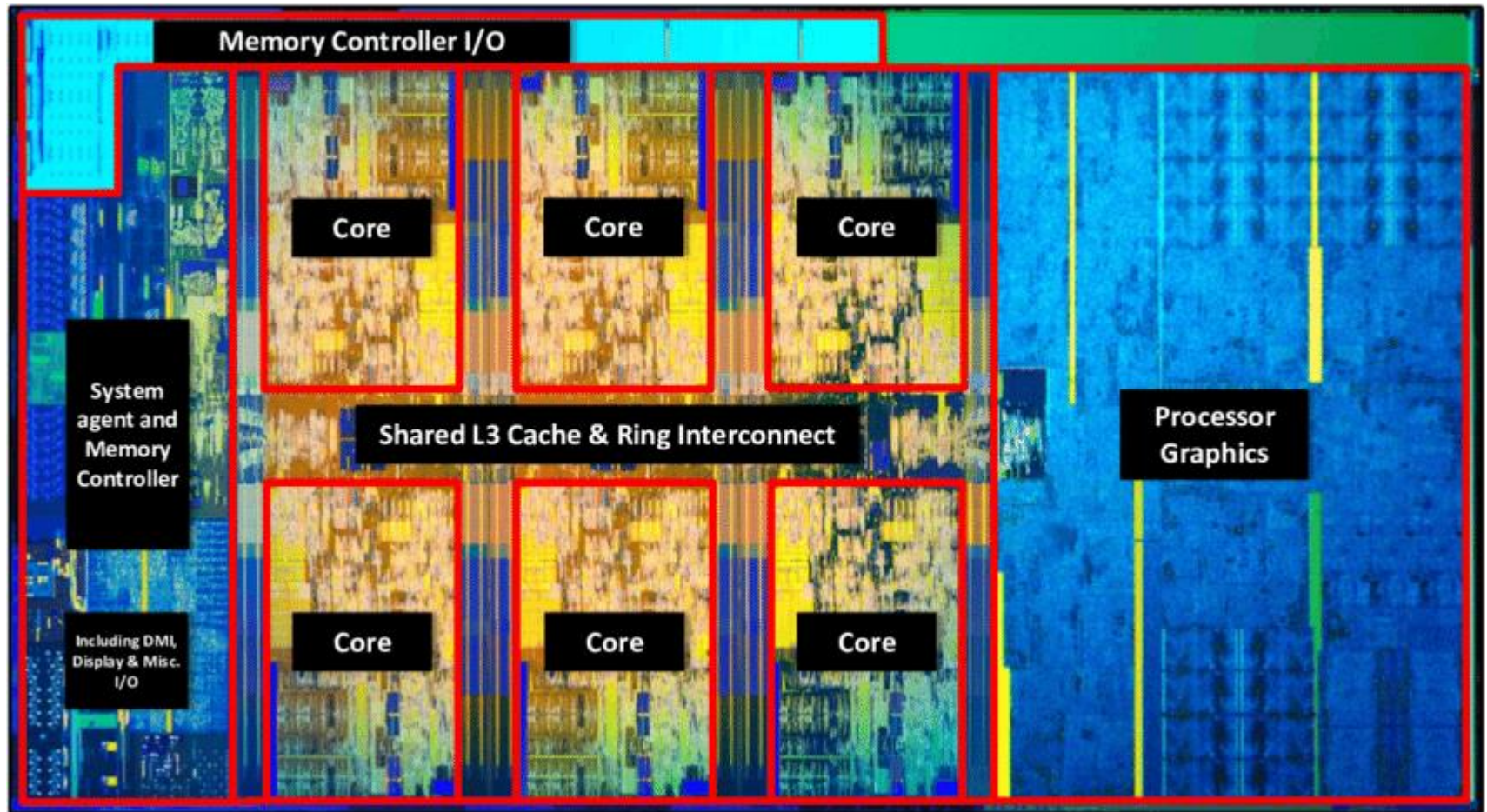| SISD Single-Instruction, Single-Data	SIMD / SIMT Single Instruction, Multiple Data (Single Instruction, Multiple Thread)
• Typical, single-threaded processor	• Vector Units (e.g. Intel MMX, SSE, SSE2) • GPU's
MISD Multiple Instruction, Single-Data	MIMD Multiple Instruction, Multiple-Data
• Less commonly used (some streaming architectures may be considered in this category)	• Multi-threaded processors • Typical CMP/Multicore system (Task parallelism with different threads executing)

GPU Chip Layout

- 2560 Small Cores

- Upwards of 7.2 billion transistors

- 8.2 TFLOPS

- 320 Gbytes/sec

Source: NVIDIA

Photo: http://www.theregister.co.uk/2010/01/19/nvidia_gf100/
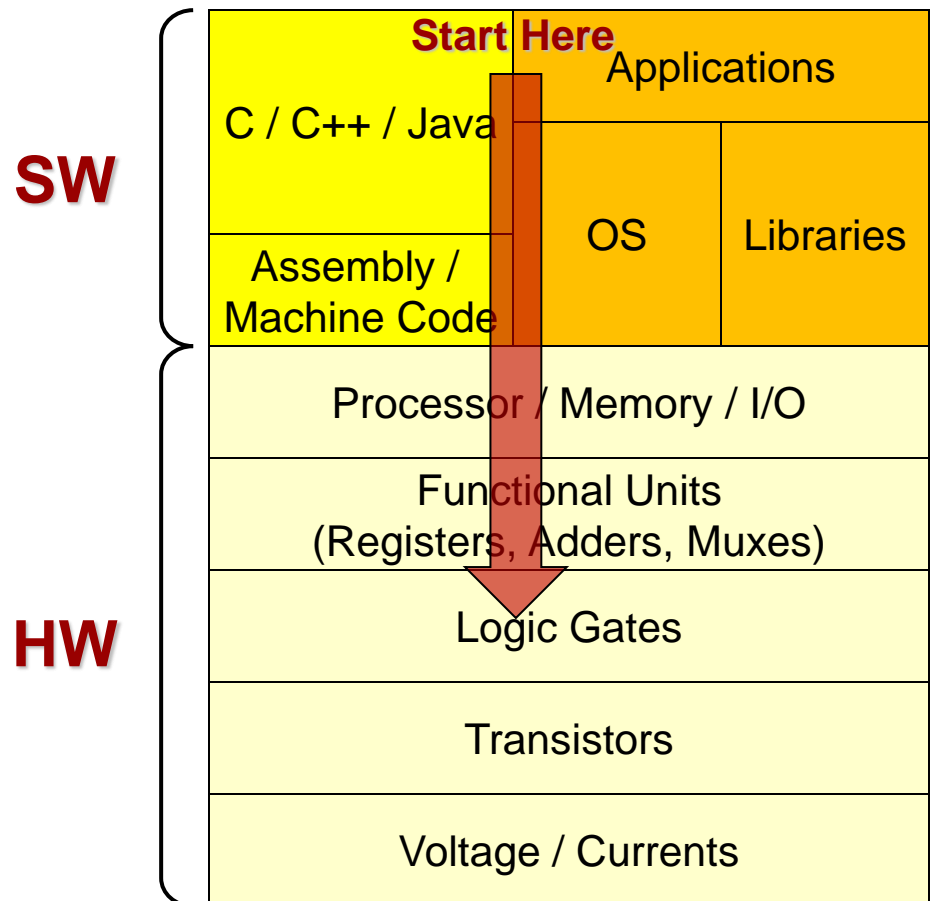
8th Gen Coffee-Lake Hex-Core Intel Processor

https://www.researchgate.net/figure/Die-Map-of-a-Hexa-Core-Coffee-Lake-Processor_fig6_332543387

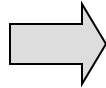In case you need a review…Look these over on your own

COMPUTER SYSTEM TOUR

Computer Systems Tour

- How does a SW program get mapped and executed on a computer

- What components make a computer system and what are their functions

- How does the architecture affect performance

SW

HW

Start Here		Applications	
C / C++ / Java			
		OS	Libraries
Assembly / Machine Code			

| Processor / Memory / I/O |
| Functional Units (Registers, Adders, Muxes) |
| Logic Gates |
| Transistors |
| Voltage / Currents |

Software Process

Software Program

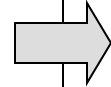**High Level Language Description**

```
if (x > 0)
  x = x + y - z;
a = b*x;
```

.c/.cpp files

Compiler

```
        MOVE.L  X,D0
        CMPI    #0,D0
        BLE     SKIP
        ADD     Y,D0
        SUB     Z,D0
SKIP  MUL       …
```
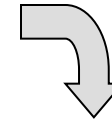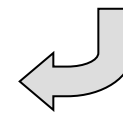
Assembly (.asm/.s files)

Assembler

```
1110 0010 0101 1001
0110 1011 0000 1100
0100 1101 0111 1111
1010 1100 0010 1011
0001 0110 0011 1000
```

Object/Machine Code (.o files)

In EE 357 you will be able to perform all the tasks of the compiler…

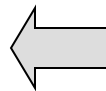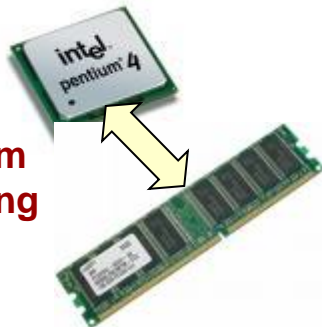**A "compiler" (i.e. gcc, VisualC++, etc.) includes the assembler & linker**

```
1110 0010 0101 1001
0110 1011 0000 1100
0100 1101 0111 1111
1010 1100 0010 1011
0001 0110 0011 1000
```

Executable Binary Image

Linker

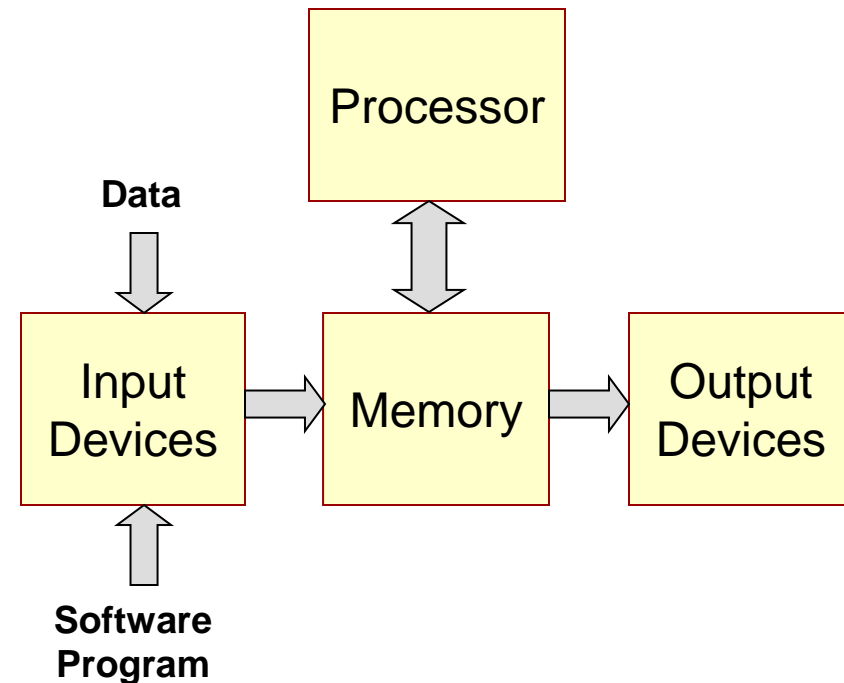**Program Executing**

Loader / OS

Compiler Process

- A compiler such as 'gcc' performs 3 tasks:
  - Compiler
    - Converts HLL (high-level language) files to assembly
  - Assembler
    - Converts assembly to object (machine) code
  - Static Linker
    - Links multiple object files into a single executable resolving references between code in the separate files
  - Output of a compiler is a binary image that can be loaded into memory and then executed.

- Loader/Dynamic Linker
  - Loads the executable image into memory and resolves dynamic calls (to OS subroutines, libraries, etc.)
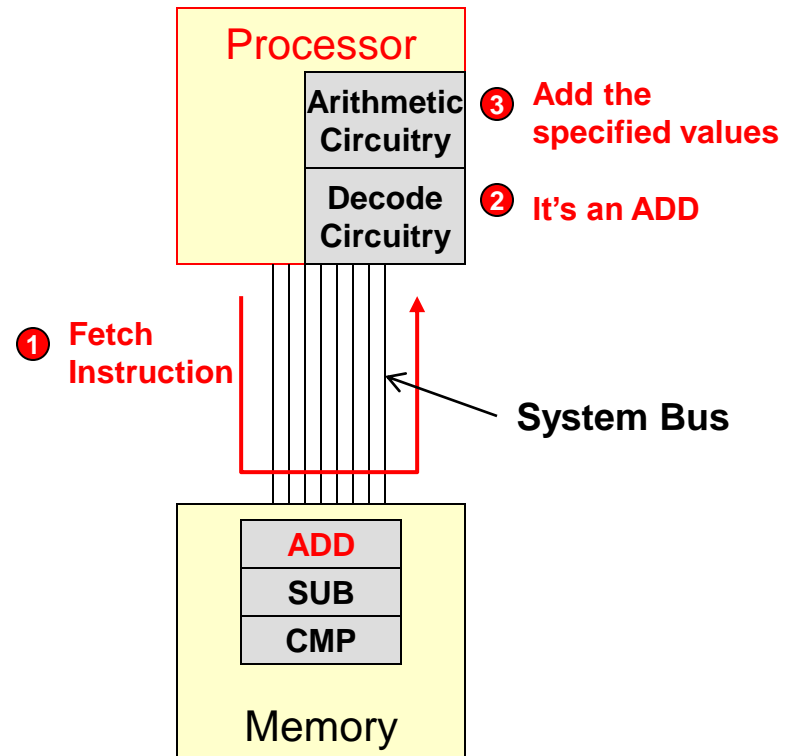
Hardware Components

- Processor
  - Executes the program and performs all the operations
  - Examples: Pentium 4, PowerPC, M68K/Coldfire
- Main Memory
  - Stores *data* and *program* (*instructions)*
  - Different forms:
    - RAM = read and write but volatile (lose values when power off)
    - ROM = read-only but non-volatile (maintains values when power off)
  - Significantly slower than the processor speeds
- Input / Output Devices
  - Generate and consume data from the system
  - Examples: Keyboard, Mouse, CD-ROM, Hard Drive, USB, Monitor display
  - MUCH, MUCH slower than the processor

Processor

Data

Input Devices

Memory

Output Devices

Software Program

Processor

- Performs the same 3-step process over and over again
  - Fetch an instruction from memory
  - Decode the instruction
    - Is it an ADD, SUB, etc.?
  - Execute the instruction
    - Perform the specified operation
- This process is known as the **Instruction Cycle**

Processor

Arithmetic Circuitry ③ Add the specified values

Decode Circuitry ② It's an ADD

① Fetch Instruction
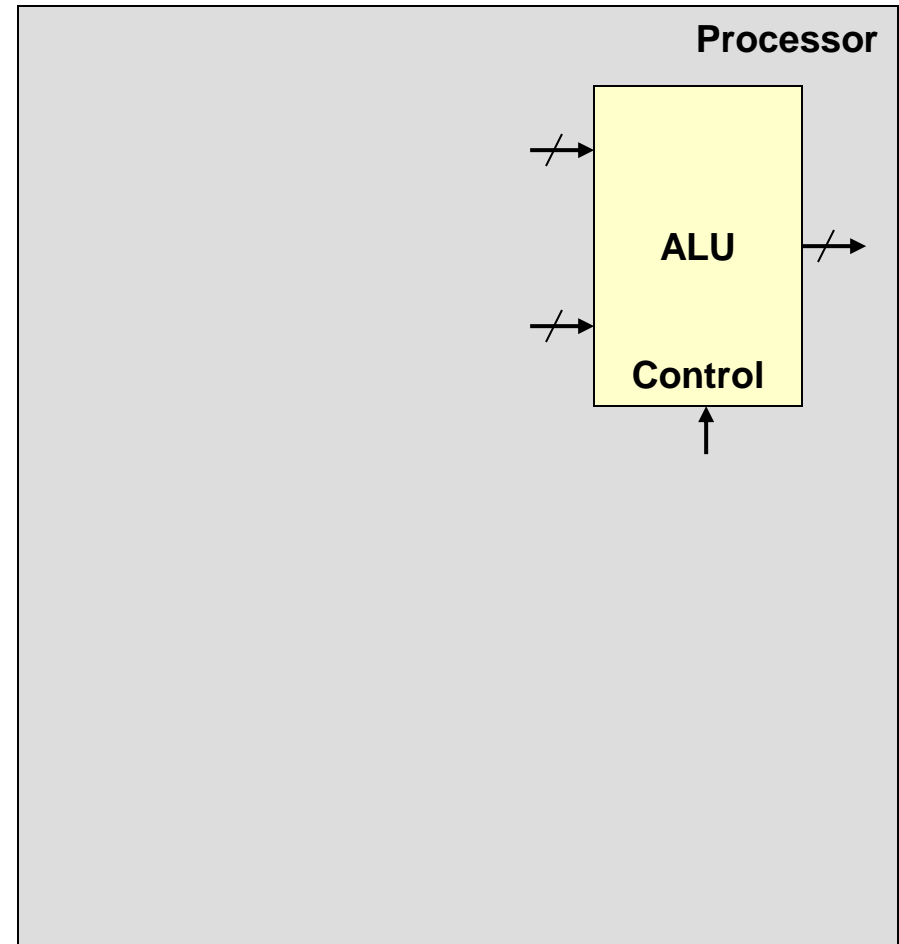
System Bus

ADD
SUB
CMP

Memory

Processors

- Processors contain 4 subcomponents
  1. ALU (Arithmetic & Logical Unit)
  2. Registers
  3. Control Circuitry & System-Bus Interface
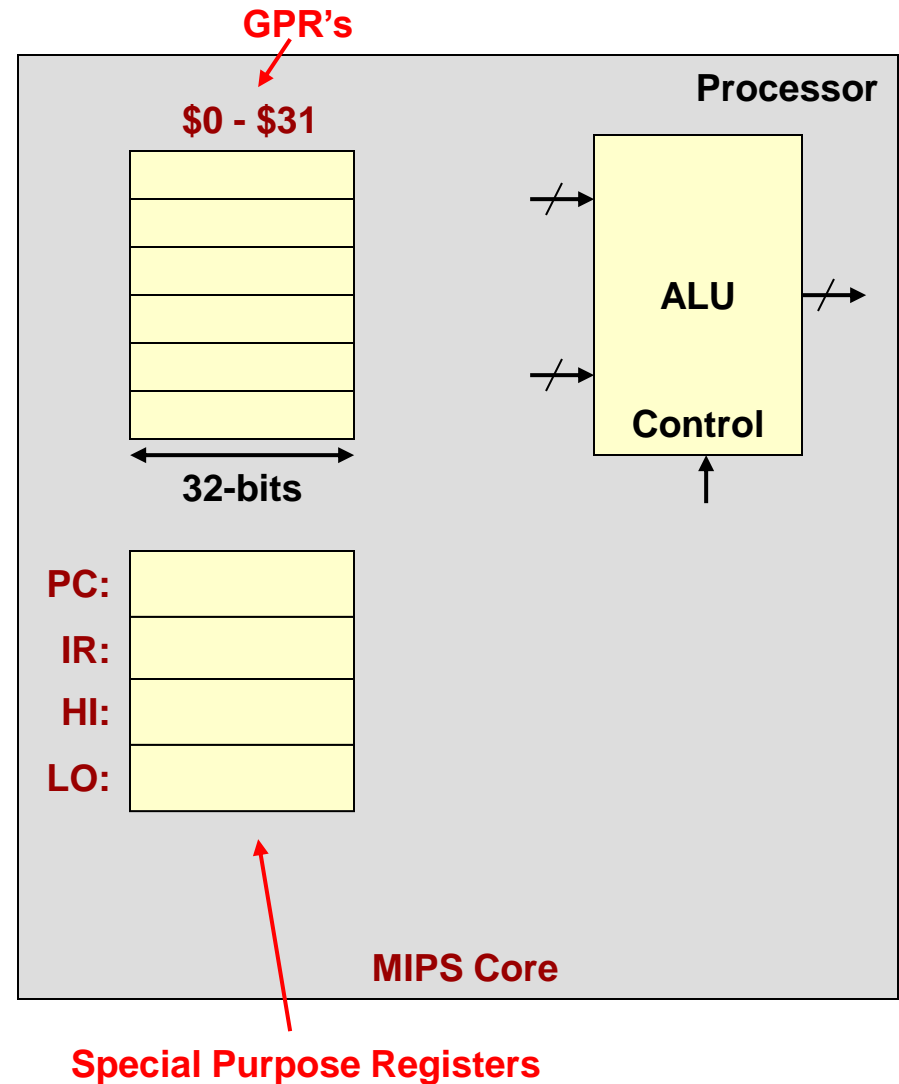 4. Cache (Optional)

ALU

- Performs arithmetic and logical operations

- 2 inputs and 1 output value

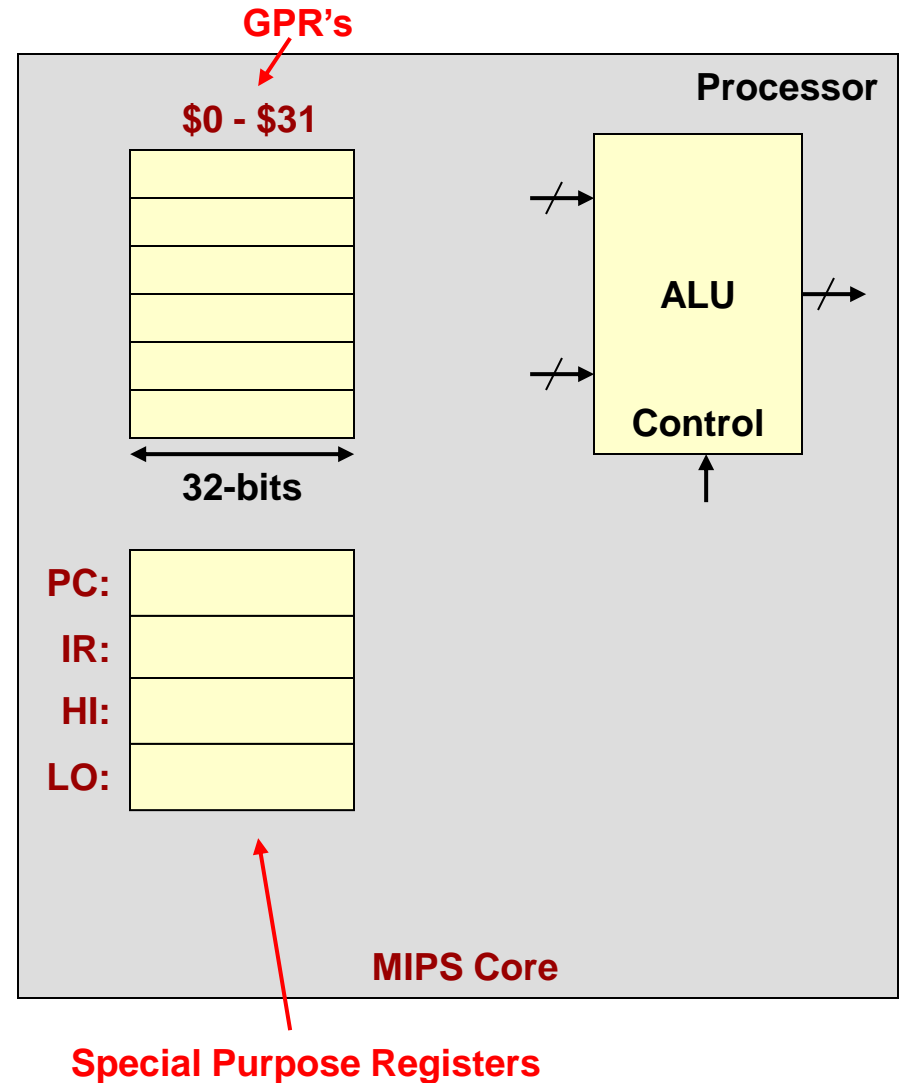- Control inputs to select operation (ADD, SUB, AND, OR...)

Registers

- Provide temporary storage for data
- 2 categories of registers
  - General Purpose Registers (GPR's)
    - for program data
    - can be used by programmer as desired
    - given names (e.g. D0-D7)
  - Special Purpose Registers
    - for internal processor operation (not for program data)

GPR's

Processor

$0 - $31

ALU

Control

32-bits

PC:

IR:

HI:

LO:

MIPS Core
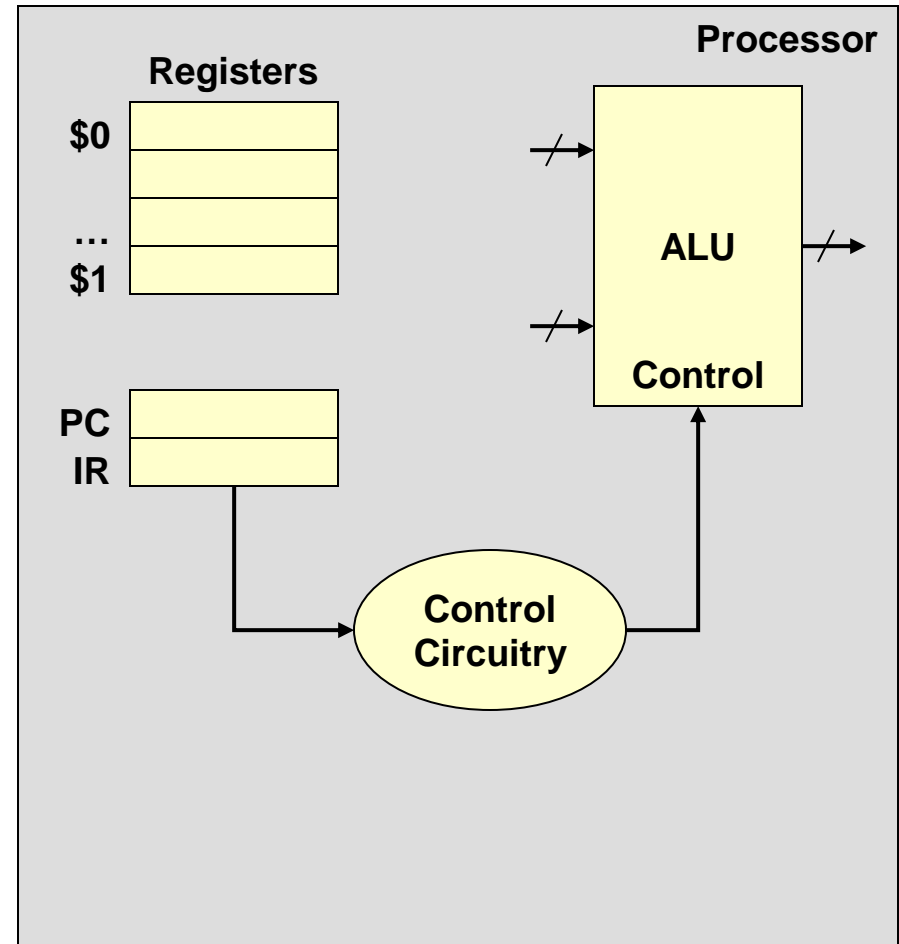
Special Purpose Registers

Registers

- GPR's
  - Faster to access than main memory
  - Keep data you are working with in registers to speed up execution
- Special Purpose Reg's.
  - Hold specific information that the processor needs to operate correctly
  - PC (Program Counter)
    - Pointer to (address of) instruction in memory that will be executed next
  - IR (Instruction Register)
    - Stores the instruction while it is being executed
  - SR (Status Register)
    - Stores status/control info

GPR's

Processor

$0 - $31

ALU

Control

32-bits

PC:

IR:

HI:

LO:

MIPS Core
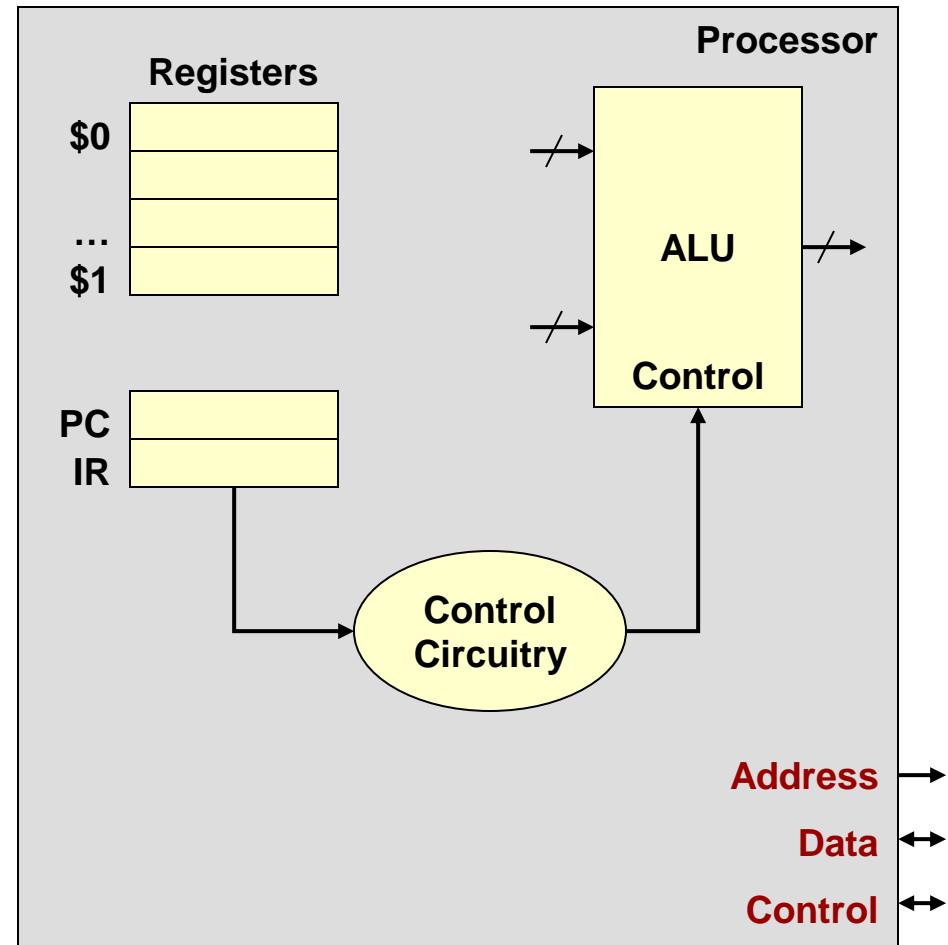
Special Purpose Registers

Control Circuitry

- Decodes each instruction

- Selects appropriate registers to use

- Selects ALU operation

- And more…

System Bus Interface

- ## System bus is the means of communication between the processor and other devices

  - Address
    - Specifies location of instruction or data
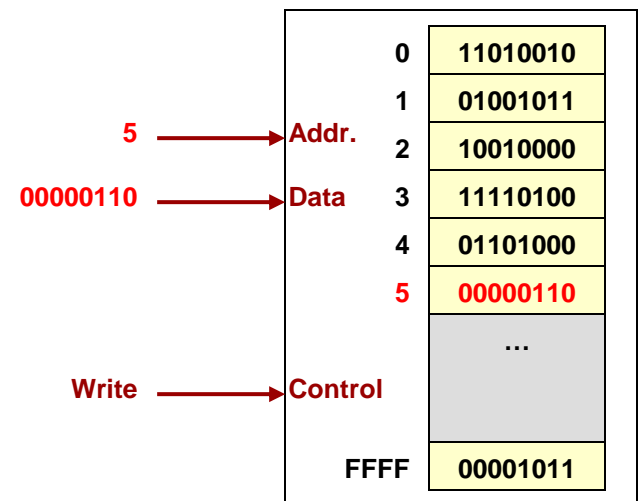 - Data
  - Control

Memory

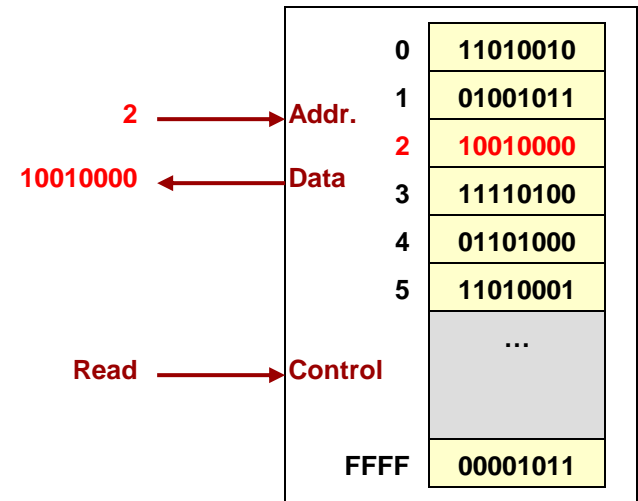- Set of cells that each store a group of bits (usually, 1 byte = 8 bits)

- Unique address assigned to each cell

  – Used to reference the value in that location

- Numbers and instructions are all represented as a string of 1's and 0's

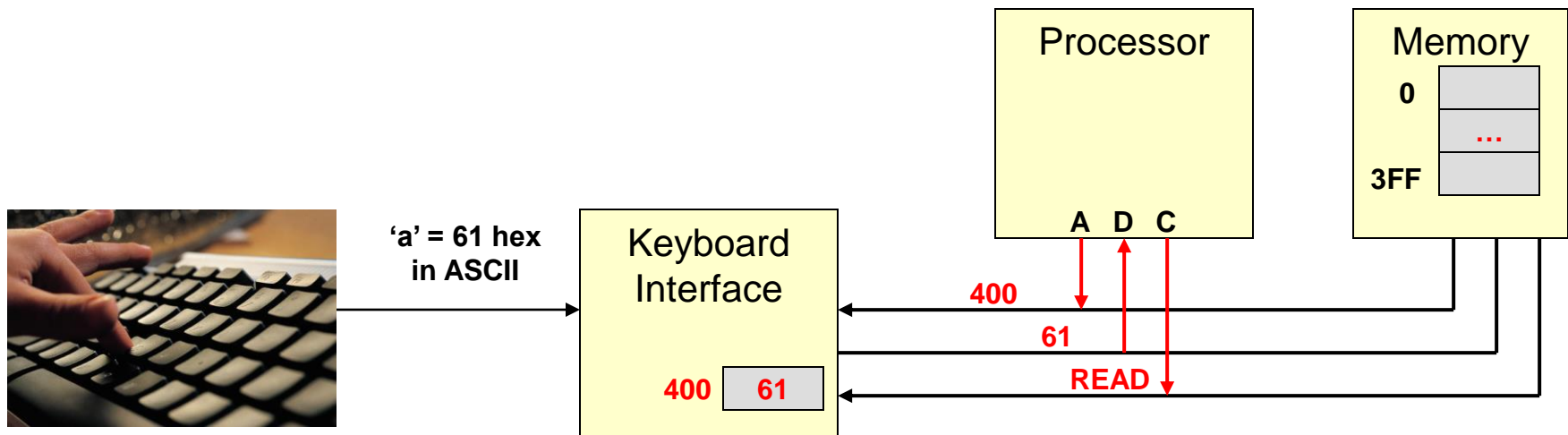| Address	Data
0	11010010
1	01001011
2	10010000
3	11110100
4	01101000
5	11010001
	...
FFFF	00001011

Memory Device

Memory Operations

- Memories perform 2 operations
  - Read: retrieves data value in a particular location (specified using the address)
  - Write: changes data in a location to a new value
- To perform these operations a set of address, data, and control inputs/outputs are used
  - Note: A group of wires/signals is referred to as a 'bus'
  - Thus, we say that memories have an address, data, and control bus.

2 → Addr.	0 — 11010010
	1 — 01001011
10010000 ← Data	2 — 10010000
	3 — 11110100
	4 — 01101000
	5 — 11010001
	...
Read → Control	
	FFFF — 00001011

A Read Operation

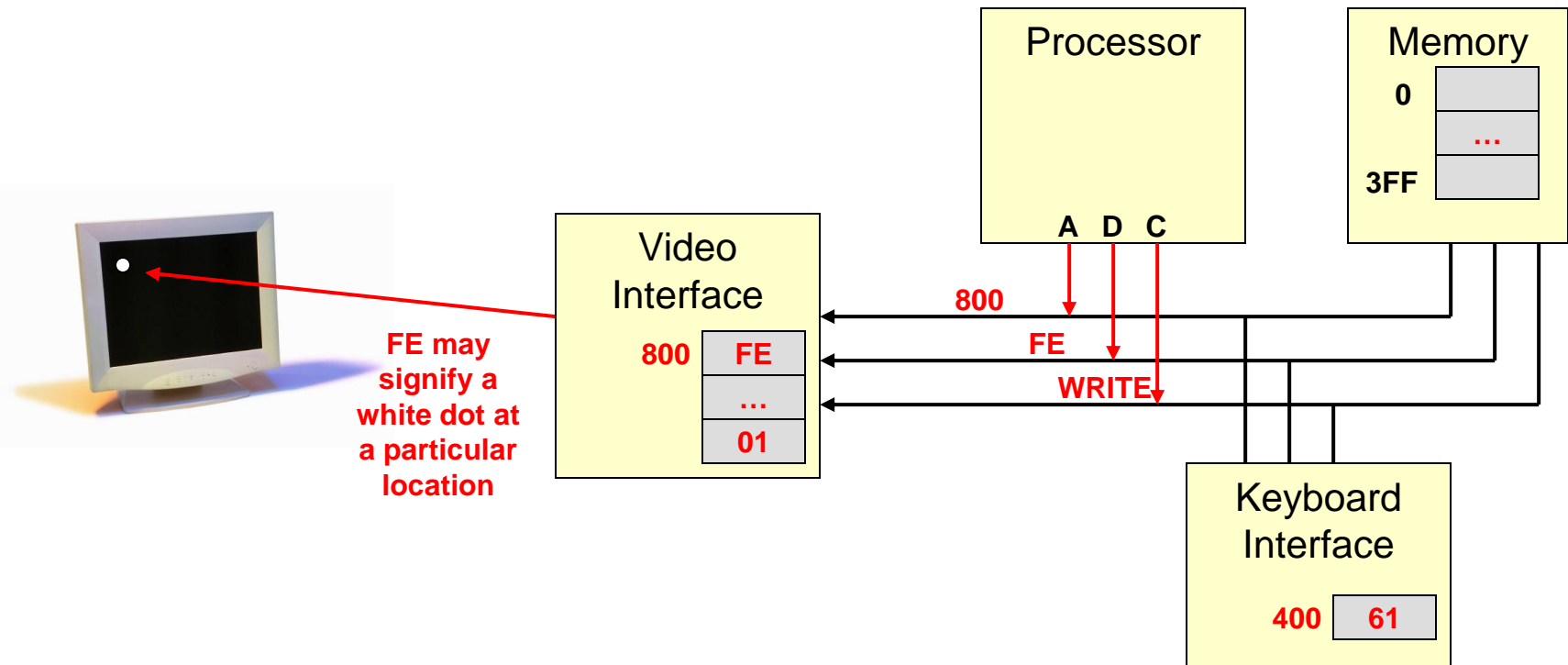|	
	0 — 11010010
	1 — 01001011
5 → Addr.	2 — 10010000
00000110 → Data	3 — 11110100
	4 — 01101000
	5 — 00000110
	...
Write → Control	
	FFFF — 00001011

A Write Operation

Input / Output

- Keyboard, Mouse, Display, USB devices, Hard Drive, Printer, etc.
- Processor can perform reads and writes on I/O devices just as it does on memory
  - I/O devices have locations that contain data that the processor can access
  - These locations are assigned unique addresses just like memory

Input / Output

- Writing a value to the video adapter can set a pixel on the screen

FE may signify a white dot at a particular location

Video Interface

800 | FE |
| ... |
| 01 |

Processor

A   D   C

800

FE

WRITE

Memory

0 | |
| ... |
3FF | |

Keyboard Interface

400 | 61 |

Computer Organization Issues

- Components run at different speeds
  - Processor can perform operations very quickly (~ 1 ns)
  - Memory is much slower (~ 50 ns) due to how it is constructed & its shear size [i.e. it must select/look-up 1 location from millions]
    - Speed is usually inversely proportional to size (i.e. larger memory => slower)
  - I/O devices are much slower
    - Hard Drive (~ 1 ms)
  - Intra-chip signals (signals w/in the same chip) run much faster than inter-chip signals
- Design HW and allocate HW resources to accommodate these inherent speed differences