

## Unit 1

### Integer Representation

## Skills & Outcomes

- You should know and be able to apply the following skills with confidence
  - Convert an unsigned binary number to and from decimal
  - Understand the finite number of combinations that can be made with n bits
  - Convert a signed (2's complement system) binary number to and from decimal
  - Convert bit sequences to and from hexadecimal
  - Predict the outcome & perform casting operations

## DIGITAL REPRESENTATION

## Information Representation


- All information in a computer system is represented as bits
  - Bit = ( \_\_\_\_\_ ) = 0 or 1
- A single bit is can only represent 2 values so to represent a wider variety of options we use a \_\_\_\_\_ of bits (e.g. 11001010)
  - Commonly sequences are 8-bits (aka a "byte"), 16-, 32- or 64-bits
- Kinds of information
  - Numbers, text, code/instructions, sound, images/videos

# Interpreting Binary Strings

- Given a sequence of 1's and 0's, you need to know the *representation system* being used, before you can understand the value of those 1's and 0's.
- Information (value) = \_\_\_\_\_


**01000001 = ?**

Unsigned  
Binary system




65 decimal

x86 Assembly  
Instruction



inc %ecx  
(Add 1 to the ecx register)

ASCII  
system



'A'<sub>ASCII</sub>

# Binary Representation Systems

- Integer Systems
  - Unsigned
    - Unsigned (Normal) binary
  - Signed
    - Signed Magnitude
    - 2's complement
    - Excess-N\*
    - 1's complement\*
- Floating Point
  - For very large and small (fractional) numbers

- Codes
  - Text
    - ASCII / Unicode
  - Decimal Codes
    - BCD (Binary Coded Decimal) / (8421 Code)

\* = Not covered in this class

# Data Representation

- In C/C++ variables can be of different types and sizes
  - Integer Types on 32-bit (64-bit) architectures

C Type (Signed)	C Type (Unsigned)	Bytes	Bits	x86 Name
char	unsigned char	1	8	byte
short	unsigned short	2	16	_____
int / int32_t †	unsigned / uint32_t †	4	32	_____
long	unsigned long	4 (8)	32 (64)	_____
long long / int64_t †	unsigned long long / uint64_t †	8	64	_____
char*	-	4 (8)	32 (64)	_____
int*	-	4 (8)	32 (64)	_____

† = defined in stdint.h

- Floating Point Types

C Type	Bytes	Bits	x86 Name
float	4	32	single
double	8	64	double

# OVERVIEW

Using power-of-2 place values

## UNSIGNED BINARY TO DECIMAL

## Number Systems

- Unsigned binary follows the rules of positional number systems
- A positional number systems consist of
  1. \_\_\_\_\_
  2. \_\_\_ coefficients [\_\_\_\_\_]
- Humans: Decimal (Base 10): 0,1,2,3,4,5,6,7,8,9
- Computers: Binary (Base 2): 0,1
- Human systems for working with computer systems (shorthand for human to read/write binary)
  - Octal (Base 8): 0,1,2,3,4,5,6,7
  - Hexadecimal (Base 16): \_\_\_\_\_

## Anatomy of a Decimal Number

- A number consists of a string of explicit coefficients (digits).
- Each coefficient has an implicit place value which is a power of the base.
- The value of a decimal number (a string of decimal coefficients) is the sum of each coefficient times its place value

$$(934)_{10} = 9 * \text{_____} + 3 * \text{_____} + 4 * \text{_____} = 934$$

Explicit coefficients

Implicit place values

$$(3.52)_{10} = 3 * 10^0 + 5 * \text{_____} + 2 * \text{_____} = 3.52$$

## Anatomy of an Unsigned Binary Number

- Same as decimal but now the coefficients are 1 and 0 and the place values are the powers of 2

$$(1011)_2 = 1 * \text{_____} + 0 * \text{_____} + 1 * \text{_____} + 1 * \text{_____}$$

Most Significant Digit (MSB)

Least Significant Bit (LSB)

coefficients

place values = powers of 2

## Binary Examples

$$\begin{array}{ccccccc} \underline{1} & \underline{0} & \underline{0} & \underline{1} & \underline{.} & \underline{1} & \\ \hline & 8 & 4 & 2 & 1 & .5 & \end{array} \quad (1001.1)_2 =$$

$$\begin{array}{ccccccc} \underline{1} & \underline{0} & \underline{1} & \underline{1} & \underline{0} & \underline{0} & \underline{0} & \underline{1} & \\ \hline & 128 & 32 & 16 & & & & 1 & \end{array} \quad (10110001)_2 =$$

## General Conversion From Unsigned Base r to Decimal

- An unsigned number in base r has place values/weights that are the powers of the base
- Denote the coefficients as:  $a_i$

$$(a_3 a_2 a_1 a_0 . a_{-1} a_{-2})_r = a_3 * r^3 + a_2 * r^2 + a_1 * r^1 + a_0 * r^0 + a_{-1} * r^{-1} + a_{-2} * r^{-2}$$

Left-most digit =  
Most Significant  
Digit (MSD)

Right-most digit =  
Least Significant  
Digit (LSD)

$$N_r \Rightarrow \text{_____} \Rightarrow D_{10}$$

Number in base r  Decimal Equivalent

## Examples

$$(746)_8 =$$

$$=$$

$$(1A5)_{16} =$$

$$=$$

$$(AD2)_{16} = 10 * 16^2 + 13 * 16^1 + 2 * 16^0$$

$$= 2560 + 208 + 2 = (2770)_{10}$$

"Making change"

## UNSIGNED DECIMAL TO BINARY

## Decimal to Unsigned Binary

- To convert a decimal number,  $x$ , to binary:
  - Only coefficients of 1 or 0. So simply find place values that add up to the desired values, starting with larger place values and proceeding to smaller values and place a 1 in those place values and 0 in all others

$$25_{10} = \frac{\quad}{32} \frac{\quad}{16} \frac{\quad}{8} \frac{\quad}{4} \frac{\quad}{2} \frac{\quad}{1}$$

For  $25_{10}$  the place value 32 is too large to include so we include 16. Including 16 means we have to make 9 left over. Include 8 and 1.

## Decimal to Unsigned Binary

$$73_{10} = \frac{\quad}{128} \frac{\quad}{64} \frac{\quad}{32} \frac{\quad}{16} \frac{\quad}{8} \frac{\quad}{4} \frac{\quad}{2} \frac{\quad}{1}$$

$$87_{10} = \frac{\quad}{128} \frac{\quad}{64} \frac{\quad}{32} \frac{\quad}{16} \frac{\quad}{8} \frac{\quad}{4} \frac{\quad}{2} \frac{\quad}{1}$$

$$145_{10} = \frac{\quad}{128} \frac{\quad}{64} \frac{\quad}{32} \frac{\quad}{16} \frac{\quad}{8} \frac{\quad}{4} \frac{\quad}{2} \frac{\quad}{1}$$

$$0.625_{10} = \frac{\quad}{.5} \frac{\quad}{.25} \frac{\quad}{.125} \frac{\quad}{.0625} \frac{\quad}{.03125}$$

## Decimal to Another Base

- To convert a decimal number,  $x$ , to base  $r$ :
  - Use the place values of base  $r$  (powers of  $r$ ). Starting with largest place values, fill in coefficients that sum up to desired decimal value without going over.

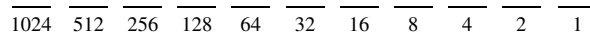
$$75_{10} = \frac{\quad}{256} \frac{\quad}{16} \frac{\quad}{1} \text{ hex}$$

The 2<sup>n</sup> rule

**UNIQUE COMBINATIONS**

## Powers of 2

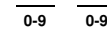
- $2^0 = 1$
- $2^1 = 2$
- $2^2 = 4$
- $2^3 = 8$
- $2^4 = 16$
- $2^5 = 32$
- $2^6 = 64$
- $2^7 = 128$
- $2^8 = 256$
- $2^9 = 512$
- $2^{10} = 1024$



## Unique Combinations

- Given  $n$  digits of base  $r$ , how many unique numbers can be formed?  $r^n$ 
  - What is the range?  $[0 \text{ to } r^n - 1]$

2-digit, decimal numbers ( $r=10, n=2$ )



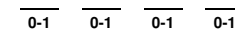
100 combinations:  
00-99

3-digit, decimal numbers ( $r=10, n=3$ )



1000 combinations:  
000-999

4-bit, binary numbers ( $r=2, n=4$ )



16 combinations:  
0000-1111

6-bit, binary numbers  
( $r=2, n=6$ )



64 combinations:  
000000-111111

Main Point: Given  $n$  digits of base  $r$ ,  $r^n$  unique numbers can be made with the range  $[0 - (r^n - 1)]$

## Range of C Data Types

- For a given integer data type we can find its range by raising 2 to the  $n$ ,  $2^n$  (where  $n$  = number of bits of the type)
  - For signed representations we break the range in half with half negative and half positive (0 is considered a positive number by common integer convention)

Bytes	Bits	Type	Unsigned Range	Signed Range
1	8	[unsigned] char	0 to 255	-128 to +127
2	16	[unsigned] short	0 to 65535	-32768 to +32767
4	32	[unsigned] int	0 to 4,294,967,295	-2,147,483,648 to +2,147,483,648
8	8	[unsigned] long long	0 to 18,446,744,073,709,551,615	-9,223,372,036,854,775,808 to +9,223,372,036,854,775,807
4 (8)	32 (64)	char*	0 to 18,446,744,073,709,551,615	

- How will I ever remember those ranges?
  - I wish I had an easy way to approximate those large numbers!

## Approximating Large Powers of 2

- Often need to find decimal approximation of a large powers of 2 like  $2^{16}$ ,  $2^{32}$ , etc.
- Use following approximations:
  - $2^{10} \approx$  \_\_\_\_\_
  - $2^{20} \approx$  \_\_\_\_\_
  - $2^{30} \approx$  \_\_\_\_\_
  - $2^{40} \approx$  \_\_\_\_\_
- For other powers of 2, decompose into product of  $2^{10}$  or  $2^{20}$  or  $2^{30}$  and a power of 2 that is less than  $2^{10}$

$$2^{16} = 2^6 * 2^{10} \approx$$

$$2^{24} = \approx$$

$$2^{28} = \approx$$

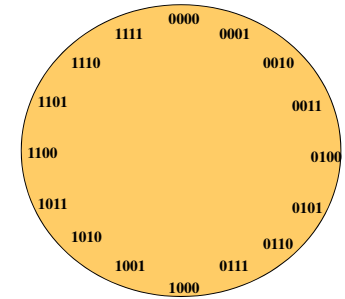
$$2^{32} = \approx$$

- 16-bit word: \_\_\_\_\_ numbers
- 32-bit dword: \_\_\_\_\_ numbers
- 64-bit qword: \_\_\_\_\_ million trillion numbers

## CONVERTING SIGNED NUMBERS TO DECIMAL

## Signed numbers

- Systems used to represent signed numbers split the possible binary combinations in half (half for positive numbers / half for negative numbers)
- Generally, positive and negative numbers are separated using the MSB
  - \_\_\_\_\_ means negative
  - \_\_\_\_\_ means positive



## 2's Complement System

- Normal binary place values except MSB has \_\_\_\_\_

– MSB of 1 = \_\_\_\_\_

4-bit  
Unsigned

Bit 3	Bit 2	Bit 1	Bit 0
8	4	2	1



0 to 15

4-bit  
2's complement

Bit 3	Bit 2	Bit 1	Bit 0
-8	4	2	1



-8 to +7

8-bit  
2's complement

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
-128	64	32	16	8	4	2	1



-128 to +127

## 2's Complement Examples

	1	0	1	1	= -5	
4-bit	-8	4	2	1		
2's complement	0	0	1	1	= +3	
	-8	4	2	1		
	1	1	1	1	= -1	
	-8	4	2	1		

Notice that +3 in 2's comp. is the same as in the unsigned system

	1	0	0	0	0	0	0	1	= -127
8-bit	-128	64	32	16	8	4	2	1	
2's complement	0	0	0	1	1	0	0	1	= +25
	-128	64	32	16	8	4	2	1	

Important: Positive numbers have the \_\_\_\_\_ representation in 2's complement as in normal unsigned binary

## 2's Complement Range

- Given n bits...
    - Max positive value = \_\_\_\_\_
      - Includes all n-1 positive place values
    - Max negative value = \_\_\_\_\_
      - Includes only the negative MSB place value
- Range with n-bits of 2's complement  
 $[-2^{n-1} \text{ to } +2^{n-1}-1]$
- Side note – What decimal value is 111...11?

## Unsigned and Signed Variables

- In C, unsigned variables use unsigned binary (normal power-of-2 place values) to represent numbers

$$\begin{array}{cccccccc} 1 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ \hline 128 & 64 & 32 & 16 & 8 & 4 & 2 & 1 \end{array} = +147$$

- In C, signed variables use the 2's complement system (Neg. MSB weight) to represent numbers

$$\begin{array}{cccccccc} 1 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ \hline -128 & 64 & 32 & 16 & 8 & 4 & 2 & 1 \end{array} = -109$$

## IMPORTANT NOTE

- All computer systems use the **2's complement system** to represent **signed integers**!
- So from now on, if we say an integer is **signed**, we are actually saying it uses the **2's complement system** unless otherwise specified
  - Other systems like "signed magnitude" or "1's complement" exist but will not be used for integers

## Zero and Sign Extension

- Extension is the process of increasing the number of bits used to represent a number without changing its value

Unsigned = Zero Extension (Always add leading 0's):

$$111011 = \underline{\quad} 111011$$

↑ Increase a 6-bit number to 8-bit number by zero extending

2's complement = Sign Extension (Replicate sign bit):

$$\begin{array}{l} \text{pos. } 011010 = \underline{\quad} 011010 \\ \text{neg. } 110011 = \underline{\quad} 110011 \end{array}$$

\_\_\_ bit is just repeated as many times as necessary



## Zero and Sign Truncation

- Truncation is the process of decreasing the number of bits used to represent a number without changing its value

Unsigned = Zero Truncation (Remove leading 0's):

~~00~~111011 = 111011 Decrease an 8-bit number to 6-bit number by truncating 0's. Can't remove a '1' because value is changed

2's complement = Sign Truncation (Remove \_\_\_\_\_ of sign bit):

pos. ~~00~~011010 = Any copies of the MSB can be removed without changing the numbers value. Be careful not to change the sign by cutting off ALL the sign bits.

neg. ~~11~~10011 =

Shortcuts for Converting Binary to Hexadecimal

## SHORTHAND FOR BINARY

## Binary and Hexadecimal

- Hex is base 16 which is  $2^4$
- 1 Hex digit ( ? )<sub>16</sub> can represent: \_\_\_\_\_
- 4 bits of binary ( ? ? ? ? )<sub>2</sub> can represent: \_\_\_\_\_
- Conclusion...  
\_\_ Hex digit = \_\_ bits

## Binary to Hex

- Make groups of 4 bits starting from radix point and working outward
- Add 0's where necessary
- Convert each group of 4 to an octal digit

101001110.11

1101011.101

## Hex to Binary

- Expand each hex digit to a group of 4 bits

14E.C<sub>16</sub>

D93.8<sub>16</sub>

## Hexadecimal Representation

- Since values in modern computers are many bits, we use hexadecimal as a shorthand notation (4 bits = 1 hex digit)
  - 11010010 = D2 hex or **0xD2** if you write it in C/C++
  - 0111011011001011 = 76CB hex or **0x76CB** if you write it in C/C++

## Interpreting Hex Strings

- What does the following hexadecimal represent?
- Just like binary, you must know the underlying *representation system* being used before you can interpret a hex value
- Information (value) = Hex + Context (System)
  - For now, best bet is to convert to \_\_\_\_\_, then translate

0x41 = ?

Unsigned Binary system



65 decimal

x86 Assembly Instruction



inc %ecx

(Add 1 to the ecx register)

ASCII system



'A'<sub>ASCII</sub>

## Hexadecimal & Sign

- If a number is represented in 2's complement (e.g. 10010110) then the MSB of its binary representation would correspond to:
  - 0 = Positive
  - 1 = Negative
- If that same 2's complement number were viewed as hex (e.g. 0x96) how could we tell if the corresponding number is positive or negative?
  - MSD of 0-7 = Positive
  - MSD of 8-F = Negative

Hex – Binary – Sign

0 = 0000 = Pos.
1 = 0001 = Pos.
2 = 0010 = Pos.
3 = 0011 = Pos.
4 = 0100 = Pos.
5 = 0101 = Pos.
6 = 0110 = Pos.
7 = 0111 = Pos.
8 = 1000 = Neg.
9 = 1001 = Neg.
A = 1010 = Neg.
B = 1011 = Neg.
C = 1100 = Neg.
D = 1101 = Neg.
E = 1110 = Neg.
F = 1111 = Neg.

Implicit and Explicit

## APPLICATION: CASTING

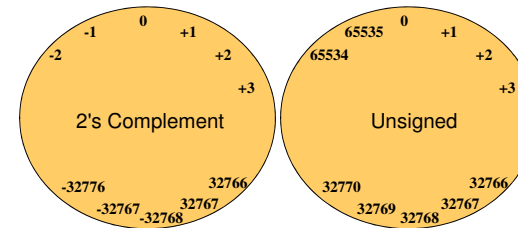
## Implicit and Explicit Casting

- Use your understanding of unsigned and 2's complement to predict the output
- Notes:  $2^{16} = 65536$ 
  - unsigned short range: 0 to 65535
  - signed short range: -32768 to +32768

```
int main()
{
    short int v = -10000; /* 0xd8f0 */
    unsigned short uv = (unsigned short) v;
    printf("v = %d, uv = %u\n", v, uv);
    return 0;
}
```

Expected Output:

v = -10000, uv = \_\_\_\_\_



```
int main()
{
    unsigned u = 4294967295u; /* UMax */
    int tu = (int) u;
    printf("u = %u, tu = %d\n", u, tu);
    return 0;
}
```

Expected Output:

u = 4294967295, tu = \_\_\_\_\_

## Implicit and Explicit Casting

- Use your understanding of zero and sign extension to predict the output

```
int main()
{
    short int v = 0xcfc7; /* -12345 */
    unsigned short uv = 0xcfc7; /* 53191 */
    int vi = v; /* ??? */
    unsigned uvi = uv; /* ??? */
    printf("vi = %x, uvi = %x\n", vi, uvi);
    return 0;
}
```

Expected Output:

vi = ffff cfc7, uvi = \_\_\_\_\_

```
int main()
{
    int x = 53191; /* 0xcfc7 */
    short sx = x;
    int y = sx;
    char z = x;

    printf("sx = %d, y = %d ", sx, y);
    printf("z = %d\n", z);
    return 0;
}
```

Expected Output:

sx = -12345, y = -12345, z = \_\_\_\_\_

## Advice

- Casting can be done implicitly and explicitly
- Casting from one system to another applies a new "interpretation" (pair of glasses) on the same bits
- Casting from one size to another will perform extension or truncation (based on the system)