

## Spiral 3-1

Hardware/Software Interfacing

## Learning Outcomes

- I understand the PicoBlaze bus interface signals: PORT\_ID, IN\_PORT, OUT\_PORT, WRITE\_STROBE
- I understand how a memory map provides the agreement between addresses the software will use and that the hardware must recognize and respond to
- I understand how to build address decoding logic to ensure only the appropriate value/register is selected for a given PORTID
- For output, I can take a memory map and the PORTID and OUTDATA bits such that the appropriate data is input or saved in a register when an OUTPUT instruction is executed
- For input, I can take a memory map and the appropriate PORTID bits to build logic and muxes such that the appropriate data value is present at INDATA when an INPUT instruction is executed

## ASICS & FPGAS REVIEW

## Digital Design Targets

- Two possible implementation targets
  - Custom Chips (ASIC's = Application Specific Integrated Circuits): Physical gates are created on silicon to implement 1 particular design
  - FPGA (Field Programmable Gate Array's): "Programmable logic" using programmable memories to implement logic functions along with other logic resources tiled on the chip. Can implement any design and then be changed to implement a new one

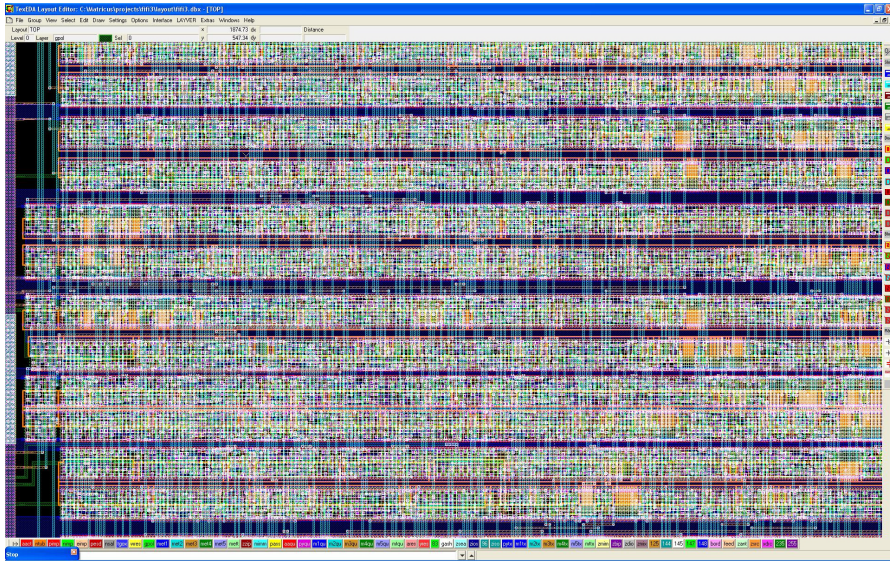


In an ASIC design, a unique chip will be manufactured that implements our design at which point the HW design is fixed & cannot be changed (example: Pentium, etc.)



FPGA's have "logic resources" on them that we can configure to implement our specific design. We can then reconfigure it to implement another design

## ASICs



## Implementation

- ASIC's
  - Use the CAD tools to synthesize and route a “netlist”
    - **Synthesis** = Takes logic description or logic schematic & converts to transistor level gates
    - **Place and Route** = Figure out where each gate should go on the chip)
  - Final “netlist” is sent to chip maker for production
  - Fabrication is very expensive (> \$1 million) so get your design right the first time.
- FPGA's
  - Synthesis converts logic description to necessary LUT contents, etc.
  - Place and route produces a configuration for the FPGA chip
  - Can reconfigure FPGA as much as you like, so less important to get it right 1<sup>st</sup> time

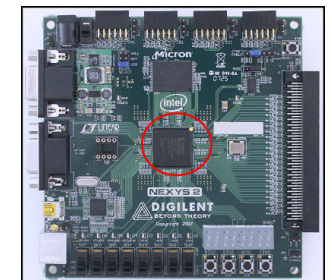
## ASIC's vs. FPGA's

- |   |                   |
|---|-------------------|
| • ASIC's  | • FPGA's          |
| – _____   | – _____           |
| – Handles _____ Designs   | – _____ Designs   |
| – _____ Expensive   | – _____ Expensive |
| – _____ Flexible (Cannot be _____ to perform a new hardware function) | – _____ Flexible  |

## Xilinx Spartan 3E

### Digilent Nexys-2 Board

- Has a Xilinx Spartan 3E FPGA (XC3S500e)
- 500K gate equivalent
- 9312 D-FF's on-board



### On-board I/O

- (4) 7-Segment Displays
- (8) LED's
- (4) Push Buttons
- (8) Switches



# Latest FPGA's

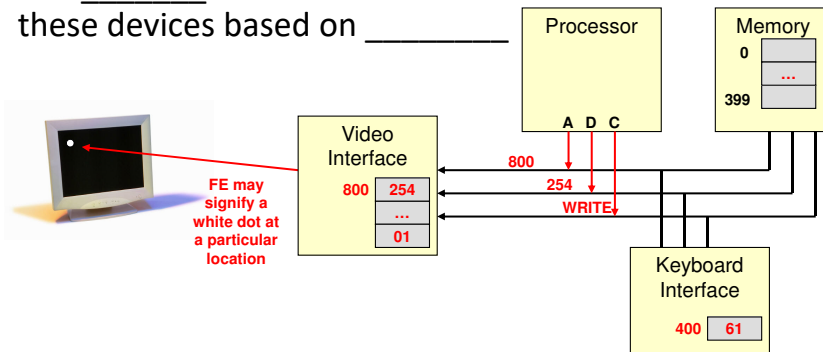
- SoC design (Xilinx Kintex [KU115])
  - Quad-Core ARM cores
  - DDR3 SDRAM Memory Interface
  - ~800 I/O Pins
  - Equiv. ~15M gate equivalent FPGA fabric
    - ~1M D-FFs + 552K LUTs
    - 1968 dedicated DSP "slices" 18x18 multiply + adder
    - 34.6 Megabits of onboard Block RAMs

Hardware/Software Interfacing

# PICOBLAZE

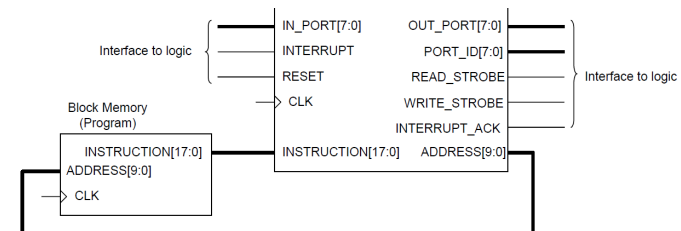
# Input / Output

- The processor connects to peripherals and other logic via the \_\_\_\_\_ (address, data, and control)
- Software running on the processor performs \_\_\_\_\_ and \_\_\_\_\_ that read and write data to and from these devices based on \_\_\_\_\_



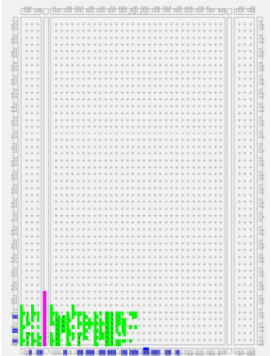
# Introduction

- Picoblaze (aka KCPSM3) is an 8-bit \_\_\_\_\_
  - The processor is not implemented directly in hardware on the FPGA but instead is just a description that is then \_\_\_\_\_ using the same process as any of our other designs
  - It provides a bus interface that can be connected to custom logic that you design and then used to control that custom logic via \_\_\_\_\_ executing on the processor



## KCPSM3 is small!

This plot from the Xilinx Floorplanner shows the same implementation of KCPSM3 in an XC3S200 Spartan-3 device. This makes it easier to appreciate the actual logic resources required by the macro without the interconnect obscuring the detail.

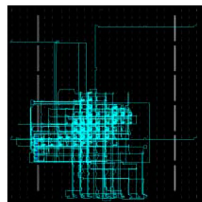


The placement in this Floorplanner view was achieved using a simple area constraint in the project UCF file.

```
INST processor_* LOC=SLICE_X0Y0:SLICE_X19Y4;
```

Such constraints are not required in normal designs and it has only been used in this case because so little of the device is occupied. Experiments have shown that placement constraints have very little effect on performance.

The FPGA Editor view shown to the right was the result when no constraints were used. The size is still 96 slices but this is now a little less obvious! The performance was actually a little higher than when using the area constraint indicating that a 'tidy' design is not always the fastest!



Taken from the KCPSM3 Manual

## Size and Performance

The following device resource information is taken from the ISE reports for the KCPSM3 macro in an XC3S200 device. The reports reveal the features that are utilised and the efficiency of the macro. The 96 'slices' reported by the MAP process in this case may reduce to the minimum of 89 'slices' when greater packing is used to fit a complete design into a device.

### XST Report

LUT1	: 2	} 109 LUTs (55 slices)
LUT2	: 6	
LUT3	: 68	
LUT4	: 33	
MUXCY	: 39	} Carry and MUX logic (Free with LUTs)
MUXF5	: 9	
XORCY	: 35	
FD	: 24	} 76 Flip_flops (Free with LUTs)
FDE	: 2	
FDR	: 30	
FDRRE	: 8	
FDRSE	: 10	
FDS	: 2	
RAM16X1D	: 8	— Register bank (8 slices)
RAM32X1S	: 10	— Call/Return Stack (10 slices)
RAM64X1S	: 8	— Scratch Pad Memory (16 slices)

Total = 89 Slices

### MAP Report

Number of occupied Slices : 96 out of 1920 5%  
 Number of Block RAMs : 1 out of 12 8%  
 Total equivalent gate count for design: 74,814  
 12 x KCPSM3 can fit into the XC3S200 device (40% of the logic slices remaining). An equivalent gate count of 897,768 gates in a 200,000 gate device!

### TRACE Report

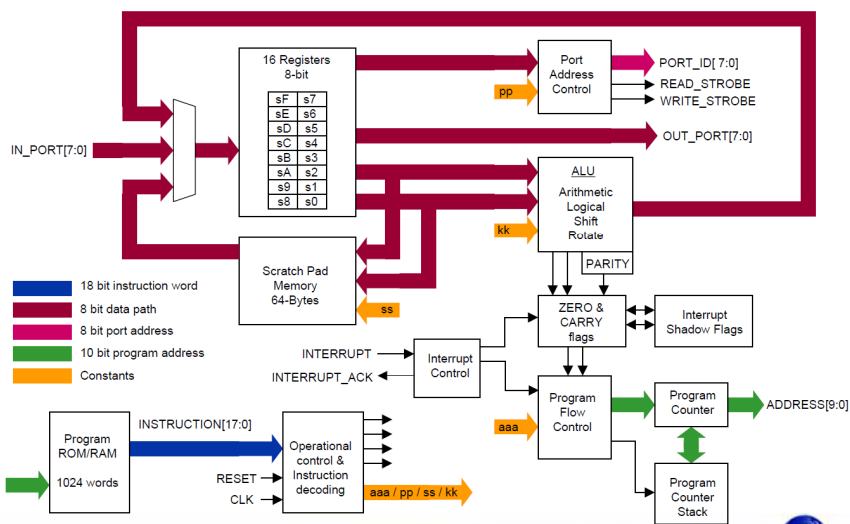
Device, speed: xc3s200,-4 (PREVIEW 1.22 2003-03-16)  
 Minimum period: 11.403ns  
 (Maximum frequency: 87.696MHz)  
 43.8 MIPS

### TRACE Report for Virtex-IIPro

Device, speed: xcvp2,-7 (ADVANCED 1.76 2003-03-16)  
 Minimum period: 7.505ns  
 (Maximum frequency: 133.245MHz) 66.6 MIPS

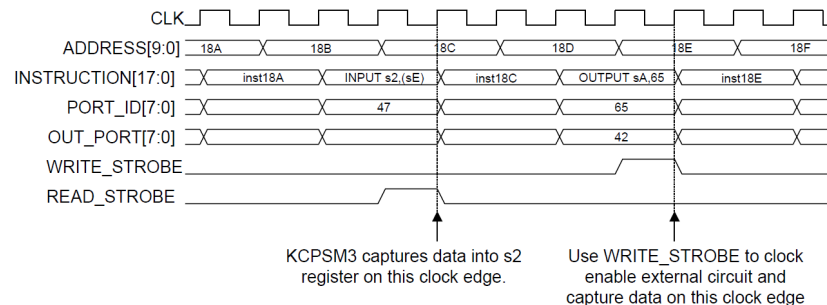
Taken from the KCPSM3 Manual

## KCPSM3 Architecture



Taken from the KCPSM3 Manual

## Input / Output Operations



KCPSM3 captures data into s2 register on this clock edge.  
 Use WRITE\_STROBE to clock enable external circuit and capture data on this clock edge

Taken from the KCPSM3 Manual

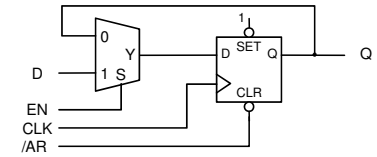
## Exercise 1

- Make the register below capture data (**out\_data**) from your Picoblaze whenever it outputs address FF hex on (**address** or **port\_id**)



## Remember: Registers w/ Enables

- Registers (D-FF's) will sample the D bit every clock edge and pass it to Q
- Sometimes we may want to hold the value of Q and ignore D even at a clock edge
- We can add an enable input and some logic in front of the D-FF to accomplish this

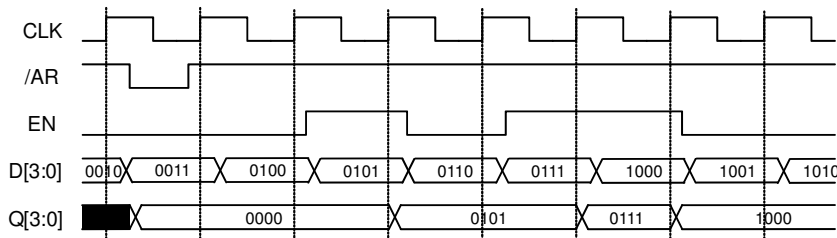


FF with Data Enable

CLK	/AR	EN	D <sub>i</sub>	Q <sub>i</sub> *
X	0	X	X	0
0,1	1	X	X	Q <sub>i</sub>
↑	1	0	X	Q <sub>i</sub>
↑	1	1	0	0
↑	1	1	1	1

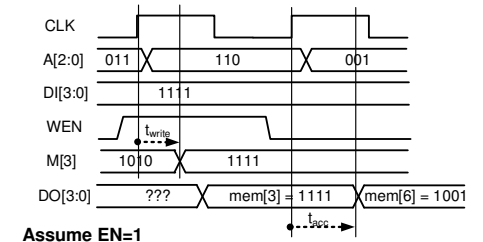
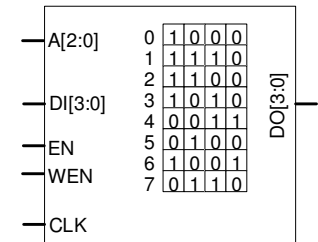
## Registers w/ Enables

- The D value is sampled at the clock edge only if the enable is active
- Otherwise the current Q value is maintained



## Recall Memory Interfaces

- We provide address and data
- EN = Overall enable (unless it is 1) the memory won't read or write (we assume EN=1)
- WEN = Write enable – 1 = Write / 0 = read

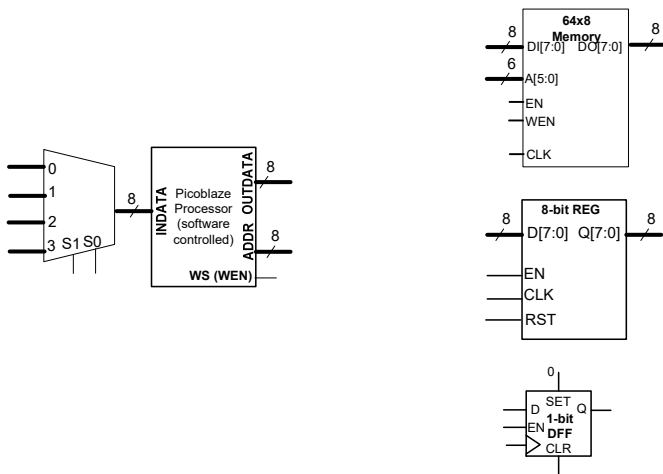


Assume EN=1





## Address Decoding Exercise 2



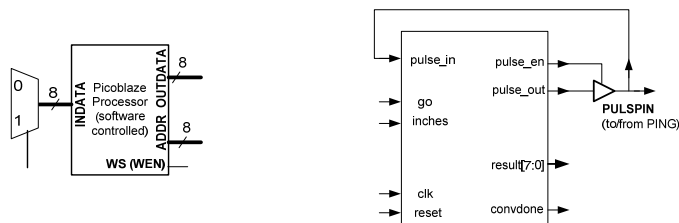
## High-Order Interleaving

- General strategy is to place devices at ranges of address divide by **more-significant bits**

Dec	A7	A6	A5	A4	A3	A2	A1	A0	Assigned Device
00			0	0	0	0	0	0	64x8
01			0	0	0	0	0	1	Memory
...									
63	0	0	1	1	1	1	1	1	
...									open
128		0	0	0	0	0	0	0	8-bit Register
...									open
192	1	1	0	0	0	0	0	0	1-bit D-FF
...									open

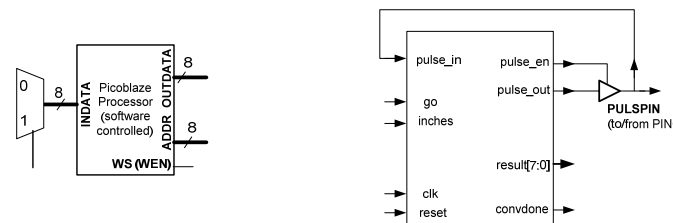
## PING))) Interfacing

- Work with your instructor to explore alternatives for interfacing the PING))) engine you created to the Picoblaze processor



## PING))) Interfacing

- Explore other alternatives...





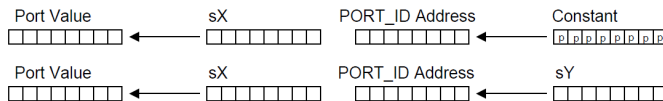
## PICOBLAZE ASSEMBLY

## Relevant Manual Pages

- Pages 4-6, (7), 8, (9-11)
- Pages 16-36, focus on
  - OUTPUT 34
  - INPUT 35
  - SHIFTS 32,33
  - JUMP 17
  - LOAD 22
  - COMPARE 31
- Input/Output design 65-68

## Output Instruction OUTPUT

The OUTPUT instruction enables the contents of any register to be transferred to logic external to KCPSM3. The port address (in the range 00 to FF) can be defined by a constant value or indirectly as the contents of any other register. The Flags are not affected by this operation.



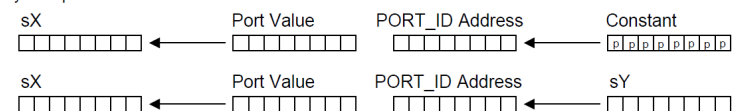
The user interface logic is required to decode the PORT\_ID port address value and capture the data provided on the OUT\_PORT. The WRITE\_STROBE is set during an output operation (see 'READ and WRITE STROBES'), and should be used to clock enable the capture register or write enable a RAM (see 'Design of Output Ports').

- Example: **output s1, FF**
  - Outputs the 8-bit number in **s1** as data on **out\_port** to the address (**port\_id**) of 0xFF

Taken from the KCPSM3 Manual

## Input Instruction INPUT

The INPUT instruction enables data values external to KCPSM3 to be transferred into any one of the internal registers. The port address (in the range 00 to FF) can be defined by a constant value or indirectly as the contents of any other register. The Flags are not affected by this operation.



The user interface logic is required to decode the PORT\_ID port address value and supply the correct data to the IN\_PORT. The READ\_STROBE is set during an input operation (see 'READ and WRITE STROBES'), but it is not always necessary for the interface logic to decode this strobe. However, it can be useful for determining when data has been read, such as when reading a FIFO buffer (see 'Design of Input Ports').

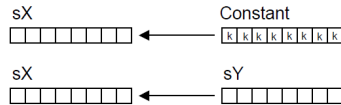
- Example: **input s8, 0c**
  - Places the address 0x0c on the **port\_id** and then grabs data from the **in\_port** at the end of the second cycle and writes it into register **s8**

Taken from the KCPSM3 Manual

# LOAD Instruction

## LOAD

The LOAD instruction provides a method for specifying the contents of any register. The new value can be a constant, or the contents of any other register. The LOAD instruction has no effect on the status of the flags.



Since the LOAD instruction does not effect the flags it may be used to reorder and assign register contents at any stage of the program execution. The ability to assign a constant with no impact to the program size or performance means that the load instruction is the most obvious way to assign a value or clear a register.

The first operand of a LOAD instruction must specify the register to be loaded as register 's' followed by a hexadecimal digit. The second operand must then specify a second register value in a similar way or specify an 8-bit constant using 2 hexadecimal digits. The assembler supports register naming and constant labels to simplify the process.

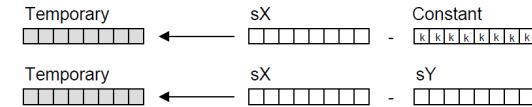
- Example: `load s3, a5`  
– Loads the constant 0xa5 into register s3

Taken from the KCPSM3 Manual

# Compare Instruction

## COMPARE

The COMPARE instruction performs an 8-bit subtraction of two operands. Unlike the 'SUB' instruction, the result of the operation is discarded and only the flags are affected. The ZERO flag is set when all the bits of the temporary result are low and indicates that both input operands were identical. The CARRY flag indicates when an underflow has occurred and indicates that the second operand was larger than the first. For example, if 's05' contains 27 hex and the instruction COMPARE s05,35 is performed, then the CARRY flag will be set (35>27) and the ZERO flag will be reset (35≠27).



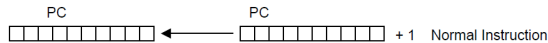
CARRY [?] Set if 'sY' or 'kk' is greater than 'sX'. Reset in all other cases.  
ZERO [?] Set if operands are equal. Reset in all other cases.

- Example: `compare sf, 2a`  
– Compares the data in register `sf` to the hex constant `0x2a`. It sets the Z flag (to determine equality) and C flag (to indicate less-than)  
Taken from the KCPSM3 Manual

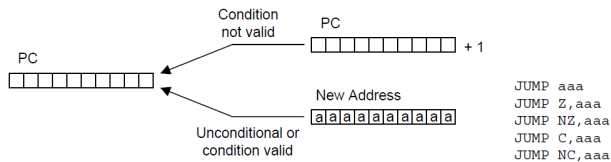
# Jump Instruction

## JUMP

Under normal conditions, the program counter (PC) increments to point to the next instruction. The address space is fixed to 1024 locations (000 to 3FF hex) and therefore the program counter is 10 bits wide. It is worth noting that the top of memory is 3FF hex and will increment to 000.



The JUMP instruction may be used to modify this sequence by specifying a new address. However, the JUMP instruction may be conditional. A conditional JUMP will only be performed if a test performed on either the ZERO flag or CARRY flag is valid. The JUMP instruction has no effect on the status of the flags.



JUMP aaa  
JUMP Z,aaa  
JUMP NZ,aaa  
JUMP C,aaa  
JUMP NC,aaa

- Example: `jump Z, label1`  
– Jumps to the location specified by `label1` if the condition bit (Z) is true  
Taken from the KCPSM3 Manual